Coloring finite subsets of uncountable sets

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Abstract

It is consistent for every $1 \leq n < \omega$ that $2^{\omega} = \omega_n$ and there is a function $F : [\omega_n]^{<\omega} \to \omega$ such that every finite set can be written at most $2^n - 1$ ways as the union of two distinct monocolored sets. If GCH holds, for every such coloring there is a finite set that can be written at least $\frac{1}{2} \sum_{i=1}^{n} {n+i \choose i} {n \choose i}$ ways as the union of two sets with the same color.

0 Introduction

In [6] we proved that for every coloring $F : [\omega_n]^{<\omega} \to \omega$ there exists a set $A \in [\omega_n]^{<\omega}$ which can be written at least $2^n - 1$ ways as $A = H_0 \cup H_1$ for some $H_0 \neq H_1$, $F(H_0) = F(H_1)$ and that for n = 1 there is in fact a function F for which this is sharp. Here we show that for every $n < \omega$ it is consistent that $2^{\omega} = \omega_n$ and for some function F as above for every finite set A there are at most $2^n - 1$ solutions of the above equation. We use historic forcing which was first used in [1] and [7] then in [5] and [4]. Under GCH, we improve the positive result of [6] by showing that for every F as above some finite set can be written at least $T_n = \frac{1}{2} \sum_{i=1}^n {n+i \choose i} {n \choose i}$ ways as the union of two sets with the same F value.

With the methods of [6] it is easy to show the following corollary of our independence result. It is consistent that $2^{\omega} = \omega_n$ and there is a function $f: \mathbf{R} \to \omega$ such that if x is a real number then x cannot be written more than $2^n - 1$ ways as the arithmetic mean of some $y \neq z$ with f(y) = f(z). ((y, z) and (z, y) are not regarded distinct.) Another idea of [6] can be used to modify our second result to the following. If GCH holds and V is a vector space over the rationals with $|V| = \omega_n$, $f: V \to \omega$ then some vector can be written at least T_n ways as the arithmetic mean of two vectors with the same f-value.

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Notation We use the standard set theory notation. If S is a set, κ a cardinal, then $[S]^{\kappa} = \{A \subseteq S : |A| = \kappa\}, [S]^{<\kappa} = \{A \subseteq S : |A| < \kappa\}, [S]^{\leq\kappa} = \{A \subseteq S : |A| < \kappa\}, [S]^{\leq\kappa} = \{A \subseteq S : |A| \leq \kappa\}.$ P(S) is the power set of S. If f is a function, A a set, then $f[A] = \{f(x) : x \in A\}.$

1 The independence result

Theorem 1 For $1 \le n < \omega$ it is consistent that $2^{\omega} = \omega_n$ and there is a function $F : [\omega_n]^{<\omega} \to \omega$ such that for every $A \in [\omega_n]^{<\omega}$ there are at most $2^n - 1$ solutions of $A = H_0 \cup H_1$ with $H_0 \neq H_1$, $F(H_0) = F(H_1)$.

For $\alpha < \omega_n$ fix a bijection $\varphi_\alpha : \alpha \to |\alpha|$. For $x \in [\omega_n]^{<\omega}$ define $\gamma_i(x)$ for $i < k = \min(n, |x|)$ as follows. $\gamma_0(x) = \max(x)$.

$$\gamma_{i+1}(x) = \varphi_{\gamma_0(x)}^{-1} \Big(\gamma_i \big(\varphi_{\gamma_0(x)}[x \cap \gamma_0(x)] \big) \Big).$$

 $\gamma(x) = \{\gamma_0(x), \dots, \gamma_{k-1}(x)\}.$

So, for example, if n = 0 then $\gamma(x) = \emptyset$, if n = 1, $x \neq \emptyset$, then $\gamma(x) = \{\gamma_0(x)\} = \{\max(x)\}.$

Lemma 1 Given $s \in [\omega_n]^{\leq n}$ there are at most countably many $x \in [\omega_n]^{<\omega}$ such that $\gamma(x) = s$.

Proof By induction on n.

Let $\Phi(s) = \bigcup \{x : \gamma(x) \subseteq s\}$, a countable set for $s \in [\omega_n]^{<\omega}$.

Definition The two sets $x, y \in [\omega_n]^{<\omega}$ are isomorphic if the structures $(x; < ,\gamma_0(x), \ldots, \gamma_{k-1}(x)), (y; <, \gamma_0(y), \ldots, \gamma_{k-1}(y))$, are isomorphic, i.e., |x| = |y| and the positions of the elements $\gamma_i(x), \gamma_i(y)$ are the same.

Notice that for every finite j there are just finitely many isomorphism types of j-element sets.

The elements of P, the applied notion of forcing will be some structures of the form p = (s, f) where $s \in [\omega_n]^{<\omega}$ and $f : P(s) \to \omega$.

The only element of P_0 is $\mathbf{1}_P = (\emptyset, \langle \emptyset, 0 \rangle)$, it will be the largest element of P. The elements of P_1 are of the form $p = (\{\xi\}, f)$ where $f(\emptyset) = 0 \neq f(\{\xi\})$ for $\xi < \omega_n$.

Given P_t , p = (s, f) is in P_{t+1} if the following is true. $s = \Delta \cup a \cup b$ is a disjoint decomposition. $p' = (\Delta \cup a, f')$ and $p'' = (\Delta \cup b, f'')$ are in P_t where $f' = f | P(\Delta \cup a), f'' = f | P(\Delta \cup b)$. There is $\pi : \Delta \cup a \to \Delta \cup b$, an isomorphism

between $(\Delta \cup a, <, P(\Delta \cup a), f')$ and $(\Delta \cup b, <, P(\Delta \cup b), f')$. $\pi | \Delta$ is the identity. For $H \subseteq \Delta \cup a$ the sets H and $\pi[H]$ are isomorphic. $a \cap \Phi(\Delta) = b \cap \Phi(\Delta) = \emptyset$. f - f' - f'' is one-to-one and takes only values outside $\operatorname{Ran}(f')$ (which is the same as $\operatorname{Ran}(f'')$). $P = \bigcup \{P_t : t < \omega\}$. We make $p \leq p', p''$ and the ordering on P is the one generated by this.

Lemma 2 (P, \leq) is ccc.

Proof Assume that $p_{\alpha} \in P$ ($\alpha < \omega_1$). We can assume by thinning and using the Δ -system lemma and the pigeon hole principle that the following hold. $p_{\alpha} \in P_t$ for the same $t < \omega$. $p_{\alpha} = (\Delta \cup a_{\alpha}, <, P(\Delta \cup a_{\alpha}), f_{\alpha})$ where the structures $(\Delta \cup a_{\alpha}, <, f_{\alpha})$ and $(\Delta \cup a_{\beta}, <, f_{\beta})$ are isomorphic for $\alpha, \beta < \omega_1$, $\{\Delta, a_{\alpha} : \alpha < \omega_1\}$ pairwise disjoint. We can also assume that if π is the isomorphism between $(\Delta \cup a_{\alpha}, <, f_{\alpha})$ and $(\Delta \cup a_{\beta}, <, f_{\beta})$ then H and $\pi[H]$ are isomorphic for $H \subseteq \Delta \cup a_{\alpha}$. Moreover, if we assume that Δ occupies the same positions in the ordered sets $\Delta \cup a_{\alpha}$ ($\alpha < \omega_1$) then π will be the identity on Δ . As $\Phi(\Delta)$ is countable, by removing countably many indices we can also assume that $\Phi(\Delta) \cap a_{\alpha} = \emptyset$ for $\alpha < \omega_1$. Now any p_{α} and p_{β} are compatible as we can take $p = (\Delta \cup a_{\alpha} \cup a_{\beta}, <, P(\Delta \cup a_{\alpha} \cup a_{\beta}), f) \leq p_{\alpha}, p_{\beta}$ where $f \supseteq f_{\alpha}, f_{\beta}$ is an appropriate extension, i.e., $f - f_{\alpha} - f_{\beta}$ is one-to-one and takes values outside $\operatorname{Ran}(f_{\alpha})$.

Lemma 3 If $(s, f) \in P$, $H_0, H_1 \subseteq s$ have $f(H_0) = f(H_1)$ then H_0, H_1 are isomorphic.

Proof Set $(s, f) \in P_t$. We prove the statement by induction on t. There is nothing to prove for t < 2. Assume now that $(s, f) \in P_{t+1}$, $s = \Delta \cup a \cup b$, $\pi : \Delta \cup a \to \Delta \cup b$ as in the definition of (P, \leq) . As $f(H_0)$ is a value taken twice by f, both H_0 and H_1 must be subsets of either $\Delta \cup a$ or $\Delta \cup b$. We are done by induction unless $H_0 \subseteq \Delta \cup a$ and $H_1 \subseteq \Delta \cup b$ (or vice versa). Now H_0 and $\pi[H_0]$ are isomorphic and $f(H_0) = f(\pi[H_0]) = f(H_1)$ so by the inductive hypothesis $\pi[H_0]$ and H_1 are isomorphic and then so are H_0 , H_1 .

Lemma 4 If $(s, f) \in P$, $H_0, H_1 \subseteq s$, $f(H_0) = f(H_1)$, $x \in H_0 \cap H_1$ then x occupies the same position in the ordered sets H_0, H_1 .

Proof Similarly to the proof of the previous Lemma, by induction on t, for $(s, f) \in P_t$. With similar steps, we can assume that $(s, f) = (\Delta \cup a \cup b, f) \leq (\Delta \cup a, f'), (\Delta \cup b, f''), H_0 \subseteq \Delta \cup a, H_1 \subseteq \Delta \cup b$. Notice that $x \in \Delta$. Now, as $\pi(x) = x, x$ is a common element of $\pi[H_0]$ and H_1 and also $f''(\pi[H_0]) = f''(H_1)$.

By induction we get that x occupies the same position in $\pi[H_0]$ and H_1 so by pulling back we get that this is true for H_0 and H_1 .

Lemma 5 If $(s, f) \in P$, $A \subseteq s$, $0 \leq j \leq n$ then A can be written at most $2^j - 1$ ways as $A = H_0 \cup H_1$ with H_0 , H_1 distinct, $f(H_0) = f(H_1)$, and $|\gamma(H_0) \cap \gamma(H_1)| \geq n - j$.

Proof By induction on j and inside that induction, by induction on t, for $(s, f) \in P_t$. The case t < 2 will always be trivial.

Assume first that j = 0. In this case our Lemma reduces to the following statement. There are no $H_0 \neq H_1$ such that $\gamma(H_0) = \gamma(H_1)$. In the inductive argument we assume as usual that $s = \Delta \cup a \cup b$ and so $(s, f) \in P_{t+1}$ was created from $(\Delta \cup a, f')$ and $(\Delta \cup b, f'')$, $H_0 \subseteq \Delta \cup a$, $H_1 \subseteq \Delta \cup b$. As $\gamma(H_0) = \gamma(H_1)$, $\gamma(H_0) \subseteq \Delta$, but then, as $\Phi(\Delta) \cap a = \emptyset$, H_0 can have no points outside Δ and similarly for H_1 , so we can go back, say to $(\Delta \cup a, f') \in P_t$ which concludes the argument.

Assume now that the statement is proved for j and we have $p = (s, f) \in P_{t+1}$, $s = \Delta \cup a \cup b$ and p was created from $p' = (\Delta \cup a, f')$ and $p'' = (\Delta \cup b, f'')$. In $A \subseteq \Delta \cup a \cup b$ we can assume that $y = A \cap a \neq \emptyset$, $z = A \cap b \neq \emptyset$ as otherwise we can pull back to p' or p''. But then, if $A = H_0 \cup H_1$, then, if, say, $H_0 \subseteq \Delta \cup a$, $H_1 \subseteq \Delta \cup b$ hold, then necessarily $H_0 \cap a = y$, $H_1 \cap b = z$, so $H_0 = x_0 \cup y$, $H_1 = x_1 \cup z$ where $x_0 \cup x_1 = x = A \cap \Delta$. We can create decompositions of $B = x \cup \pi[y] \cup z$ by taking $B = \pi[H_0] \cup H_1$. But some of these decompositions will not be different and it may happen that we get non-proper (i.e., one-piece) decomposition. This can only happen if $\pi[y] = z$, and then the two decompositions $A = (x_0 \cup y) \cup (x_1 \cup z)$ and $A = (x_1 \cup y) \cup (x_0 \cup y)$ produce the same decomposition of B, namely, $B = (x_0 \cup z) \cup (x_1 \cup z)$ and there is but one decomposition, $A = (x \cup y) \cup (x \cup z)$ which cannot be mapped to a decomposition of B. If this (i.e., $\pi[y] = z$) does not happen, we are done by induction. If this does happen, we know that $\gamma(H_0) = \gamma(x_0 \cup y)$ has an element in y (by the argument at the beginning of the proof). As $f(x_0 \cup y) = f(x_1 \cup z)$, by Lemmas 3 and 4, both $H_0 = x_0 \cup y$ and $H_1 = x_1 \cup z$ have an element in the γ -subset, at the same positions which are mapped onto each other by π . We get that $\gamma(x_0 \cup z) \cap \gamma(x_1 \cup z)$ has at least n-j element, so by our inductive assumption we have at most $2^{j} - 1$ decompositions, which gives at most $2 \cdot (2^j - 1) + 1 = 2^{j+1} - 1$ decompositions of A.

Let $G \subseteq P$ be a generic subset. Set $S = \bigcup \{s : (s, f) \in G\}, F = \bigcup \{f : (s, f) \in G\}.$

Lemma 6 There is a $p \in P$ such that $p \parallel - \mid S \mid = \aleph_n$.

Proof Otherwise $\mathbf{1} \parallel - \sup(S) < \omega_n$. By ccc, there is an ordinal $\xi < \omega_n$ for which $\mathbf{1} \parallel - \sup(S) < \xi$, but this is impossible as there are conditions in P_1 forcing that $\xi \in S$.

Now we can conclude the proof of the Theorem. If G is generic, and $p \in G$ with the condition p of Lemma 6, then in V[G] F witnesses the theorem by Lemma 5 (for j = n) on the ground set S. As $|S| = \omega_n$ we can replace it by ω_n . \Box

2 The GCH result

Set

$$T_n = \frac{1}{2} \sum_{i=1}^n \binom{n+i}{n} \binom{n}{i}.$$

So $T_1 = 1$, $T_2 = 6$, $T_3 = 31$. In general, T_n is asymptotically $c(3 + 2\sqrt{2})^n / \sqrt{n}$ for some c.

Theorem 2 (GCH) If $F : [\omega_n]^{<\omega} \to \omega$ then some $A \in [\omega_n]^{<\omega}$ has at least T_n decompositions as $A = H_0 \cup H_1$, $H_0 \neq H_1$, $F(H_0) = F(H_1)$.

Proof By the Erdős-Rado theorem (see [2, 3]) there is a set $\{x_{\alpha} : \alpha < \omega_1\}$ which is (n-1)-end-homogeneous, i.e., for some $g : [\omega_1]^{<\omega} \to \omega$, if $\alpha_1 < \cdots < \alpha_k < \beta_1 < \cdots < \beta_{n-1} < \omega_1$ then

$$f(\{x_{\alpha_1},\ldots,x_{\alpha_k},x_{\beta_1},\ldots,x_{\beta_{n-1}}\}) = g(\alpha_1,\ldots,\alpha_k).$$

Select $S_1 \in [\omega_1]^{\omega_1}$ in such a way that $g(\alpha) = c_0$ for $\alpha \in S_1$. Set $\gamma_1 = \min(S_1)$. In general, if γ_i , S_i are given $(1 \le i < n)$ pick $S_{i+1} \in [S_i - (\gamma_i + 1)]^{\omega_1}$ so that $g(\gamma_1, \ldots, \gamma_i, \alpha) = c_i$ for $\alpha \in S_{i+1}$ and set $\gamma_{i+1} = \min(S_{i+1})$. Given $\gamma_1, \ldots, \gamma_n$ and S_n let $\gamma_{n+1}, \ldots, \gamma_{2n}$ be the *n* least elements of $S_n - (\gamma_n + 1)$.

Our set will be $A = \{x_{\gamma_1}, \ldots, x_{\gamma_{2n}}\}$. For $0 \le i < n$ the color of any (n + i)element subset of A containing $x_{\gamma_1}, \ldots, x_{\gamma_i}$ will be c_i . We can select $\frac{1}{2} \binom{2n-i}{n} \binom{n}{i}$ different pairs of those sets which cover A. In toto, we get T_n decompositions
of A.

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