LORDS OF THE ITERATION

ANDRZEJ ROSLANOWSKI AND SAHARON SHELAH

Abstract. We introduce several properties of forcing notions which imply that their \( \lambda \)–support iterations are \( \lambda \)–proper. Our methods and techniques refine those studied in [RS01], [RS07], [RS05] and [RS], covering some new forcing notions (though the exact relation of the new properties to the old ones remains undecided).

0. Introduction

Since the beginning of 1980s it has been known that the theory of proper forcing does not admit naive generalization to the context of larger cardinals and iterations with larger supports. The evidence of that was given already in Shelah [She82] (see [She98, Appendix 3.6(2)]). It seems that the first steps towards developing the theory of forcing iterated with uncountable supports were done in Shelah [She03a], [She03b], but the properties introduced there were aimed at situations when we do not want to add new subsets of \( \lambda \) (corresponding to the case of no new reals in CS iterations of proper forcing notions). Later Roslanowski and Shelah [RS01] introduced an iterable property called properness over semi-diamonds and then Eisworth [Eis03] proposed an iterable relative of it. These properties work nicely for \( \lambda \)–support iterations (where \( \lambda = \lambda^{<\lambda} \) is essentially arbitrary) and forcings adding new subsets of \( \lambda \), but the price to pay is that many natural examples are not covered. If we restrict ourselves to inaccessible \( \lambda \), then the properties given by Roslanowski and Shelah [RS07, RS05, RS] may occur useful. Those papers give both iteration theorems and new examples of forcing notions for which the theorems apply.

In the present paper we further advance the theory and we give results applicable to both the case of inaccessible \( \lambda \) as well as those working for successor cardinals. The tools developed here may be treated as yet another step towards comparing and contrasting the structure of \( \lambda^{\lambda} \) with that of \( \omega^{\omega} \). That line of research already has received some attention in the literature (see e.g., Cummings and Shelah [CS95], Shelah and Spasojević [SS02] or Zapletal [Zap97]). Also with better iteration theorems one may hope for further generalizations of Roslanowski and Shelah [RS99] to the context of uncountable cardinals. (Initial steps in the latter direction were presented in Roslanowski and Shelah [RS07].) However, while we do give some examples of forcing notions to which our properties apply, we concentrate on the
development of the theory of forcing leaving the real applications for further investigations. The need for the development of such general theory was indirectly stated by Hytinnen and Rautila in [HR01], where they commented:

Our proof is longer than the one in [MS93] partly because we are not able to utilize the general theory of proper forcing, especially the iteration lemma, but we have to prove everything “from scratch”.

We believe that the present paper brings us substantially closer to the right general iteration theorems for iterations with uncountable supports.

In the first section we introduce $D\ell$–parameters (which will play an important role in our definitions) and a slight generalization of the B–bounding property from [RS05]. We also define a canonical example for testing usefulness of our iteration theorems: the forcing $Q^E_E$ in which conditions are complete $\lambda$–trees in which along each $\lambda$-branch the set of splittings forms a set from a filter $E$ (and the splitting at $\nu$ is into a set from a filter $E_\nu$ on $\lambda$). The main result of the first section (Theorem 1.10) says that we may iterate with $\lambda$–supports forcing notions $Q^E_E$, provided $\lambda$ is inaccessible and $E$ is always the same and has some additional properties.

If we want to iterate forcing notions like $Q^E_E$ but with different $E$ on each coordinate (when the result of the first section is not applicable), we may decide to use very orthogonal filters. Section 2 presents an iteration theorem 2.7 which is tailored for such situation. Also here we need the assumption that $\lambda$ is inaccessible.

The following section introduces $B$–noble forcing notions and the iteration theorem 3.3 for them. The main gain here is that it allows us to iterate (with $\lambda$–supports) forcing notions like $Q^E_E$ even if $\lambda$ is not inaccessible. The fourth section gives more examples of forcing notions and shows a possible application. In Corollary 4.5 we substantially improve a result from [RS05] showing that dominating numbers associated with different filters may be distinct even if $\lambda$ is a successor.

The fifth section shows that some of closely related forcing notions may have different properties. Section 6 presents yet another property that is useful in $\lambda$–support iterations (for inaccessible $\lambda$): reasonably merry forcing notions. This property has the flavour of putting together being $B$–bounding (of [RS05]) with being fuzzy proper (of [RS07]). We also give an example of a forcing notion which is reasonably merry but which was not covered by earlier properties. We conclude the paper with a section listing open problems.

This research is a natural continuation of papers mentioned earlier ([She03a], [She03b], [RS01], [RS07], [RS05] and [RS]). All our iteration proofs are based on trees of conditions and the arguments are similar to those from the earlier works. While we tried to make this presentation self-contained, the reader familiar with the previous papers will definitely find the proofs presented here easier to follow (as several technical aspects do re-occur).

0.1. Notation. Our notation is rather standard and compatible with that of classical textbooks (like Jech [Jec03]). In forcing we keep the older convention that a stronger condition is the larger one.

(1) Ordinal numbers will be denoted be the lower case initial letters of the Greek alphabet ($\alpha, \beta, \gamma, \delta$ . . .) and also by $i, j$ (with possible sub- and superscripts).
Cardinal numbers will be called $\kappa, \lambda, \mu$; $\lambda$ will be always assumed to be a regular uncountable cardinal such that $\lambda^{<\lambda} = \lambda$ (we may forget to mention this).

Also, $\chi$ will denote a sufficiently large regular cardinal; $\mathcal{H}(\chi)$ is the family of all sets hereditarily of size less than $\chi$. Moreover, we fix a well ordering $<_\chi$ of $\mathcal{H}(\chi)$.

(2) We will consider several games of two players. One player will be called *Generic* or *Complete* or just *COM*, and we will refer to this player as “she”. Her opponent will be called *Antigeneric* or *Incomplete* or just *INC* and will be referred to as “he”.

(3) For a forcing notion $\mathbb{P}$, almost all $\mathbb{P}$–names for objects in the extension via $\mathbb{P}$ will be denoted with a tilde below (e.g., $\tau, X$). There will be some exceptions to this rule, however. $\Gamma_\mathbb{P}$ will stand for the canonical $\mathbb{P}$–name for the generic filter in $\mathbb{P}$. Also some (names for) normal filters generated in the extension from objects in the ground model will be denoted by $D$, $D^\mathbb{P}$ or $D[\mathbb{P}]$.

The weakest element of $\mathbb{P}$ will be denoted by $\emptyset_\mathbb{P}$ (and we will always assume that there is one, and that there is no other condition equivalent to it). All forcing notions under considerations are assumed to be atomless.

By “$\lambda$–support iterations” we mean iterations in which domains of conditions are of size $\leq \lambda$. However, on some occasions we will pretend that conditions in a $\lambda$–support iteration $\mathbb{Q} = (P_\zeta, Q_\zeta : \zeta < \zeta^*)$ are total functions on $\zeta^*$ and for $p \in \text{lim}(\mathbb{Q})$ and $\alpha \in \zeta^* \setminus \text{dom}(p)$ we will let $p(\alpha) = \emptyset_{\mathbb{Q}_\alpha}$.

(4) By “a sequence” we mean “a function defined on a set of ordinals” (so the domain of a sequence does not have to be an ordinal). For two sequences $\zeta, \nu$ we write $\nu \triangleleft \eta$ whenever $\nu$ is a proper initial segment of $\eta$, and $\nu \triangleleft \eta$ when either $\nu \triangleleft \eta$ or $\nu = \eta$. The length of a sequence $\eta$ is the order type of its domain and it is denoted by $\text{lh}(\eta)$.

(5) A tree is a $\triangleleft$–downward closed set of sequences. A complete $\lambda$–tree is a tree $T \subseteq {}^{<\lambda}\lambda$ such that every $\triangleleft$–chain of size less than $\lambda$ has an $\triangleleft$–bound in $T$ and for each $\eta \in T$ there is $\nu \in T$ such that $\eta \triangleleft \nu$.

Let $T$ be a $\lambda$–tree. For $\eta \in T$ we let

$$\text{suc}_T(\eta) = \{\alpha : \lambda^\frown(\alpha) \in T\} \quad \text{and} \quad (T)_\eta = \{\nu \in T : \nu \triangleleft \eta \text{ or } \eta \triangleleft \nu\}.$$

We also let root($T$) be the shortest $\eta \in T$ such that $|\text{suc}_T(\eta)| > 1$ and $\text{lim}_\lambda(T) = \{\eta \in {}^{<\lambda}\lambda : (\forall \alpha < \lambda)(\eta|\alpha \in T)\}$.

### 0.2. Background on trees of conditions.

**Definition 0.1.** Let $\mathbb{P}$ be a forcing notion.

1. For a condition $r \in \mathbb{P}$, let $\mathcal{O}^\lambda_\mathbb{P}(\mathbb{P}, r)$ be the following game of two players, Complete and Incomplete:

   - the game lasts at most $\lambda$ moves and during a play the players construct a sequence $\langle(p_i, q_i) : i < \lambda\rangle$ of pairs of conditions from $\mathbb{P}$ in such a way that $(\forall j < i < \lambda)(r \leq p_j \leq q_j \leq p_i)$ and at the stage $i < \lambda$ of the game, first Incomplete chooses $p_i$ and then Complete chooses $q_i$.

   Complete wins if and only if for every $i < \lambda$ there are legal moves for both players.
(2) We say that the forcing notion \( \mathbb{P} \) is strategically \( (<\lambda)\)-complete if Complete has a winning strategy in the game \( \tilde{\omega}_0^\lambda(\mathbb{P}, r) \) for each condition \( r \in \mathbb{P} \).

(3) Let \( N < (\mathcal{H}(\chi), \in, <_\chi) \) be a model such that \( <^\lambda N \subseteq N, |N| = \lambda \) and \( \mathbb{P} \in N \). We say that a condition \( p \in \mathbb{P} \) is \( (N, \mathbb{P}) \)-generic in the standard sense (or just: \( (N, \mathbb{P}) \)-generic) if for every \( \mathbb{P} \)-name \( \tau \in N \) for an ordinal we have \( p \Vdash \tau \in N \).”.

(4) \( \mathbb{P} \) is \( \lambda \)-proper in the standard sense (or just: \( \lambda \)-proper) if there is \( x \in \mathcal{H}(\chi) \) such that for every model \( N < (\mathcal{H}(\chi), \in, <_\chi) \) satisfying

\[ <^\lambda N \subseteq N, \quad |N| = \lambda \quad \text{and} \quad \mathbb{P}, x \in N, \]

and every condition \( q \in N \cap \mathbb{P} \) there is an \( (N, \mathbb{P}) \)-generic condition \( p \in \mathbb{P} \) stronger than \( q \).

Remark 0.2. Let us recall that if \( \mathbb{P} \) is either strategically \( (<\lambda^+) \)-complete or \( \lambda^+ \)-cc, then \( \mathbb{P} \) is \( \lambda \)-proper. Also, if \( \mathbb{P} \) is \( \lambda \)-proper then

- \( \lambda^+ \) is not collapsed in forcing by \( \mathbb{P} \), moreover
- for every set of ordinals \( A \in V^\mathbb{P} \) of size \( \lambda \) there is a set \( A^+ \in V \) of size \( \lambda \) such that \( A \subseteq A^+ \).

Definition 0.3 (Compare [RS07, Def. A.1.7], see also [RS05, Def. 2.2]).

(1) Let \( \gamma \) be an ordinal, \( \emptyset \neq w \subseteq \gamma \). A \( (w, 1)^\gamma \)-tree is a pair \( T = (T, \text{rk}) \) such that

- \( \text{rk} : T \rightarrow w \cup \{\gamma\} \),
- if \( t \in T \) and \( \text{rk}(t) = \varepsilon \), then \( t \) is a sequence \( \langle (t)_\zeta : \zeta \in w \cap \varepsilon \rangle \),
- \( (T, <) \) is a tree with root \( \langle \rangle \) and
- if \( t \in T \), then there is \( t' \in T \) such that \( t \leq t' \) and \( \text{rk}(t') = \gamma \).

(2) If, additionally, \( T = (T, \text{rk}) \) is such that every chain in \( T \) has a \(<\)-upper bound in \( T \), we will call it a standard \( (w, 1)^\gamma \)-tree.

We will keep the convention that \( T^\gamma_y \) is \( (T^\gamma_y, \text{rk}^\gamma_y) \).

(3) Let \( \bar{Q} = \langle \bar{P}_i, \bar{Q}_i : i < \gamma \rangle \) be a \( \lambda \)-support iteration. A tree of conditions in \( \bar{Q} \) is a system \( \bar{p} = \langle p_t : t \in T \rangle \) such that

- \( (T, \text{rk}) \) is a \( (w, 1)^\gamma \)-tree for some \( w \subseteq \gamma \),
- \( p_t \in P_{\text{rk}(t)} \) for \( t \in T \), and
- if \( s, t \in T \), \( s < t \), then \( p_s = p_t |_{\text{rk}(s)} \).

If, additionally, \( (T, \text{rk}) \) is a standard tree, then \( \bar{p} \) is called a standard tree of conditions.

(4) Let \( \bar{p}^0, \bar{p}^1 \) be trees of conditions in \( \bar{Q} \), \( \bar{p}^i = \langle p^i_t : t \in T \rangle \). We write \( \bar{p}^0 \leq \bar{p}^1 \) whenever for each \( t \in T \) we have \( p^0_t \leq p^1_t \).

Note that our standard trees and trees of conditions are a special case of that [RS07, Def. A.1.7] when \( \alpha = 1 \).

Proposition 0.4. Assume that \( \bar{Q} = \langle \bar{P}_i, \bar{Q}_i : i < \gamma \rangle \) is a \( \lambda \)-support iteration such that for all \( i < \gamma \) we have

\[ \Vdash_{\bar{P}_i} " \bar{Q}_i \ is \ strategically \ (<\lambda)\)-complete \". \]

Suppose that \( \bar{p} = \langle p_t : t \in T \rangle \) is a tree of conditions in \( \bar{Q} \), \( |T| < \lambda \), and \( T \subseteq \bar{P}_\gamma \) is open dense. Then there is a tree of conditions \( \bar{q} = \langle q_t : t \in T \rangle \) such that \( \bar{p} \leq \bar{q} \) and \( (\forall t \in T)(\text{rk}(t) = \gamma \Rightarrow q_t \in T) \).
Proof. This is essentially [RS07, Proposition A.1.9] and the proof there applies here without changes. \qed

1. $\mathcal{D}_\ell$–parameters

In this section we introduce $\mathcal{D}_\ell$–parameters and we use them to get a possible slight improvement of [RS05, Theorem 3.1] (in Theorem 1.10). We also define our canonical testing forcing $Q_{\mathcal{D}_\ell}^P$ to which this result can be applied.

Definition 1.1. (1) A pre-$\mathcal{D}_\ell$–parameter on $\lambda$ is a triple $p = (\dot{P}, S, D) = (\dot{P}^p, S^p, D^p)$ such that
   \begin{itemize}
     \item $D$ is a proper uniform normal filter on $\lambda$, $S \in D$,
     \item $\dot{P} = \langle P_\delta : \delta \in S \rangle$ and $P_\delta \in [\delta]^{<\lambda}$ for each $\delta \in S$.
   \end{itemize}

(2) For a function $f \in \lambda^\lambda$ and a pre-$\mathcal{D}_\ell$–parameter $p = (\dot{P}, S, D)$ we let
   \[ \text{set}^P(f) = \{ \delta \in S : f(\delta) \in P_\delta \}. \]

(3) We say that a pre-$\mathcal{D}_\ell$–parameter $p = (\dot{P}, S, D)$ is a $\mathcal{D}_\ell$–parameter on $\lambda$ if
   \[ \text{set}^P(f) \in D \text{ for every } f \in \lambda^\lambda. \]

Example 1.2. (1) If $\lambda$ is strongly inaccessible, $D$ is the filter generated by club subsets of $\lambda$ and $P_\delta = \delta \delta$, $\dot{P} = \langle P_\delta : \delta < \lambda \rangle$, then $(\dot{P}, \lambda, D)$ is a $\mathcal{D}_\ell$–parameter on $\lambda$.

(2) $\wedge^*_\lambda$ is a statement asserting existence of a $\mathcal{D}_\ell$–parameter with the filter generated by clubs of $\lambda$.

(3) $\wedge^*_\lambda$ implies the existence of a $\mathcal{D}_\ell$–parameter $(\dot{P}, S, D)$ such that $|P_\delta| = |\delta|$.

(4) For more instances of the existence of $\mathcal{D}_\ell$–parameters we refer the reader to Shelah [She00, §3].

Definition 1.3. Let $p$ be a pre-$\mathcal{D}_\ell$–parameter on $\lambda$ and $\mathbb{Q}$ be a forcing notion not collapsing $\lambda$. In $V^\mathbb{Q}$ we define
   \begin{itemize}
     \item $D^p[\mathbb{Q}] = D^p[\mathbb{Q}]$ is the normal filter generated by $D^p \cup \{ \text{set}^P(f) : f \in \lambda^\lambda \}$,
     \item $p[\mathbb{Q}] = (\dot{P}^p, S^p, D^p[\mathbb{Q}]).$
   \end{itemize}

Remark 1.4. If $\mathbb{Q}$ is a strategically $(<\lambda)$–complete forcing notion and $D$ is a (proper) normal filter on $\lambda$, then in $V^\mathbb{Q}$ the normal filter on $\lambda$ generated by $D \cap \mathbb{V}$ is also a proper filter. Abusing notation, we will denote this filter by $D$ (or $D^\mathbb{Q}$). The filter $D^p[\mathbb{Q}]$ can be larger, but it is still a proper filter, provided $p$ is a $\mathcal{D}_\ell$–parameter.

Lemma 1.5. Assume that $p = (\dot{P}, S, D)$ is a $\mathcal{D}_\ell$–parameter on $\lambda$ and $\mathbb{Q}$ is a strategically $(<\lambda)$–complete forcing notion. Then $\Vdash_{\mathbb{Q}} \emptyset \notin D^p[\mathbb{Q}]$. Consequently, $\Vdash_{\mathbb{Q}} " p[\mathbb{Q}] ^{\mathcal{D}_\ell} \text{ is a } \mathcal{D}_\ell \text{–parameter on } \lambda \".$

Proof. Assume that $p \in \mathbb{Q}$ and $A_\delta$ is a $\mathbb{Q}$–name for an element of $D \cap \mathbb{V}$ and $f_\delta$ is a $\mathbb{Q}$–name for an element of $\lambda^\lambda$ (for $\delta < \lambda$). Using the strategic completeness of $\mathbb{Q}$ build a sequence $\langle p_\alpha, A_\delta, f_\delta : \delta < \lambda \rangle$ such that for each $\delta < \lambda$:

\begin{enumerate}
  \item $p_\alpha \in \mathbb{Q}$, $p \leq p_0 \leq p_\alpha \leq p_\delta$ for $\alpha < \delta$,
  \item $A_\delta \in D \cap \mathbb{V}$, $f_\delta \in \lambda^\lambda$ and
  \item $p_\alpha \Vdash_{\mathbb{Q}} " A_\delta = A_\delta "$ and $f_\alpha |\delta = f_\alpha |\delta$ for all $\alpha \leq \delta "$.
\end{enumerate}

Since $p$ is a $\mathcal{D}_\ell$–parameter, we know that $B = \bigtriangleup_{\delta < \lambda} A_\delta \cap \bigtriangleup_{\delta < \lambda} \text{set}^P(f_\delta) \in D$. Let $\delta \in B$. Then

\[ p_\delta \Vdash_{\mathbb{Q}} " \delta \in \bigtriangleup_{\alpha < \lambda} A_\alpha \text{ and } f_\alpha |\delta = f_\alpha |\delta \in P_\delta \text{ for all } \alpha < \delta ". \]
We may find an increasing continuous sequence $\langle \delta < \lambda \rangle$ such that $\mathbb{P} = (\mathbb{P}, S, D)$ is a strategically $\langle \lambda \rangle$–complete forcing notion and $N \prec (\mathcal{H}(\chi), \in, \langle \lambda \rangle)$ is such that $\mathbb{P} \in N$, $|N| = \lambda$ and $\langle \lambda \rangle \subseteq N$. Let $\langle N_\delta : \delta < \lambda \rangle$ be an increasing continuous sequence of elementary submodels of $N$ such that $\mathbb{P} \in N_0$, $\delta \subseteq N_\delta$, $P_\delta \subseteq N_{\delta + 1}$, $\langle N_\varepsilon : \varepsilon \leq \delta \rangle \in N_{\delta + 1}$ and $|N_\delta| < \lambda$ (for $\delta < \lambda$). Then
\[ \models_{\mathbb{Q}} \left( \forall A \subseteq N \right) \left( \{ \delta < \lambda : A \cap N_\delta \in N_{\delta + 1} \} \in D[Q] \right). \]

**Proof.** We may find an increasing continuous sequence $\langle \alpha_\delta : \delta < \lambda \rangle \subseteq \lambda$ and a bijection $f : N \rightarrow \lambda$ such that $f|N_\delta = \alpha_\delta$ and $f|N_\delta \in N_{\delta + 1}$ (for $\delta < \lambda$). For $A \subseteq N$ let $\varphi_A : \lambda \rightarrow 2$ be such that $\varphi_A(\alpha) = 1$ if and only if $f^{-1}(\alpha) \in A$. Plainly, if $\delta = \alpha_\delta$ and $\varphi_A|\delta \in P_\delta$, then $A \cap N_\delta \in N_{\delta + 1}$. \( \square \)

Let $\mathbb{C}_0^\lambda$ be a forcing notion consisting of all pairs $(\alpha, f)$ such that $\alpha < \lambda$ and $f \in \prod_{\beta < \alpha} (\beta + 1)$ ordered by the extension (so $(\alpha, f) \leq (\alpha', f')$ if and only if $f \subseteq f'$). Thus it is a $\langle \lambda \rangle$–complete forcing notion which is an incarnation of the $\lambda$–Cohen forcing notion.

**Proposition 1.7.** Assume $\lambda$ is strongly inaccessible. If $\mathbb{P} = (\mathbb{P}, S, D)$ is a $\mathcal{D}$–parameter on $\lambda$ such that $(\forall \delta \in S)(|P_\delta| \leq |\delta|)$, then $\models_{\mathbb{C}_0^\lambda} \left( D[\mathbb{C}_0^\lambda] \neq D[\mathbb{C}_0^\lambda] \right)$. \( \square \)

**Proof.** Let $f$ be the canonical $\mathbb{C}_0^\lambda$–name for the generic function in $\prod_{\alpha < \lambda} (\alpha + 1)$, so $(\alpha, f) \models_{\mathbb{C}_0^\lambda} f \subseteq f$. Plainly, $\models_{\mathbb{C}_0^\lambda} \langle \alpha, f \rangle \models_{\mathbb{C}_0^\lambda} \text{set}^\mathbb{P}(f) \in D[\mathbb{C}_0^\lambda]$ and we are going to argue that $\models_{\mathbb{C}_0^\lambda} \lambda \setminus \text{set}^\mathbb{P}(f) \in (D[\mathbb{C}_0^\lambda])^+$. To this end, suppose that $p \in \mathbb{C}_0^\lambda$ and $\delta_\alpha$ is a $\mathbb{C}_0^\lambda$–name for an element of $D \cap V$ (for $\delta < \lambda$). By induction on $\xi < \lambda$ choose $\langle \alpha_\xi, B_\xi, p_\xi : \xi < \lambda \rangle$ so that

1. $(\alpha) \langle \alpha_\xi : \xi < \lambda \rangle$ is an increasing continuous sequence of ordinals below $\lambda$,
2. $(\beta) B_\xi \subseteq D$, $p_\xi^0 = (p_\xi^0 : \sigma \in \prod_{\zeta < \xi} \alpha_\zeta + 1) \subseteq \mathbb{C}_0^\lambda$, $p_\xi^0 = p = (\alpha_0, f_0^\xi)$,
3. $(\gamma)$ if $\sigma \in \prod_{\zeta < \xi} \alpha_\zeta + 1$, then $p_\xi^\sigma = (\alpha_\xi, f_\xi^\sigma)$ and $f_\xi^\sigma(\alpha_\zeta) = \sigma(\zeta)$ for $\zeta < \xi$,
4. $(\delta)$ if $\xi < \xi'$, $\sigma' \in \prod_{\zeta < \xi'} \alpha_\zeta + 1$ and $\sigma = \sigma'|\xi$, then $p_\xi^\sigma \leq p_{\xi'}^{\sigma'}$,
5. $(\varepsilon)$ if $\xi < \lambda$ is limit, $\sigma \in \prod_{\zeta < \xi} \alpha_\zeta + 1$, then $(\alpha_\xi = \sup(\alpha_\zeta : \zeta < \xi))$ and $f_\xi^\sigma = \bigcup_{\zeta < \xi} f_\xi^{\sigma|\zeta}$, and
6. $(\zeta) p_{\xi + 1}^\sigma \models B_\xi \subseteq A_\xi$ for every $\sigma \in \prod_{\zeta < \xi} \alpha_\zeta + 1$.

(Remember that $\lambda$ is inaccessible, so $|\prod_{\zeta < \xi} \alpha_\zeta + 1| < \lambda$ for each $\xi$.) Next, consider the set $B = \bigcup_{\xi < \lambda} B_\xi \in D$. Let $\delta \in B \cap S$ be a limit ordinal. Since $|P_\delta| \leq |\delta| < \prod_{\xi < \delta} |\alpha_\xi|$, we may pick $\sigma \in \prod_{\zeta < \delta} (\alpha_\zeta + 1)$ such that $f_\xi^\sigma \notin P_\delta$. Then $p_\xi^\sigma \models \delta \in \bigcup_{\xi < \lambda} A_\xi \setminus \text{set}^\mathbb{P}(f)$. \( \square \)

**Definition 1.8.** Let $\mathbb{P} = (\mathbb{P}, S, D)$ be a $\mathcal{D}$–parameter on $\lambda$, $\mathbb{Q}$ be a strategically $\langle \lambda \rangle$–complete forcing notion.
Theorem 1.10. Assume that $b$–bounding (see [RS, Definition 2.9(2),(4)]) and get the parallel iteration theorem. If not reasonably $B$–bounding over $D$ and not reasonably $B$–bounding over $\mathbb{Q}$ however, if there exists a forcing notion which is reasonably $B$–bounding over $\mathbb{Q}$, then $D$ is a normal filter on $\lambda$. We define a forcing notion $(\otimes)_{\mathbb{B}}$ there is a condition $p^* \in Q$ stronger than $p$ and such that $p^* \vdash_{\mathbb{Q}} \{ \alpha < \lambda : (\exists t \in I_\alpha)(q_\alpha^t \in \Gamma_t) \} \in D[\mathbb{Q}]$.

Remark 1.9. The notion introduced in 1.8 is almost the same as the one of [RS05, Definition 3.1(2),(5)]. The difference is that in $(\otimes)_{\mathbb{B}}$ we use the filter $D[\mathbb{Q}]$ and not $D[\mathbb{Q}] = D$, so potentially we have a weaker property here. We do not know, however, if there exists a forcing notion which is reasonably $B$–bounding over $\mathbb{P}$ and not reasonably $B$–bounding over $D$. (See Problem 7.1.)

In a similar fashion we may also modify the property of being nicely double $b$–bounding (see [RS, Definition 2.9(2),(4)]) and get the parallel iteration theorem.

Theorem 1.10. Assume that

1. $\lambda$ is a strongly inaccessible cardinal and $\mathbb{P}$ is a $\mathcal{D}_\ell$–parameter on $\lambda$,
2. $\mathbb{Q} = (\mathbb{P}_\alpha, \mathbb{Q}_\alpha : \alpha < \gamma)$ is a $\lambda$–support iteration,
3. for every $\alpha < \lambda$, $\mathbb{P}_\alpha$ “$\mathbb{Q}_\alpha$ is reasonably $B$–bounding over $\mathbb{P}[\mathbb{P}_\alpha]$.”

Then

(a) $\mathbb{P}_\gamma = \operatorname{lim}(\mathbb{Q})$ is $\lambda$–proper,
(b) if $\tau$ is a $\mathbb{P}_\gamma$–name for a function from $\lambda$ to $V$, $p \in \mathbb{P}_\gamma$, then there are $q \geq p$ and $\langle A_\xi : \xi < \lambda \rangle$ such that $\langle \forall \xi < \lambda \rangle (\langle A_\xi \rangle \leq q)$ and $q \vdash (\{ \xi < \lambda : \tau(\xi) \in A_\xi \}) \in D[\mathbb{P}_\gamma]$.

Proof. The proof is essentially the same as that of [RS05, Theorem 3.1] with a small modification at the end (in Claim 3.1 there); compare with the proof of Theorem 2.7 here and specifically with 2.7.1.

Definition 1.11. Let $\tilde{E} = (E_\nu : \nu \in <\lambda \lambda)$ be a system of ($<\lambda$)–complete non-principal filters on $\lambda$ and let $E$ be a normal filter on $\lambda$. We define a forcing notion $Q^E_E$ as follows.

A condition $p$ in $Q^E_E$ is a complete $\lambda$–tree $p \subseteq <\lambda \lambda$ such that

- for every $\nu \in p$, either $|\text{succ}_p(\nu)| = 1$ or $\text{succ}_p(\nu) \in E_\nu$, and
- for every $\eta \in \lim\lambda \lambda(p)$ the set $\{ \alpha < \lambda : \text{succ}_p(\eta|\alpha) \in E_{\eta|\alpha} \}$ belongs to $E$.

Note that no relation between $p_\beta^\mu$ and $p_\beta^\mu$ for $\beta < \alpha$ is required to hold.
The order $\preceq = \preceq_{Q_E}$ is the reverse inclusion: $p \preceq q$ if and only if $(p,q) \in Q_E$ and $q \preceq p$.

**Proposition 1.12.** Assume that $E,E$ are as in 1.11. Let $p = (P,S,D)$ be a $\mathcal{D}t$–parameter on $\lambda$ such that $\lambda \setminus S \in E$.

1. $Q_E^P$ is a $<\lambda$–complete forcing notion of size $2^\lambda$.
2. $Q_E^P$ is reasonably $B$–bounding over $p$.
3. If $\lambda$ is strongly inaccessible and $(\forall \delta \in S)(|P_\delta| \leq |\delta|)$, then $\Vdash_{Q_E^P} D^{Q_E^P} \neq D[Q_E^P]$.

**Proof.**

(1) Should be clear.

(2) Let $p \in Q_E^P$. We are going to describe a strategy $st$ for Generic in $\hat{\varnothing}_{\mathcal{D}p}(p, Q_E^P)$. In the course of the play, Generic constructs a sequence $(T_\xi : \xi < \lambda)$ so that if $\langle I_\xi, q_\xi^E, q_\xi^F : t \in I_\xi \rangle : \xi < \lambda$ is the sequence formed by the innings of the two players, then the following conditions are satisfied.

(a) $T_\xi \in Q_E^P$ and if $\xi < \zeta < \lambda$ then $p = T_0 \supseteq T_\xi \supseteq T_\zeta$ and $T_\xi \cap \xi = T_\xi \cap \xi$.

(b) If $\zeta < \lambda$ is limit, then $T_\zeta = \bigcap_{\xi < \zeta} T_\xi$.

(c) If $\xi \in S$ then

- $I_\xi = P_\xi \cap T_\xi$ and $p_\xi^E = (T_\xi)_t$ for $t \in I_\xi$,
- $T_{\xi+1} = \bigcup \{q_\xi^E : t \in I_\xi\} \cup \bigcup \{T_\eta_\nu : \nu \in T_\xi \cap \xi \}$. 

(d) If $\xi \notin S$, then $I_\xi = \emptyset$ and $T_{\xi+1} = T_\xi$.

Conditions (a)–(d) fully describe the strategy $st$. Let us argue that it is a winning strategy and to this end suppose that $\langle I_\xi, q_\xi^E, q_\xi^F : t \in I_\xi \rangle : \xi < \lambda$ is a play of $\hat{\varnothing}_{\mathcal{D}p}(p, Q_E^P)$ in which Generic uses $st$ and constructs the sequence $(T_\xi : \xi < \lambda)$ so that (a)–(d) are satisfied. Put $p^* = \bigcap_{\xi < \lambda} T_\xi \subseteq <^\lambda \lambda$. It follows from (a)+(b) that $p^*$ is a complete $<\lambda$–tree and for each $\nu \in p^*$ either $\|\text{succ}_p(\nu)\| = 1$ or $\text{succ}_p(\nu) \in E_{\nu}$. Suppose now that $\eta \in \text{lim}_\lambda(p^*)$ and for $\xi < \lambda$ let $B_\xi = \{\alpha < \lambda \geq \text{succ}_{T_\xi}(\eta|\alpha) \in E_{\eta|\alpha}\}$. Since $\eta \in \text{lim}_\lambda(T_\xi)$ for each $\xi < \lambda$, we know that $B_\xi \in E$.

$$B = \bigtriangleup_{\xi < \lambda} B_\xi \cap \{\xi < \lambda : \xi \text{ is limit and } \xi \notin S\}.$$ 

It follows from our assumptions that $B \in E$. For each $\alpha \in B$ we know that $\text{succ}_{T_\xi}(\eta|\alpha) \in E_{\eta|\alpha}$ for $\xi < \alpha$ and $T_\alpha = \bigcap_{\xi < \alpha} T_\xi$, so $\text{succ}_{T_\alpha}(\eta|\alpha) \in E_{\eta|\alpha}$. Moreover, $T_\beta \cap \alpha+1 = T_\alpha \cap \alpha+1$ for all $\beta > \alpha$ (remember (a)+(d)) and consequently $\text{succ}_{T_\alpha}(\eta|\alpha) = \text{succ}_{T_\xi}(\eta|\alpha) \in E_{\eta|\alpha}$. Thus we have shown that $p^* \in Q_E^P$.

Let $W$ be a $Q_E^P$–name given by $\Vdash_{Q_E^P} W = \bigcup \{\text{root}(p) : p \in \Gamma_{Q_E^P}\}$. It should be clear that $\|\text{set}^P(W)\| = \lambda$ and thus $\Vdash_{Q_E^P} \text{set}^P(W) \in D[Q_E^P]$.

Plainly, if $\xi \in S$ and $t \in \xi \cap p^*$, then $(p^*)_t \geq q_\xi^E$ and hence

$$(p^*)^\xi_{|\xi} \in \Gamma_{Q_E^P} \Rightarrow W|_\xi \in I_\xi \text{ and } q_\xi^E|_{|\xi} \in \Gamma_{Q_E^P}$$,

so Generic won the play.

(3) We are going to show that $\Vdash_{Q_E^P} \text{set}^P(W) \notin D[Q_E^P]$. To this end suppose that $p \in Q_E^P$ and $A_\xi$ (for $\xi < \lambda$) are $Q_E^P$–names for elements of $D$. Let $st$ be the winning

strategy of Generic in $\otimes_{p}^{\text{rb}}(p, Q_{E}^{F})$ described in part (2) above. Consider a play $\langle I_{\xi}, (q^{t}_{\xi}, q_{\xi}^{\xi} : t \in I_{\xi}) : \xi < \lambda \rangle$ of $\otimes_{p}^{\text{rb}}(p, Q_{E}^{F})$ in which

\begin{itemize}
  \item $(*)_{1}$ Generic follows st and constructs aside a sequence $(T_{\xi} : \xi < \lambda)$,
  \item $(*)_{2}$ Antigeneric plays so that at a stage $\xi \in S$ he picks a set $B_{\xi} \in D$ and conditions $q_{\xi}^{t} \geq p_{\xi}^{t}$ (for $t \in I_{\xi}$) such that
  \[ (\forall t \in I_{\xi})(q_{\xi}^{t} \models B_{\xi} \subseteq \bigcap_{\xi \leq t} A_{\xi}). \]
\end{itemize}

Let $p^{*} = \bigcap_{\xi < \lambda} T_{\xi}$ be the condition determined at the end of part (2) and let $B = \bigtriangleup B_{\xi}$. Choose an increasing continuous sequence $\langle \gamma_{\xi} : \xi < \lambda \rangle \subseteq \lambda$ and a complete $\lambda$–tree $T \subseteq p^{*}$ such that for every $\xi < \lambda$ we have

\begin{itemize}
  \item $(*)_{3}$ if $\nu \in T \cap \gamma_{\xi}$, then $\{\rho \in T \cap \gamma_{\xi+1} : \nu < \rho\} = \{\gamma_{\xi}\}$ and
  \item $(*)_{4}$ if $\rho \in T \cap \gamma_{\xi+1}$, then $\rho \models \gamma_{\xi}$.
\end{itemize}

(The choice can be done by induction on $\xi$; remember that $p$ is a $\mathcal{D}^\ell$–parameter and $\lambda$ is assumed to be inaccessible.) Pick a limit ordinal $\xi \in B \cap S$ such that $\xi = \gamma_{\xi}$. Since $|T \cap \xi| > |\gamma_{\xi}|$, we may choose $\nu \in T \cap \xi \setminus P_{\xi}$. Put $q = (p^{*})_{\nu}$. Then $q \geq p^{*} \geq p$ and $q \models Q_{\xi} \xi \in \bigtriangleup A_{\xi} \setminus \text{set}^p(\mathcal{F})$ (remember $(*)_{2} + (*)_{4}$).

\section{2. Iterations with Lords}

Theorem 1.10 can be used for $\lambda$-support iteration of forcing notions $Q_{E}^{F}$ when on each coordinate we have the same filter $E$. But if we want to use different filters on various coordinates we have serious problems. However, if we move to the other extreme: having very orthogonal filters we may use a different approach to argue that the limit of the iteration is $\lambda$–proper.

\begin{definition}
  \begin{enumerate}
    \item A forcing notion with $\lambda$–complete $(\kappa, \mu)$–purity is a triple $(Q, \leq, \leq_{pr})$ such that $\leq, \leq_{pr}$ are transitive reflexive (binary) relations on $Q$ such that
      \begin{enumerate}
        \item both $(Q, \leq)$ and $(Q, \leq_{pr})$ are 2-strategically $(<\lambda)$–complete,
        \item for every $p \in Q$ and a $(Q, \leq_{pr})$–name $\tau$ for an ordinal below $\kappa$, there is a set $A$ of size less than $\mu$ and a condition $q \in Q$ such that $p \leq_{pr} q$ and $q$ forces (in $(Q, \leq_{pr})$) that “$\tau \in A$”.
      \end{enumerate}
  \end{enumerate}
\end{definition}

\begin{definition}
  \begin{enumerate}
    \item (2) If $(Q, \leq, \leq_{pr})$ is a forcing notion with $\lambda$–complete $(\kappa, \mu)$–purity for every $\kappa$, then we say that it has $\lambda$–complete $(\kappa, \mu)$–purity.
    \item (3) If $(Q, \leq, \leq_{pr})$ is a forcing notion with $\lambda$–complete $(\kappa, \mu)$–purity, then all our forcing terms (like “forces”, “name” etc) refer to $(Q, \leq)$. The relation $\leq_{pr}$ has an auxiliary character only and if we want to refer to it we add “purely” (so $q$ is stronger than $p$) means $p \leq q$, and “$q$ is purely stronger than $p$” means $p \leq_{pr} q$.
  \end{enumerate}
\end{definition}

\begin{definition}
  \begin{enumerate}
    \item Let $Q = (Q, \leq, \leq_{pr})$ be a forcing notion with $\lambda$–complete $(\kappa, \lambda^{+})$–purity, $p = (P, S, D)$ be a $\mathcal{D}^\ell$–parameter on $\lambda$, $\mathcal{U}$ be a normal filter on $\lambda$ and $\tilde{\mu} = \langle \mu_{\alpha} : \alpha < \lambda \rangle$ be a sequence of cardinals below $\lambda$.
      \begin{enumerate}
        \item For a condition $p \in Q$ we define a game $\otimes_{\mathcal{U}, p}^{pr}(p, Q)$ between two players, Generic and Antigeneric, as follows. A play of $\otimes_{\mathcal{U}, p}^{pr}(p, Q)$ lasts $\lambda$ steps
  \end{enumerate}
\end{definition}
and during a play a sequence
\[ \langle \ell_\alpha, \langle p^\alpha_t, q^\alpha_t : t \in \mu_\alpha \rangle : \alpha < \lambda \rangle \]
is constructed. So suppose that the players have arrived to a stage \( \alpha < \lambda \) of the game. Now,

1. \( \exists \bar{p}_\alpha \) first Antigeneric picks \( \ell_\alpha \in \{0, 1\} \).
2. After this, Generic chooses a system \( \langle p^\alpha_t : t \in \mu_\alpha \rangle \) of pairwise incompatible conditions from \( Q_\alpha \), and
3. Antigeneric answers with a system of conditions \( q^\alpha_t \in Q_\alpha \) (for \( t \in \mu_\alpha \)) such that for each \( t \in \mu_\alpha \):
   - \( p^\alpha_t \leq q^\alpha_t \), and
   - if \( \ell_\alpha = 1 \), then \( p^\alpha_t \leq_{pr} q^\alpha_t \).

At the end, Generic wins the play
\[ \langle \ell_\alpha, \langle p^\alpha_t, q^\alpha_t : t \in \mu_\alpha \rangle : \alpha < \lambda \rangle \]
if and only if either \( \{ \alpha < \lambda : \ell_\alpha = 1 \} \not\in U \), or
\( \forall \bar{p}_\alpha \) there is a condition \( p^* \in Q \) stronger than \( p \) and such that
\[ p^* \Vdash \{ \alpha < \lambda : (\exists t \in \mu_\alpha)(q^\alpha_t \in \Gamma_\alpha) \} \in D[Q] \].

(2) We say that the forcing notion \( \bar{Q} \) (with \( \lambda \)-complete \((\ast, \lambda^+)\)-purity) is purely \( B^*\)-bounding over \( U, p, \bar{\mu} \) if for any \( p \in Q \), Generic has a winning strategy in the game \( \bar{\mathcal{O}}_{U, p, \bar{\mu}}(p, Q) \).

Remark 2.3. Note that in the definition of the game \( \bar{\mathcal{O}}_{U, p, \bar{\mu}}(p, Q) \) the size of the index set used at stage \( \alpha \) is declared to be \( \mu_\alpha \) (while in the related game \( \mathcal{O}^{pr}_U(p, Q) \) we required just \( \mu_\alpha < \lambda \)). The reason for this is that otherwise in the proof of the iteration theorem for the current case we could have problems with deciding the size of the set \( I_\alpha \); compare clause \((*)_4\) of the proof of Theorem 2.7.

Observation 2.4. Assume \( \bar{E}, E \) are as in 1.11. For \( p, q \in Q^E_\bar{E} \) let \( p \leq_{pr} q \) mean that \( p \leq q \) and \( \text{root}(p) = \text{root}(q) \). Then

1. \( (Q^E_\bar{E}, \leq, \leq_{pr}) \) is a forcing notion with \( \lambda \)-complete \((\ast, \lambda^+)\)-purity,
2. if, additionally, each \( E_\nu \) (for \( \nu < \lambda \)) is an ultrafilter on \( \lambda \), then \( (Q^E_\bar{E}, \leq, \leq_{pr}) \) has \((\kappa, 2)\)-purity for every \( \kappa < \lambda \).

Proposition 2.5. Assume that \( \bar{E}, E \) are as in 1.11, \( p = (P, S, D) \) is a \( \mathcal{D}\)-parameter on \( \lambda \) and \( \bar{\mu} = \langle \mu_\alpha : \alpha < \lambda \rangle \) is a sequence of non-zero cardinals below \( \lambda \) such that \( (\forall \alpha \in S)(|P_\alpha| \leq \mu_\alpha) \). Then \( (Q^E_\bar{E}, \leq, \leq_{pr}) \) is purely \( B^*\)-bounding over \( E, p, \bar{\mu} \).

Proof. Let \( p \in Q^E_\bar{E} \) and let \( st \) be the strategy described in the proof of 1.12(2) with a small modification that we start the construction with \( \xi_0 = \text{lh}(\text{root}(p)) + 1 \) (so \( T_{\xi_0} = p \) and the first \( \xi_0 \) steps of the play are not relevant). Then we also replace clauses (c)+(d) there by

(c)\( \xi \geq \xi_0 \) then
- \( I_\xi \subseteq \xi \cap T_\xi \) is of size \( \mu_\xi \), and \( p^\xi_t = (T_\xi)_t \) for \( t \in I_\xi \), and
- if \( \xi \in S \) then \( P_\xi \cap T_\xi \subseteq I_\xi \), and
- \( T_{\xi+1} = \{ q^\xi_t : t \in I_\xi \} \cup \{ (T_\xi)_\nu : \nu \in \xi \cap T_\xi \setminus I_\xi \} \).
(So, in particular, Antigeneric’s choice of \( \ell, \xi \) has no influence on the answers by Generic.) We are going to show that \( \text{st} \) is a winning strategy for Generic in \( O_E^{pr}(p, Q_E) \). To this end suppose that \( \langle \ell, (p^t(X), t : t \in \ell) : \xi < \lambda \rangle \) is a play of \( O_E^{pr}(p, Q_E) \) in which Generic follows \( \text{st} \) (we identify \( |\ell| = |\mu| \) and \( |T| : \xi < \lambda \) is the sequence of side objects constructed in the course of the play. Assume \( A = \{ \xi : \lambda : \xi < \xi_0 \} \in E \) (otherwise Generic wins by default). Like in 1.12(2), put \( p^* \in Q_E^{\xi} \) we note that if \( \eta \in \text{lim}_\lambda(p^*) \) and

\[
\delta \in \bigtriangleup_{\xi < \lambda} \{ \alpha < \lambda : \text{succ}_T(\eta[\alpha]) \in E_{\eta[\alpha]} \} \cap A \cap \{ \xi < \lambda : \xi > \xi_0 \text{ is limit} \},
\]

then

\[
\text{succ}_{p^*}(\eta[\delta]) = \begin{cases} 
\text{succ}_T(\eta[\delta]) & \text{if } \eta[\delta] \notin I_\delta \\
\text{succ}_{p^*_{\eta[\delta]}}(\eta[\delta]) & \text{if } \eta[\delta] \in I_\delta \in E_{\eta[\delta]}.
\end{cases}
\]

Exactly as in 1.12(2) we justify that \( p^* \) witnesses \( (\otimes)^{p_0} \).

\[\square\]

**Lemma 2.6.** Assume that

\( (1) \) \( \lambda \) is strongly inaccessible,

\( (2) \) \( Q = (P_0, Q_0 : \alpha < \gamma) \) is a \( \gamma \)-support iteration, \( w \subseteq \gamma, \ |w| < \lambda, \alpha_0 \in w, \)

\( (3) \) for every \( \alpha < \gamma, \)\n
\( \| F_{\alpha} \text{ “ } Q_\alpha = (Q_\alpha, \leq, \leq_{pr}) \text{ is a forcing notion with } \lambda \text{-complete } (\ast, \lambda^+) \text{-parity “} , \)

\( (4) \) \( P_{\alpha_0} \) is \( \lambda \)-proper,

\( (5) \) \( T = (T, \text{rk}) \) is a standard \((w, 1)\gamma\)-tree, \( |T| < \lambda, \)

\( (6) \) \( p = (p_\ell : \ell \in T) \) is a standard tree of conditions in \( P_\gamma, \)

\( (7) \) \( \tau \) is a \( P_\gamma \)-name for an ordinal.

Then there are a set \( A \) of size \( \lambda \) and a standard tree of conditions \( q = (q_\ell : \ell \in T) \subseteq P_\gamma \) such that

\( (a) \) \( (\forall \ell \in T)(\text{rk}(\ell) = \gamma \Rightarrow q_\ell \models \tau \in A), \) and

\( (b) \) \( p \leq q \) and if \( t \in T, \text{rk}(t) > \alpha_0 \) then \( q_\ell(q_{t|\alpha_0} \models_{P_{\alpha_0}} \tau(t) \leq_{pr} q_{t}(\alpha_0) \).

**Proof.** Let us start with the following observation.

**Claim 2.6.1.** If \( p \in P_\gamma \) then there are a set \( A_0 \) of size \( \lambda \) and a condition \( q \geq p \) such that \( q \models_{P_\gamma} \tau \in A_0 \) and \( q_{|\alpha_0} \models_{P_\alpha_0} p(\alpha_0) \leq_{pr} q(\alpha_0) \).

**Proof of the Claim.** Let us look at \( P_\gamma \) as the result of 3 stage composition \( P_{\alpha_0} \ast Q_{\alpha_0} \ast \mathbb{P}_{\alpha_0+1} \), \( \mathbb{P}_{\alpha_0+1} \gamma \) where \( \mathbb{P}_{\alpha_0+1} \gamma \) is a \( \mathbb{P}_{\alpha_0} \)-name for the following forcing notion. The set of conditions in \( \mathbb{P}_{\alpha_0+1} \gamma \) is \( \{ \tau(\alpha, \gamma) : r \in \mathbb{P}_\gamma \} \) (so it belongs to \( V \)); the order of \( \mathbb{P}_{\alpha_0+1} \gamma \) is such that if \( G_{\alpha_0+1} \subseteq P_{\alpha_0+1} \) is generic over \( V \), then

\[
V[G_{\alpha_0+1}] \models \text{ “ } r \leq_{P_{\alpha_0+1}1|\gamma} G_{\alpha_0} \text{ if and only if there is } q \in G_{\alpha_0+1} \text{ such that } q ^ \prec r \leq_{\mathbb{P}_{\alpha_0}} q ^ \prec s.“ \]

Now, pick a \( \mathbb{P}_{\alpha_0} \)-name \( (\tau, \alpha) \) such that

\[
p(\alpha_0 + 1) \models_{\mathbb{P}_{\alpha_0+1}} \text{ “ } p(\alpha_0, \gamma) \leq \tau \text{ and } \tau \models_\alpha \tau = \tau \text{” }
\]

and then choose a \( \mathbb{P}_{\alpha_0} \)-name \( A^\ast \) for a subset of \( \mathbb{P}_{\alpha_0+1} \gamma \times ON \) and a \( \mathbb{P}_{\alpha_0} \)-name \( q(\alpha_0) \) for a condition in \( Q_{\alpha_0} \) such that

\[
p(\alpha_0) \models_{\mathbb{P}_{\alpha_0}} \text{ “ } p(\alpha_0) \leq_{pr} q(\alpha_0) \text{ and } |A^\ast| = \lambda \text{ and } q(\alpha_0) \models_{\mathbb{P}_{\alpha_0}} (\exists (s, \beta) \in A^\ast)(\tau = s \& \alpha = \beta \text{” }) .
\]
Since $\mathbb{P}_{\alpha_0}$ is $\lambda$-proper, we may choose a set $A^+ \subseteq \mathbb{P}(\alpha_0 + 1), \gamma \times \text{ON}$ of size $\lambda$ and a condition $q|\alpha_0 \geq p|\alpha_0$ such that $q|\alpha_0 \Vdash A^* \subseteq A^+$. Then

$$q|\alpha_0 + 1 \Vdash \mathbb{P}_{\alpha_0 + 1} \left( \exists (s, \beta) \in A^+ \right) \left( \tau = s \land \alpha = \beta \right).$$

Put $A = \{ \beta : (\exists s)((s, \beta) \in A^+) \}$. Now we may easily define $q|\alpha_0, \gamma$ so that $\text{dom}(q|\alpha_0, \gamma) = \bigcup \{ \text{dom}(s) : (\exists s)((s, \beta) \in A^+) \}$ and $q|\alpha_0 + 1 \Vdash \mathbb{P}_{\alpha_0 + 1} \left( \exists \xi \leq \mathbb{P}(\alpha_0 + 1), \gamma \right) q|\alpha_0, \gamma \text{ and } \alpha \in A^+$. \hfill \Box

Fix an enumeration $(t_\xi : \zeta \leq \zeta^*)$ of $\{ t \in T : \text{rk}(t) = \zeta \}$ (so $\zeta^* < \lambda$). For each $\alpha \in \gamma \setminus \{ \alpha_0 \}$ fix a $\mathbb{P}_\alpha$-name $\mathbf{st}^\alpha_{\text{pr}}$ for a winning strategy of Complete in the game $\mathbb{G}_\alpha^0 \left( (\mathbb{Q}_\alpha, \leq, \mathbb{V}_{\alpha_0}) \right)$ such that as long as Incomplete plays $\mathbb{V}_{\alpha_0}$, Complete answers with $\mathbb{V}_{\alpha_0}$ as well. Let $\mathbf{st}^\zeta_{\text{pr}}$ be the $<_\zeta^*-$first $\mathbb{P}_{\alpha_0}$-name for a winning strategy of Complete in $\mathbb{G}_\alpha^0 \left( (\mathbb{Q}_{\alpha_0}, \leq_{\text{pr}}), p_\epsilon(\alpha_0) \right)$ (for $\zeta \leq \zeta^*$). Note that if $\zeta, \zeta' \leq \zeta^*$ and $t_\xi \big| (\alpha_0 + 1) = t_\xi', (\alpha_0 + 1)$, then $\mathbf{st}^\zeta_{\text{pr}} = \mathbf{st}^{\zeta'}_{\text{pr}}$.

By induction on $\zeta \leq \zeta^*$ we choose a sequence $\langle \tilde{p}^\xi, \tilde{q}^\xi, A^\xi : \zeta \leq \zeta^* \rangle$ so that the following demands are satisfied.

(i) $\tilde{p}^\xi = \langle \tilde{p}^\xi_t : t \in T \rangle, \tilde{q}^\xi = \langle \tilde{q}^\xi_t : t \in T \rangle$ are standard trees of conditions, $A^\xi$ is a set of ordinals of size $\lambda$.

(ii) If $\epsilon < \zeta \leq \zeta^*$, then $\tilde{p}^\epsilon \leq \tilde{p}^\xi \leq \tilde{q}^\xi \leq \tilde{p}^\zeta$ and $A^\xi \subseteq A^\zeta$.

(iii) $p^\xi_\epsilon \Vdash_{\mathbb{P}_\epsilon} \tau \in A^\xi$.

(iv) If $\alpha \in \gamma \setminus \{ \alpha_0 \}, \zeta, \xi \leq \zeta^*$, then $q^\zeta_{\epsilon \alpha} \Vdash_{\mathbb{P}_\alpha} \langle \tilde{p}^\zeta_{\epsilon \alpha}(\alpha), q^\zeta_{\epsilon \alpha}(\alpha) : \epsilon \leq \xi \rangle$ is a result of a play of $\mathbb{G}_\alpha^0 \left( (\mathbb{Q}_\alpha, \leq, \mathbb{V}_{\alpha_0}) \right)$ in which Complete uses $\mathbf{st}^\alpha_{\text{pr}}$.

(v) If $\zeta, \xi \leq \zeta^*$, then $q^\zeta_{\epsilon \alpha} \Vdash_{\mathbb{P}_\alpha} \langle \tilde{p}^\zeta_{\epsilon \alpha}(\alpha), q^\zeta_{\epsilon \alpha}(\alpha) : \epsilon \leq \xi \rangle$ is a result of a play of $\mathbb{G}_\alpha^0 \left( (\mathbb{Q}_{\alpha_0}, \leq_{\text{pr}}), p_\epsilon(\alpha_0) \right)$ in which Complete uses $\mathbf{st}^\alpha_{\text{pr}}$.

Suppose that we have determined $\tilde{p}^\xi, \tilde{q}^\xi, A^\xi$ for $\epsilon < \zeta \leq \zeta^*$. First we choose $\tilde{p}^\zeta = \langle \tilde{p}^\zeta_t : t \in T \rangle \subseteq \mathbb{P}_\gamma$. If $\xi = 0$ then we set $\tilde{p}^\zeta = \tilde{p}^\epsilon$. Otherwise we choose $\tilde{p}^\zeta$ so that for $t \in T$ we have:

(vi) $\text{dom}(p^\zeta_t) = \bigcup_{\epsilon < \zeta} \text{dom}(q^\xi_t)$, and

(vii) if $\alpha \in \text{dom}(p^\zeta_t) \setminus \{ \alpha_0 \}$, then $p^\zeta(\alpha)$ is the $<_{\zeta^*}$-first $\mathbb{P}_\alpha$-name for a condition in $\mathbb{Q}_\alpha$ such that $p^\xi_t|\alpha \Vdash_{\mathbb{P}_\epsilon}, \forall \epsilon < \zeta (q^\xi_t(\alpha) \leq p^\xi_t(\alpha))$, and

(viii) $p^\zeta(\alpha) = \langle p^\zeta_{\epsilon \alpha}(\alpha), q^\zeta_{\epsilon \alpha}(\alpha) : \epsilon < \zeta, \alpha \rangle$ is a result of a play of $\mathbb{G}_\alpha^0 \left( (\mathbb{Q}_{\alpha_0}, \leq_{\text{pr}}), p_\epsilon(\alpha_0) \right)$ in which Complete uses $\mathbf{st}^\alpha_{\text{pr}}$.

The choice is possible by (iv)+(v), and since we pick "the $<_{\zeta^*}$-first names" we easily see that $\tilde{p}^\zeta$ is a standard tree of conditions. Now we use 2.6.1 to find a set $A^\zeta$ of size $\lambda$ and a condition $p^\zeta_{\epsilon \alpha} \in \mathbb{P}_\gamma$ such that

$$\bigcup_{\epsilon < \zeta} A^\xi \subseteq A^\zeta, \quad p^\zeta_{\epsilon \alpha} \leq p^\zeta_{\epsilon \alpha}, \quad p^\zeta_{\epsilon \alpha} \Vdash_{\mathbb{P}_\alpha} p^\zeta_{\epsilon \alpha} \leq_{\text{pr}} p^\zeta_{\epsilon \alpha}(\alpha_0) \quad \text{and} \quad p^\zeta_{\epsilon \alpha} \Vdash_{\mathbb{P}_\alpha} \tau \in A^\xi.$$

Next, for each $t \in T$ we let $p^\zeta_t \in \mathbb{P}\text{rk}(t)$ be such that

(ix) if $s = t \cap t_\xi$, then $p^\zeta_t|\text{rk}(s) = p^\zeta_s|\text{rk}(s)$ and $p^\zeta_t|\text{rk}(t) = p^\zeta_s|\text{rk}(s, \text{rk}(t))$. 


Clearly, \( p^\natural = (p^\natural_t : t \in T) \) is a standard tree of conditions satisfying the relevant parts of the demands in (ii)–(v). Now we choose a tree of conditions \( q^\natural = (q^\natural_t : t \in T) \) so that the requirements of (iv)+(v) hold (for this we proceed like in (vi)–(viii) above).

After the construction is carried out we note that \( q^\natural \) and \( A^\natural \) are as required in the assertion of the lemma. \( \Box \)

**Theorem 2.7.** Assume that

1. \( \lambda \) is strongly inaccessible, \( \bar{\mu} = (\mu_\alpha : \alpha < \lambda) \) is a sequence of cardinals below \( \lambda \), \( p = (P, S, D) \) is a \( \Delta \ell \)-parameter on \( \lambda \), and
2. \( \bar{Q} = \{Q_\alpha : \alpha < \gamma\} \) is a \( \lambda \)-support iteration,
3. \( \mathcal{U}_\alpha \) is a \( \mathcal{P}_\alpha \)-name for a normal filter on \( \lambda \) (for \( \alpha < \gamma \)),
4. \( A_{\alpha, \beta} \subseteq \lambda \) is such that \( \models \mathcal{P}_\alpha A_{\alpha, \beta} \in \mathcal{U}_\alpha \) and \( \models \mathcal{P}_\beta \lambda \setminus A_{\alpha, \beta} \in \mathcal{U}_\beta \) (for \( \alpha < \beta < \gamma \)), and
5. for every \( \alpha < \gamma \), \( \models \mathcal{P}_\alpha ^{\beta_\gamma} \) is purely \( B^\ast \)-bounding over \( \mathcal{U}_\alpha, p[\mathcal{P}_\alpha], \bar{\mu} \).

Then \( \mathcal{P}_\gamma = \lim(\bar{Q}) \) is \( \lambda \)-proper.

**Proof.** The arguments follow closely the lines of the arguments for [RS05, Thm. 3.1, 3.2] and [RS, Thm. 2.12]. The proof is by induction on \( \gamma \), so assume that we know also that each \( \mathcal{P}_\alpha \) is \( \lambda \)-proper for \( \alpha < \lambda \).

Let \( N = (\mathcal{H}(\chi), \in, < \gamma) \) be such that \( \langle \mathcal{H}(\chi), \in, < \gamma \rangle \models \langle \lambda, \mathcal{P}_\alpha, \mathcal{Q}_\alpha, \alpha < \gamma, \mathcal{P}_\gamma \rangle \) and \( \models \langle \alpha < \beta < \gamma \rangle \) be such that \( \langle \alpha < \beta < \gamma \rangle \models \langle \mathcal{H}(\chi), \in, < \gamma \rangle \) is a \( \lambda \)-support iteration, \( \mathcal{U}_\alpha \) is a \( \mathcal{P}_\alpha \)-name for a normal filter on \( \lambda \) (for \( \alpha < \gamma \)), and \( A_{\alpha, \beta} \subseteq \lambda \) is such that \( \models \mathcal{P}_\alpha A_{\alpha, \beta} \in \mathcal{U}_\alpha \) and \( \models \mathcal{P}_\beta \lambda \setminus A_{\alpha, \beta} \in \mathcal{U}_\beta \) (for \( \alpha < \beta < \gamma \)),

so that the following demands are satisfied.

\((*)_0\) All objects listed in \((\circ)_0^\alpha + (\circ)^{\beta_\gamma}_0\) belong to \( N \). After stage \( \delta < \lambda \) of the construction, the objects in \((\circ)_0^\alpha \) are known as well as those in \((\circ)^{\beta_\gamma}_0\) for \( \xi \in w_\delta \).

\((*)_1\) \( r_\delta, r_\gamma \in \mathcal{P}_\gamma, r_0(0) = r_0(0) = p(0), w_\delta \subseteq \gamma, |w_\delta| = |\delta| + 1, \bigcup_{\alpha < \lambda} \text{dom}(r_\alpha) = w_\alpha = N \cap \gamma, w_0 = \{0\}, w_\delta \subseteq w_{\delta + 1} \) and if \( \delta \) is limit then \( w_\delta = \bigcup_{\alpha < \delta} w_\alpha \).

\((*)_2\) For each \( \alpha < \beta < \lambda \) we have \( (\forall \xi \in w_{\alpha + 1})(r_\alpha(\xi) = r_\gamma(\xi)) \) and \( p \leq r_\alpha \leq r_\gamma \leq r_\delta \).

\((*)_3\) If \( \xi \in (\gamma \setminus w_\delta) \cap N \), then \( r_\delta(\xi) \vdash \langle \mathcal{P}_\alpha, \mathcal{Q}_\xi, \bar{\mu}_\xi \rangle \) is a standard \( (w_\delta, 1) \)-tree, \( T_\delta = \bigcup_{\alpha \leq \gamma} \bigcap_{\xi \in w_\delta \cap \alpha} \mu_\xi \).
and for each $\xi \in \text{dom}(p_{\delta,t}^\delta) \setminus w_\delta$:

\[ p_{\delta,t}^\delta(\xi) \models \varphi_\gamma \quad \text{if the set } \{r_\alpha(\xi) : \alpha < \delta\} \cup \{p(\xi)\} \text{ has an upper bound in } Q_\xi, \]

then $p_{\delta,t}^\delta(\xi)$ is such an upper bound ".

(*7) For $\xi \in N \cap \gamma$, $\ell_\delta, \xi \in \{0,1\}$ and $\bar{p},\xi, \bar{q},\xi$ are $P_\xi$–names for sequences of conditions in $Q_\xi$ of length $\mu_\delta$.

(*8) If either $\xi = 0 = \beta$ or $\xi \in w_{\beta+1} \setminus w_\beta$, $\beta < \lambda$, then

\[ \models \varphi_\gamma \quad \langle \ell_\alpha, \xi, \bar{p}_\alpha, \xi, \bar{q}_\alpha, \xi : \alpha < \lambda \rangle \text{ is a play of } D^{br}_{\ell_\xi, P, \xi}(r_\beta(\xi), Q_\xi) \]

in which Generic uses $st_\xi$ ".

(*9) $\alpha_\delta \in w_\delta$ ($\alpha_\delta$ will be called the lord of stage $\delta$) and

\[ (\forall \beta \in w_\delta \setminus \{\alpha_\delta\})(\delta \notin \bigcap \{\lambda \setminus A_{\xi, \beta} : \xi \in w_\delta \cap \beta\}) \cap \bigcap \{A_{\beta, \xi} : \xi \in w_\delta \setminus (\beta + 1)\}). \]

(*10) $\ell_\delta, \xi = 0$ for $\xi \in N \cap \gamma \setminus \{\alpha_\delta\}$ and $\ell_\delta, \alpha_\delta = 1$.

(*11) If $t \in T_\delta$, $rk_\delta(t) = \xi < \gamma$, then for each $\epsilon < \mu_\delta$

\[ q_{\epsilon,t}^\delta \models \varphi_\gamma \quad \langle \bar{p}_\delta(\epsilon) = p_{\epsilon,t-\epsilon(\epsilon)}^\delta(\xi) \rangle \quad \text{and } \bar{q}_\delta(\epsilon) = q_{\epsilon,t-\epsilon(\epsilon)}^\delta(\xi) \quad \text{".} \]

(*12) If $t_0, t_1 \in T_\delta$, $rk_\delta(t_0) = rk_\delta(t_1)$ and $\xi \in w_\delta \cap rk_\delta(t_0)$, $t_0|\xi = t_1|\xi$ but

\[ (t_0|\xi) \neq (t_1|\xi), \quad \text{then } \]

\[ p_{\delta,t_0|\xi}^\delta \models \varphi_\gamma \quad \text{the conditions } p_{\delta,t_0|\xi}^\delta(\xi), p_{\delta,t_1|\xi}^\delta(\xi) \text{ are incompatible "} \]

(*13) $Z_\delta$ is a set of ordinals, $|Z_\delta| = \lambda$ and for each $t \in T_\delta$ with $rk_\delta(t) = \gamma$ we have

\[ q_{\epsilon,t}^\delta \models \varphi_\gamma \quad (\forall \alpha \leq \delta)(\exists \tau_\alpha \in Z_\delta). \]

(*14) $\text{dom}(r_\delta^\alpha) = \text{dom}(r_\delta) = \bigcup_{\tau \in T_\delta} \text{dom}(q_{\epsilon,t}^\delta \cup \text{dom}(\rho))$ and if $t \in T_\delta$, $\xi \in \text{dom}(r_\delta) \cap$ \begin{array}{c}
\text{rk}_\delta(t) \setminus w_\delta, \quad \text{and } q_{\epsilon,t|\xi}^\delta \leq q \in P_\xi, \, r_\delta|\xi \leq q, \, \text{then}
\end{array}

\[ q \models \varphi_\gamma \quad \text{if the set } \{r_\alpha(\xi) : \alpha < \delta\} \cup \{q_{\epsilon,t}^\delta(\xi), p(\xi)\} \text{ has an upper bound in } Q_\xi, \]

then $r_\delta(\xi)$ is such an upper bound ".

First we fix an increasing continuous sequence $\langle w_\alpha : \alpha < \lambda \rangle$ of subsets of $N \cap \gamma$ such that the relevant demands in (*1) are satisfied. Now, suppose that we have arrived to a stage $\delta < \lambda$ of the construction and all objects listed in $(\otimes)_\alpha$ and relevant cases of $(\otimes)_\alpha$ (see (*9)) have been determined for $\alpha < \delta$.

To ensure (*9), all choices below are made in $N$ (e.g., each time we choose an object with some properties, we pick the $<_\lambda$–first such object).

If $\delta$ is a successor ordinal and $\xi \in w_\delta \setminus w_{\delta-1}$, then we let $st_\xi \in N$ be a $P_\xi$–name for a winning strategy of Generic in $D^{br}_{\ell_\xi, P, \xi}(r_{\delta-1}(\xi), Q_\xi)$. We also put $\ell_\delta, \xi = 0$ for all $\alpha < \delta$ and we pick $\bar{p}_\delta, \xi, \bar{q}_\delta, \xi$ (for $\alpha < \delta$) so that the suitable parts of (*7) + (*8) at $\xi$ are satisfied.

Clause (*4) fully describes $T_\delta$. Now we choose the lord of stage $\delta$. If for some $\beta \in w_\delta$ we have

\[ \delta \in \bigcap \{\lambda \setminus A_{\xi, \beta} : \xi \in w_\delta \cap \beta\} \cap \bigcap \{A_{\beta, \xi} : \xi \in w_\delta \setminus (\beta + 1)\}, \]

then $\alpha_\delta = \beta$.
then $\alpha_\delta$ is equal to this $\beta$ (note that there is at most one $\beta \in w_\delta$ with the required property). Otherwise we let $\alpha_\delta = 0$. Then we put $\ell_{\delta,\alpha_\delta} = 1$ and $\ell_{\delta,\xi} = 0$ for all $\xi \in w_\delta \setminus \{\alpha_\delta\}$.

Next, for each $\xi \in w_\delta$ we choose a $\mathbb{P}_\xi$-name $\bar{p}_{\delta,\xi}$ such that

\[
\models \mathbb{P}_\xi \quad \bar{p}_{\delta,\xi} = \langle \bar{p}_{\delta,\xi}(\varepsilon) : \varepsilon < \mu_\delta \rangle \text{ is given to Generic by } \mathbf{st}_\xi \text{ as an answer to } \langle \ell_{\alpha,\delta}, \bar{p}_{\alpha,\xi}, \bar{q}_{\alpha,\xi} : \alpha < \delta \rangle^{-\langle \ell_{\delta,\xi} \rangle}.
\]

(Note that $\models \mathbb{P}_\xi$ “conditions $\bar{p}_{\delta,\xi}(\varepsilon_0), \bar{p}_{\delta,\xi}(\varepsilon_1)$ are incompatible” whenever $\varepsilon_0 < \varepsilon_1 < \mu_\delta$ and $\xi \in w_\delta$.) After this we may choose a tree of conditions $p^\delta_{\varepsilon, t} = \langle p^\delta_{\varepsilon, t} : t \in T_\delta \rangle$ such that for each $t \in T_\delta$:

- $\text{dom}(p^\delta_{\varepsilon, t}) = (\text{dom}(p) \cup \bigcup_{\alpha < \delta} \text{dom}(r_{\alpha}) \cup w_\delta) \cap \text{rk}_\delta(t)$ and
- for $\xi \in \text{dom}(p^\delta_{\varepsilon, t}) \setminus w_\delta$, $p^\delta_{\varepsilon, t}(\xi)$ is the $<_\chi$-first $\mathbb{P}_\xi$-name for a condition in $Q_\xi$ such that

\[
p^\delta_{\varepsilon, t}(\xi) \models \mathbb{P}_\xi \quad \text{if the set } \{r_\alpha(\xi) : \alpha < \delta\} \cup \{p(\xi)\} \text{ has an upper bound in } Q_\xi,
\]

then $p^\delta_{\varepsilon, t}(\xi)$ is such an upper bound ”,

- $p^\delta_{\varepsilon, t}(\xi) = \bar{p}_{\alpha,\xi}(\langle t \rangle) \text{ for } \xi \in \text{dom}(p^\delta_{\varepsilon, t}) \cap w_\delta$.

Using Lemma 2.6 we may pick a tree of conditions $q^\delta_{\varepsilon, t} = \langle q^\delta_{\varepsilon, t} : t \in T_\delta \rangle$ and a set $Z_\delta$ of ordinals such that

- $p^\delta_{\varepsilon, t} \leq q^\delta_{\varepsilon, t}, |Z_\delta| = \lambda$,
- if $t \in T_\delta$, $\text{rk}_\delta(t) = \gamma$ then $q^\delta_{\varepsilon, t} \models \mathbb{P}_\xi$ (\forall \alpha \leq \delta)(\tau_\alpha \in Z_\delta),
- if $t \in T_\delta$, $\text{rk}_\delta(t) > \alpha_\delta$ then $q^\delta_{\varepsilon, t}[\alpha_\delta] \models \mathbb{P}_\delta$. $p^\delta_{\varepsilon, t}(\alpha_\delta) \leq \mathbf{pr} q^\delta_{\varepsilon, t}(\alpha_\delta)$.

Note that if $\xi \in w_\delta$ and $\varepsilon_0 < \varepsilon_1 < \mu_\delta$, $t \in T_\delta$, $\text{rk}_\delta(t) = \xi$, then

\[
q^\delta_{\varepsilon, t} \models \mathbb{P}_\xi \quad \text{the conditions } q^\delta_{\varepsilon, t \langle \varepsilon_0 \rangle}(\xi), q^\delta_{\varepsilon, t \langle \varepsilon_1 \rangle}(\xi) \text{ are incompatible }.
\]

Hence we have no problems with finding $\mathbb{P}_\xi$-names $\bar{q}_{\delta, \xi}$ (for $\xi \in w_\delta$) such that

- $\models \mathbb{P}_\xi$ “$\bar{q}_{\delta, \xi} = \langle \bar{q}_{\delta, \xi}(\varepsilon) : \varepsilon < \mu_\delta \rangle$ is a sequence of conditions in $Q_\xi$ ”,
- $\models \mathbb{P}_\xi$ “(\forall \varepsilon < \mu_\delta)(\bar{q}_{\delta, \xi}(\varepsilon) \leq \bar{q}_{\delta, \xi}(\varepsilon))” and $\models \mathbb{P}_\delta$ “(\forall \varepsilon < \mu_\delta)(\bar{q}_{\delta, \alpha_\delta}(\varepsilon) \leq \mathbf{pr} \bar{q}_{\delta, \alpha_\delta}(\varepsilon))”
- if $t \in T_\delta$, $\text{rk}_\delta(t) > \xi$, then $q^\delta_{\varepsilon, t}[\xi] \models \mathbb{P}_\xi q^\delta_{\varepsilon, t}(\xi) = \bar{q}_{\delta, \xi}(\langle t \rangle)$.

Now we define $r^-_\delta, r_\delta \in \mathbb{P}_\gamma$, so that

\[
\text{dom}(r^-_\delta) = \text{dom}(r_\delta) = \bigcup_{t \in T_\delta} \text{dom}(q^\delta_{\varepsilon, t}) \cup \text{dom}(p)
\]

and

- if $\xi \in w_\alpha + 1$, $\alpha < \delta$, then $r^-_\delta(\xi) = r_\delta(\xi) = r_\alpha(\xi)$,
- if $\xi \in \text{dom}(r^-_\delta) \setminus w_\delta$, then $r^-_\delta(\xi)$ is the $<^*_\chi$-first $\mathbb{P}_\xi$-name for an element of $Q_\xi$ such that

\[
r^-_\delta[\xi] \models \mathbb{P}_\xi \quad r^-_\delta(\xi) \text{ is an upper bound of } \{r_\alpha(\xi) : \alpha < \delta\} \cup \{p(\xi)\} \text{ and if } t \in T_\delta, \text{ } \text{rk}_\delta(t) > \xi, \text{ and } q^\delta_{\varepsilon, t}[\xi] \in \Gamma_{P_\xi} \text{ and the set }
\]

\[
\{r_\alpha(\xi) : \alpha < \delta\} \cup \{q^\delta_{\varepsilon, t}(\xi), p(\xi)\} \text{ has an upper bound in } Q_\xi,
\]

then $r^-_\delta(\xi)$ is such an upper bound ”,

and $r_\delta(\xi)$ is the $<^*_\chi$-first $\mathbb{P}_\xi$-name for an element of $Q_\xi$ such that

\[
r_\delta[\xi] \models \mathbb{P}_\xi \quad r_\delta(\xi) \text{ is given to Complete by } \mathbf{st}_\xi \text{ as the answer to }
\]

\[
\langle r^-_\alpha(\xi), r_\alpha(\xi) : \alpha < \delta \rangle.
\]
For each limit ordinal

Claim 2.7.1. Proof of the Claim.

Note that by a straightforward induction on \( \xi \in \text{dom}(r_\delta) \) one easily applies \((*)_3\) from previous stages to show that \( r_\delta^\gamma, r_\delta \) are well defined and \( r_\delta \geq r_\delta^\gamma \geq r_\alpha, p \) for \( \alpha < \delta \). If \( \delta = 0 \) we also stipulate \( r_0(0) = r_0(0) = p(0) \).

This completes the description of stage \( \delta \) of our construction. One easily verifies that the demands \((*)_0\)–\((*)_{14}\) are satisfied.

After completing all \( \lambda \) stages of the construction, for each \( \xi \in N \cap \gamma \) we look at the sequence \langle \ell, p, \tilde{q}, q_\delta : \delta < \lambda \rangle. \) For \( \delta < \lambda \) such that \( \xi \in w_\delta \) let

\[
B_\delta^\xi = \bigcap \{ \lambda \setminus A_\xi, \gamma : \xi \in w_\delta \cap \gamma \} \cap \{ \lambda \setminus A_\xi, \gamma : \xi \in w_\delta \} 
\]

and for \( \lambda < \) such that \( \xi \not\in w_\delta \) put \( B_\delta^\xi = \lambda. \) It follows from our assumptions that \( \models_{\alpha < \lambda} (\forall \delta < \lambda)(B_\delta^\xi \subseteq U_\xi) \) and thus also \( \models_{\alpha < \lambda} B_\delta^\xi \subseteq U_\xi \). Note that if \( \delta \) is a limit ordinal, \( \xi \in w_\delta \) and \( \delta \in \bigtriangleup B_\delta^\xi \), then also \( \delta \in B_\delta^\xi \) and hence \( \xi = \alpha_\delta \) (remember \((*)_0\)) and \( \ell_\delta, \gamma = 1 \) (by \((*)_{15}\)). Consequently for each \( \xi \in N \cap \gamma \)

\[ \models_{\alpha < \lambda} \{ \delta < \lambda : \ell_\delta, \gamma = 1 \} \in U_\xi \] .

Therefore, for every \( \xi \in N \cap \gamma \) we may pick a \( P_\xi \)-name \( q(\xi) \) for a condition in \( Q_\xi \) such that

- if \( \xi \in w_{\beta+1} \setminus w_\beta, \beta < \lambda \) (or \( \xi = 0 = \beta \)), then

\[ \models_{\alpha < \lambda} q(\xi) \geq r_\beta(\xi) \text{ and } q(\xi) \models_{\alpha < \lambda} \{ \delta < \lambda : (\exists \varepsilon < \mu_\delta)(\tilde{q}_{\delta, \xi}(\varepsilon) \in \Gamma_{\tilde{G}_\xi}) \} \in D_{\models_{\alpha < \lambda}}. \]

This determines a condition \( q \in \prod_{\alpha < \lambda} \) (with \( \text{dom}(q) = N \cap \gamma \)) and easily \((\forall \beta < \lambda)(p \leq r_\beta \leq q)\) (remember \((*)_{15}\)). For each \( \xi \in N \cap \gamma \) we may pick a \( \tilde{C}_\xi, f_\xi^i \) (for \( i < \lambda \)) such that

\[ q(\xi + 1) \models_{\alpha < \lambda} \{ \delta < \lambda : (\exists \varepsilon < \lambda, \mu_\delta)(\tilde{q}_{\delta, \xi}(\varepsilon) \in \Gamma_{\tilde{G}_\xi}) \} \in D_{\models_{\alpha < \lambda}}. \]

Claim 2.7.1. For each limit ordinal \( \delta < \lambda \), the condition \( q \) forces (in \( \prod_{\alpha < \lambda} \) that

\[ (\forall \xi \in w_\delta)(\exists \delta < \lambda)(C_\xi^\delta \cap \Delta \subseteq \bigtriangleup \text{set}^P(f_\xi^i)) \Rightarrow (\exists \bar{t} \in T_\delta)(\text{rk}_\delta(t) = \gamma \land q_{\bar{t}}(\xi) \in \prod_{\alpha < \lambda}) \].

Proof of the Claim. The proof is essentially the same as that for [RS05, Claim 3.1], however for the sake of completeness we will present it fully. Suppose that \( r \geq q \) and a limit ordinal \( \delta < \lambda \) are such that

\[ (\forall \xi \in w_\delta)(\exists \delta < \lambda)(C_\xi^\delta \cap \Delta \subseteq \bigtriangleup \text{set}^P(f_\xi^i)) \].

For each \( \xi < \gamma \) fix a \( \prod_{\alpha < \lambda} \)-name \( s_{\tilde{t}}^{\lambda} \) for a winning strategy of Complete in \( \prod_{\alpha < \lambda} (Q_\xi, \emptyset_{\tilde{G}_\xi}) \) such that as long as incomplete plays \( \emptyset_{\tilde{G}_\xi} \), Complete answers with \( \emptyset_{\tilde{G}_\xi} \) as well.

We are going to show that there is \( t \in T_\delta \) such that \( \text{rk}_\delta(t) = \gamma \) and the conditions \( q_{\bar{t}}^\delta, r_\alpha \) are compatible. Let \( \langle \varepsilon_\alpha : \alpha \leq \alpha^* \rangle = w_\delta \cap \gamma \) be the increasing enumeration. By induction on \( \alpha \leq \alpha^* \) we can choose conditions \( r_\alpha^\sigma, r_\beta^{\sigma^*} \in \prod_{\alpha} \) and \( t = \langle (t_\varepsilon : \alpha < \alpha^*) \in T_\delta \rangle \) such that letting \( t_\varepsilon^\sigma = \langle (t_\varepsilon : \beta < \alpha) \in T_\delta \rangle \) we have

\[ (\exists \varepsilon_\alpha : \alpha \leq \alpha^*) = w_\delta \cap \gamma \] .

For each \( \beta < \alpha \) and \( \zeta < \varepsilon_\alpha \)

\[ r_\beta^{\sigma^*} \models_{\alpha < \lambda} (r_\beta^{\gamma}(\varepsilon), r_\beta^{\gamma}(\zeta) : \beta < \beta) \] is a legal partial play of \( \prod_{\alpha < \lambda} (Q_\xi, \emptyset_{\tilde{G}_\xi}) \) in which Complete uses her winning strategy \( s_{\tilde{t}}^{\lambda} \).
Suppose that $\alpha \leq \alpha^*$ is a limit ordinal and we have already defined $t^\alpha_\alpha = \langle (t)_{\varepsilon,\beta} : \beta < \alpha \rangle$ and $(r^*_{\beta}, r^*_\beta : \beta < \alpha)$. Let $\xi = \sup(\varepsilon, \beta < \alpha)$. It follows from (H)$_\beta$ that we may find a condition $r^*_\alpha \in P_{\varepsilon,\beta}$ such that $r^*_\alpha | \xi \in P_\xi$ is stronger than all $r^*_\beta$ (for $\beta < \alpha$) and also $r^*_\alpha | (\xi, \varepsilon, \alpha) = r | (\xi, \varepsilon, \alpha)$. Clearly $r | \varepsilon, \alpha \subseteq r^*_\alpha$ and also $q^t_{\alpha, \varepsilon} \subseteq r^*_\alpha$. (remember (H)$_\beta$ for $\beta < \alpha$). Now by induction on $\xi \in [\xi, \varepsilon, \alpha)$ we argue that $q^t_{\alpha, \varepsilon} | \xi \subseteq r^*_\alpha | \xi$. So suppose that $\xi \leq \zeta < \varepsilon, \alpha$ and we know $q^t_{\alpha, \varepsilon} | \zeta \leq r^*_\alpha | \zeta$. It follows from (*)$_3$ + (*)$_5$ + (*)$_6$ that $r^*_\alpha | \xi \models (\forall i < \delta)(r^*_\alpha(\xi) \leq p^*_{\alpha, \varepsilon}(\xi) \leq q^t_{\alpha, \varepsilon}(\xi))$ and therefore we may use (*)$_{14}$ to conclude that

$$r^*_\alpha | \xi \models p^*_{\alpha, \varepsilon}(\xi) \leq q^t_{\alpha, \varepsilon}(\xi) \leq q(\xi) \leq r(\xi) = r^*_\alpha(\xi).$$

Finally we let $r^*_{\alpha, \varepsilon} \in P_{\varepsilon, \alpha}$ be a condition such that for each $\zeta < \varepsilon, \alpha$

$$r^*_{\alpha, \varepsilon} | \zeta \models p^*_{\alpha, \varepsilon}(\zeta) : \beta < \zeta \models (r^*_{\alpha, \varepsilon}(\zeta))$$

Now suppose that $\alpha = \beta + 1 \leq \alpha^*$ and we have already defined $r^*_{\beta, \varepsilon}, r^*_{\beta} \in P_{\varepsilon, \beta}$ and $i^\beta \in T_\beta$. It follows from the choice of $q$ and (H)$_a + (H)$_b + (**)$_{11}$ that

$$r^*_{\alpha, \varepsilon} | \varepsilon, \alpha \models (\forall \varepsilon < \mu\beta)(q^t_{\alpha, \varepsilon}(\varepsilon, \beta) \in \Gamma_{\varepsilon, \beta})$$

Therefore we may choose $\varepsilon = (t)_{\varepsilon, \alpha} < \mu\beta$ (thus defining $t^\alpha_\alpha$) and a condition $r^*_\alpha \in P_{\varepsilon, \alpha}$ such that

- $r^*_\beta \leq r^*_\alpha | \varepsilon, \beta$ and

$$r^*_\alpha | \varepsilon, \beta \models r^*_\alpha(\varepsilon, \beta) \geq r(\varepsilon, \beta) \& r^*_\alpha(\varepsilon, \beta) \geq q^t_{\alpha, \varepsilon}(\varepsilon, \beta),$$

- $r^*_\alpha | \varepsilon, \alpha \models r(\varepsilon, \alpha) = r | \varepsilon, \alpha$.

Exactly like in the limit case we argue that $r^*_\alpha$ has the desired properties and then in the same manner as there we define $r^*_\alpha^*$.

We finish the proof of the claim noting that $t = t^\alpha_\alpha \in T_\beta$ and the condition $r^*_\alpha$ are such that $r^*_\alpha \geq r$ and $r^*_\beta \geq q^t_{\alpha, \varepsilon}$. \qed

Let us use 2.7.1 to argue that $q$ is $(N, P_{\varepsilon, \gamma})$–generic. To this end suppose $\tau \in N$ is a $P_{\varepsilon, \gamma}$–name for an ordinal, say $\tau = \tau_{\alpha, \gamma}, \alpha < \lambda$, and let $q^\tau \models q$. Since $P_{\varepsilon, \gamma}$ is strategically (<$\lambda$)–complete we may build an increasing sequence $\langle q^\iota_i : i < \lambda \rangle$ of conditions above $q^\tau$ and a sequence $\langle C^\iota_i, f^\iota_i : i \in N \cap \gamma, i < \lambda \rangle$ such that $C^\iota_i \in D \cap V, f^\iota_i \in \delta^\lambda$ and for each $\xi \in w_i$

$$q^\iota_i \models (\forall j \leq i)(C^\iota_j = C^\iota_j \& f^\iota_j | i = f^\iota_j | i).$$

The set $\{ \delta \in \Delta : (\forall \xi \in w_i)(\forall j < \delta)(\delta \in (i^\iota_i) \cap \Delta \text{ set}(f^\iota_i)) \}$ is in $D$, so we may choose a limit ordinal $\delta > \alpha$ such that for each $\xi \in w_\delta$ we have

$$\delta \in \Delta \text{ set}(f^\iota_i) \forall i \in \lambda$$

Then $q^\iota_\delta \models (\forall \xi \in w_\delta)(\delta \in (i^\iota_i) \cap \Delta \text{ set}(f^\iota_i))$ and therefore, by 2.7.1,

$$q^\iota_\delta \models (\exists t \in T_\delta)(rk_\delta(t) = \gamma \& q^\iota_\delta | \xi \in P_{\varepsilon, \gamma}).$$

Using (*)$_{13}$ we conclude $q^\iota_\delta \models \tau \in Z_\delta$ and hence $q^\iota_\delta \models \tau \in N$. \qed
Remark 2.8. Naturally, we want to apply Theorem 2.7 to $\gamma = \lambda^{++}$ in a model where $2^\lambda = \lambda^+$, so one may ask if the assumptions (3) + (4) of 2.7 can be satisfied in such a universe. But they are not so unusual: suppose that we start with $V \models \lambda^{<\lambda} = \lambda$ & $2^\lambda = \lambda^+$. Consider the following forcing notion $\mathbb{C}_\lambda^\lambda$.

A condition $p \in \mathbb{C}_\lambda^\lambda$ is a function $p : \text{dom}(p) \rightarrow 2$ such that $\text{dom}(p) \subseteq \lambda^+ \times \lambda$ and $\|\text{dom}(p)\| < \lambda$.

The order is the inclusion.

Plainly, $\mathbb{C}_\lambda^\lambda$ is a $(<\lambda)$–complete $\lambda^+$–cc forcing notion of size $\lambda^+$. Suppose now that $G \subseteq \mathbb{C}_\lambda^\lambda$ is generic over $V$ and let us work in $V[G]$. Put $f = \bigcup G$ (so $f : \lambda^+ \times \lambda \rightarrow 2$) and for $\alpha < \lambda^+$ and $i < 2$ let $A^i_\alpha = \{ \xi < \lambda : f(\alpha, \xi) = i \}$. For a function $h : \lambda^+ \rightarrow 2$ let $U_h$ be the normal filter generated by the family $\{ A^i_\alpha : \alpha < \lambda^+ \}$. One easily verifies that each $U_h$ is a (proper) normal filter and plainly if $h, h' : \lambda^+ \rightarrow 2$ are distinct, then $h(\alpha) = 0, h'(\alpha) = 1$, then $A^0_\alpha \in U_h$ and $\lambda \setminus A^0_\alpha = A^1_\alpha \in U_{h'}$.

3. Noble iterations

The iteration theorems 1.10 and 2.7 have one common drawback: they assume that $\lambda$ is strongly inaccessible. In this section we introduce a property slightly stronger than being $B$–bounding over $p$ and we show the corresponding iteration theorem. The main gain is that the only assumption on $\lambda$ is $\lambda = \lambda^{<\lambda}$.

Definition 3.1. Let $Q = (Q, \leq)$ be a forcing notion and $p = (\bar{P}, \lambda, D)$ be a $\mathcal{D}\ell$–parameter on $\lambda$.

(1) For a condition $p \in Q$ we define a game $\mathbb{D}_{\bar{P}}^{B+}(p, Q)$ between two players, Generic and Antigeneric, as follows. A play of $\mathbb{D}_{\bar{P}}^{B+}(p, Q)$ lasts $\lambda$ steps during which the players construct a sequence $\langle f_\alpha, X_\alpha, \bar{a}_\alpha, q_\alpha : \alpha < \lambda \rangle$ such that

(a) $f_\alpha : \alpha \rightarrow Q$ and $f_\beta \subseteq f_\alpha$ for $\beta < \alpha$,

(b) $X_\alpha \subseteq P_\alpha$ and for every $\eta \in X_\alpha$ the sequence $\langle f_\alpha(\eta(\xi)) : \xi < \alpha \rangle \subseteq Q$ has an upper bound in $Q$ and if $\eta_0, \eta_1 \in X_\alpha$ are distinct, then for some $\xi < \alpha$ the conditions $f_\alpha(\eta_0(\xi)), f_\alpha(\eta_1(\xi))$ are incompatible,

(c) $p_\alpha = (p^\alpha_\eta : \eta \in X_\alpha) \subseteq Q$ is a system of conditions in $Q$ such that

\[ (\forall \xi < \alpha) (f_\alpha(\eta(\xi)) \leq p^\alpha_\eta) \] for $\eta \in X_\alpha$,

(d) $q_\alpha = (q^\alpha_\eta : \eta \in X_\alpha) \subseteq Q$ is a system of conditions in $Q$ such that

\[ (\forall \eta \in X_\alpha) (p^\alpha_\eta \leq q^\alpha_\eta) \] The choices of the objects listed above are done so that at stage $\alpha < \lambda$ of the play:

(8) first Generic picks a function $f_\alpha : \alpha \rightarrow Q$ with the property described in (a) above (so if $\alpha$ is limit, then $f_\alpha = \bigcup_{\beta < \alpha} f_\beta$). She also chooses $X_\alpha$,

$\bar{a}_\alpha$ satisfying the demands of (b)+(c) (note that $X_\alpha$ could be empty).

(9) Then Antigeneric decides a system $q_\alpha$ as in (d).

At the end, Generic wins the play $\langle f_\alpha, X_\alpha, \bar{a}_\alpha, q_\alpha : \alpha < \lambda \rangle$ if and only if

$\mathbb{D}_{\bar{P}}^{B+}$ there is a condition $p^* \in Q$ stronger than $p$ and such that

$\models Q \forall \alpha < \lambda : (\exists \eta \in X_\alpha) (q^\alpha_\eta \in \Gamma_Q) \in D[Q]$.

(2) A forcing notion $Q$ is $B$–noble over $p$ if it is strategically $(<\lambda)$–complete and Generic has a winning strategy in the game $\mathbb{D}_{\bar{P}}^{B+}(p, Q)$ for every $p \in Q$. 
Note that in the above definition we assumed that $\bar{P} = (P_\delta : \delta < \lambda)$. This was caused only to simplify the description of the game – if the domain of $\bar{P}$ is $S \in D$, then we may extend it to $\lambda$ in some trivial way without changing the resulting properties.

Observation 3.2. If $p$ is a $\mathcal{Dl}$-parameter and a forcing notion $Q$ is $B$-noble over $p$, then $Q$ is reasonably $B$-bounding over $p$.

Theorem 3.3. Assume that

1. $\lambda = \lambda^+\lambda$ and $p = (\bar{P}, \lambda, D)$ is a $\mathcal{Dl}$-parameter on $\lambda$, and
2. $\bar{Q} = \langle \mathbb{P}_\xi, \bar{Q}_\xi : \xi < \gamma \rangle$ is a $\lambda$-support iteration such that for every $\xi < \gamma$,

$$\Vdash_{\mathbb{P}_\xi} "\bar{Q}_\xi \text{ is } B\text{-noble over } p \bar{P}[\mathbb{P}_\xi] ".$$

Then

(a) $\mathbb{P}_\gamma = \lim(\bar{Q})$ is $\lambda$-proper,

(b) for each $\mathbb{P}_\gamma$-name $\bar{\tau}$ for a function from $\lambda$ to $V$ and a condition $p \in \mathbb{P}_\gamma$ there are $q \in \mathbb{P}_\gamma$ and $(A_\alpha : \alpha < \lambda)$ such that $|A_\alpha| < \lambda$ (for $\alpha < \lambda$) and

$$q \Vdash_{\mathbb{P}_\gamma} "\{ \alpha < \lambda : \bar{\tau}(\alpha) \in A_\alpha \} \in D[\mathbb{P}_\gamma] ".$$

Proof. (a) Assume that $N \prec (\mathcal{H}(\chi), \in, <^*_\gamma)$ is such that $<^\lambda N \subseteq N$, $|N| = \lambda$ and $\bar{Q}, \bar{P}, \ldots \in N$. Let $p \in N \cap \mathbb{P}_\gamma$. Choose $\bar{N} = \langle N_\delta : \delta < \lambda \rangle$ and $\bar{\alpha} = \langle \alpha_\delta : \delta < \lambda \rangle$ such that

- $\bar{N}$ is an increasing continuous sequence of elementary submodels of $N$,

- $\bar{\alpha}$ is an increasing continuous sequence of ordinals below $\lambda$,

- $N = \bigcup N_\delta$, $\bar{Q}, \bar{P}, \ldots \in N_0$, $\delta \subseteq N_\delta$, $P_\delta \subseteq N_{\delta+1}$, $N(|\delta + 1) \in N_{\delta+1}$, $|N_\delta| < \lambda$ and

- $\alpha_\delta + \text{otp}(N_\delta \cap \gamma) + 888 < \alpha_{\delta+1}, \bar{\alpha}(\delta + 2) \in N_{\delta+1}$.

Put $w_\delta = N_\delta \cap \gamma$ and for each $\xi < \gamma$ let $\bar{\text{st}}_{\xi}^0$ be the $<^\lambda_{\xi, \gamma}$-first $\mathbb{P}_\xi$-name for a winning strategy of Complete in $\mathcal{Dl}_\lambda^0(\bar{Q}_\xi, \bar{\theta}_\xi)$ such that it instructs Complete to play $\bar{\theta}_\xi$, as long as her opponent plays $\bar{\theta}_\xi$. We also assume that whenever possible, $\bar{\theta}_\xi$ is the $<^\lambda_{\xi, \gamma}$-first name for the answer by $\bar{\text{st}}_{\xi}^0$ to a particular sequence of names. Note that $\langle \bar{\text{st}}_{\xi}^0 : \xi < \gamma \rangle \in N_0$.

By induction on $\delta < \lambda$ we will construct

$$(\otimes)_\delta \bar{T}_\delta, \bar{p}_\delta^\bar{q}, \bar{q}_\delta, r_\delta, r, \bar{r}_\delta \text{ and } f_{\delta, \xi, \bar{\Delta}_\xi, \bar{\nu}_\xi, \bar{\mu}_\xi, \bar{g}_\xi, \bar{g}_\xi, \bar{\text{st}}_{\xi}^0, \bar{p}_\delta^{\bar{\text{st}}_{\xi}^0} \text{ for } \xi \in w_\delta}$$

so that the following conditions $(*)_0-(*)_16$ are satisfied.

$(*)_0$ Objects listed in $(\otimes)_\delta$ form the $<^\lambda_{\xi, \gamma}$-first tuple with the properties described in $(*)_1-(*)_16$ below. Consequently, the sequence

$$(\text{objects listed in } (\otimes)_\delta : \epsilon < \delta)$$

is definable from $N[\alpha_\delta, \bar{\alpha}, \bar{Q}, \bar{P}, p$ (in the language $\mathcal{L}(\in, <^\lambda)$), so if $\delta = \alpha_\delta$ is limit, then this sequence belongs to $N_{\delta+1}$. Also, objects listed in $(\otimes)_\delta$ are known after stage $\delta$ (and they all belong to $N$).

$(*)_1 \bar{r}_\delta, r_\delta \in \mathbb{P}_\gamma, w_\delta \subseteq \text{dom}(\bar{r}_\delta) = \text{dom}(r_\delta)$ and $r_\delta (\xi) = \bar{r}_\delta (\xi) = r_\delta (\xi) = p(\xi)$ for $\xi \in w_\delta$.

$(*)_2$ For each $\epsilon < \delta < \lambda$ we have $(\forall \xi \in w_{\epsilon+1})(r_\delta (\xi) = \bar{r}_\delta (\xi))$ and $p \leq r_\delta \leq r_\delta \leq r_\delta \leq r_\delta$.
The demands \((*)\) of Theorem \(2\) are satisfied if \(r_\delta|\xi| \preceq \varphi_\xi\)
the sequence \(\{r_\varepsilon(\xi), r_\varepsilon(\xi) : \varepsilon \leq \delta\}\) is a legal partial play of 
\(\bar{\delta}(\xi, \varphi_\xi)\) and Complete follows \(s(t)\).

and if \(\xi \in w_{3+n}\), then \(\varphi_\xi\) is the \(<\chi\)-first \(\varphi_\xi\)-name for a winning strategy of 
Generic in \(\bar{\delta}(\xi, \varphi_\xi)\). (And for \(\xi \in w_0\), \(\varphi_\xi\) is the \(<\chi\>-\text{first} \(\varphi_\xi\)-name for a winning strategy of Generic in \(\bar{\delta}(\xi, \varphi_\xi)\).)
Note that \(\varphi_\xi \in N_0\) for 
\(\xi \in w_0\) and \(\varphi_\xi \in N_{3+n}\) for \(\xi \in w_3\).

\((*)\) of Theorem \(2\) is a \((w_3, 1)\)-tree, \(T_3 \subseteq \bigcup_{\alpha \leq \gamma} T_\alpha\). (Note that we
do not require here that \(T_3\) is standard, so some chains in \(T_3\) may have no
\(<\alpha\)-bounds.)

\((*)\) of Theorem \(2\) is a \((w_3, 1)\)-tree, \(T_3 \subseteq \bigcup_{\alpha \leq \gamma} T_\alpha\).
\(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\). (And for \(\xi \in w_0\), \(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\).)

For \(\xi \in w_3\), we have that \(\varphi_\delta\) is such an upper bound.

\((*)\) of Theorem \(2\) is a \((w_3, 1)\)-tree, \(T_3 \subseteq \bigcup_{\alpha \leq \gamma} T_\alpha\).
\(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\). (And for \(\xi \in w_0\), \(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\).)

The demands \((*)\) of Theorem \(2\) are satisfied if \(\delta = \alpha_\delta\) is a limit ordinal.

\((*)\) of Theorem \(2\) is a \((w_3, 1)\)-tree, \(T_3 \subseteq \bigcup_{\alpha \leq \gamma} T_\alpha\).
\(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\). (And for \(\xi \in w_0\), \(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\).)

The demands \((*)\) of Theorem \(2\) are satisfied if \(\delta = \alpha_\delta\) is a limit ordinal.

\((*)\) of Theorem \(2\) is a \((w_3, 1)\)-tree, \(T_3 \subseteq \bigcup_{\alpha \leq \gamma} T_\alpha\).
\(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\). (And for \(\xi \in w_0\), \(\varphi_\delta\) is the \(<\chi\>-\text{first} \(\varphi_\delta\)-name for a winning strategy of
Generic in \(\bar{\delta}(\xi, \varphi_\delta)\).)

The demands \((*)\) of Theorem \(2\) are satisfied if \(\delta = \alpha_\delta\) is a limit ordinal.
(*15) If \( t_0, t_1 \in T_3 \), \( \text{rks}(t_0) = \text{rks}(t_1) \) and \( \xi \in w_3 \cap \text{rks}(t_0) \), \( t_0 \upharpoonright \xi = t_1 \upharpoonright \xi \) but \( (t_0)_{\xi} \neq (t_1)_{\xi} \), then

\[ p^\xi_{t_0|\xi} \upharpoonright p^\xi_{t_1|\xi} \] "the conditions \( p^\xi_{t_0}(\xi), p^\xi_{t_1}(\xi) \) are incompatible ".

(*16) If \( \tau \in N_3 \) is a \( \mathbb{P}_{\gamma} \)-name for an ordinal and \( t \in T_3 \) satisfies \( \text{rks}(t) = \gamma \), then the condition \( q^\gamma_{t} \) forces a value to \( \tau \).

The rule (*0) (and conditions (*1)−(*16)) actually fully determines our objects, but we should argue that at each stage there exist objects with properties listed in (*1)−(*16).

Suppose we have arrived to a stage \( \delta < \lambda \) of the construction and all objects listed in \((=\delta)\) for \( \beta < \delta \) have been determined so that all relevant demands are satisfied, in particular the sequence \( \text{otp}(\xi) \) is a tree of conditions, \( \overline{\xi} \) and \( X^\xi \bar{\xi} \) are given to \( \text{Generic} \) by \( \mathbf{st}_\xi \). We also pick

\[ \text{objects listed in } (\geq \delta) : \gamma < \delta \] as the store for \( \text{Generic} \).

Note that

\[ \text{objects listed in } (\geq \delta) : \gamma < \delta \]

Now by induction on \( \xi \in w_3 \cup \{ \gamma \} \) we will choose \( \{ t \in T_3 : \text{rks}(t) \leq \xi \} \) and \( p^\beta \xi \) and auxiliary objects \( p^\beta \xi \) so that, in addition to demands (*3)−(*13) we also have:

\[ (\geq \delta) \iff [p^\beta \xi : t \in T_3 \& \text{rks}(t) \leq \xi] \subseteq \mathbb{P}_\xi \text{ is a tree of conditions, } p^\beta \xi \leq p^\delta \xi \]

and \( \text{dom}(p^\beta \xi) \supseteq (\text{dom}(p) \cup \bigcup_{\gamma < \delta} \text{dom}(r_{\gamma}) \cup w_3) \cap \text{rks}(t) \) whenever \( t \in T_3 \),

\[ \text{rk}_3(t) \leq \xi, \text{ and} \]

(*18) if \( \xi_0 < \xi_1 \) are from \( w_3 \cup \{ \gamma \} \), \( t \in T_3 \), \( \text{rk}_3(t) = \xi_0 \), then \( p^\beta \xi_0 ? p^\beta \xi_1 \), and

(*19) if \( t \in T_3 \), \( \text{rk}_3(t) = \xi \), then \( \text{dom}(p^\beta \xi) = \text{dom}(p^\delta \xi) \) and for \( \beta \in \text{dom}(p^\delta \xi) \) we have

\[ p^\beta \xi \upharpoonright \beta \upharpoonright p^\beta \xi \] "the sequence \( \langle p^\beta \xi(\beta), p^\delta \xi(\beta) : \xi \in (w_3 \cup \{ \gamma \}) \cap (\xi + 1) \rangle \) is a legal partial play of \( 2^\delta(\xi_0) \) \( \text{in which Complete uses the winning strategy } \mathbf{st}_0^\beta \)."

To take care of clause (*13), each time we pick an object, we choose the \( \langle \xi \rangle \)−first one with the respective property.

CASE 1: \( \text{otp}(w_3 \cap \xi) = \xi + 1 \) is a successor ordinal. Let \( \xi_0 = \max(w_3 \cap \xi) \) and suppose that we have defined \( T^\xi = \{ t \in T_3 : \text{rks}(t) \leq \xi_0 \} \) \( \text{and } \text{rks}(\xi_0) = \text{rk}^\delta \xi_0 \text{ satisfying the relevant demands of (*1)−(*10)). Let } t \in T_3 \text{ be such that } \text{rk}_3(t) = \xi_0 \). It follows from (*9) that either \( p^\beta \xi_0 \upharpoonright \text{" } X^\delta_{\xi_0} = 0 \text{" or } p^\beta \xi_0 \upharpoonright \text{" } X^\delta_{\xi_0} = X^\delta \text{" for some non-empty set } X^\delta \subseteq P_3. \) In the former case stipulate
$X_t = \{*\}$. Note that necessarily $X_t^d \subseteq N_{\delta+1}$ and $X_t^c \subseteq N_{\delta+\xi+2}$ (remember $(*)_{13}$).

We declare that

$$\{t \in T_a : r_k(t) \leq \xi\} = T^* \cup \{t \cup \{(\xi_0, a)\} : t \in T^* \& \ r_k(t) = \xi_0 \& a \in X_t^d\}.$$ 

Plainly, $|\{t \in T_a : r_k(t) \leq \xi\}| < \lambda$ and even $\{t \in T_a : r_k(t) \leq \xi\} \subseteq N_{\delta+\xi+1}$ and $\{t \in T_a : r_k(t) \leq \xi\} \subseteq N_{\delta+\xi+2}$ (again, by $(*)_{13}$). Choose a tree of conditions $\bar{p}^+ = \{p_t^+ : t \in T_a \& \ r_k(t) \leq \xi\} \subseteq \mathbb{P}_\xi$ so that

1. $\text{dom}(p_t^+) \supseteq \left(\text{dom}(p) \cup \bigcup_{\varepsilon < \delta} \text{dom}(r_x) \cup w_0\right) \cap r_k(t)$ for $t \in T_a$, $r_k(t) \leq \xi$,
2. if $t \in T_a$, $r_k(t) < \xi$ then $p_t^+ = p_t^{\xi, \alpha}$,
3. if $t \in T_a$, $r_k(t) = \xi$ and $(t)_{\xi_0} \neq *$, then $p_t^+(\xi_0)$ is a $\mathbb{P}_\xi$-name such that $p_t^{\xi, \alpha} \forces \check{\xi}_{\xi_0}(t)_{\xi_0} \leq p_t^+(\xi_0)$, 
4. if $t \in T_a$, $r_k(t) = \xi$ and $\beta \in \text{dom}(p_t^+) \setminus (\xi_0 + 1)$, then $p_t^+ \forces \beta \in \check{\mathbb{P}}_\beta$ if the set $\{r_x(\beta) : \varepsilon < \delta\} \cup \{p(\beta)\}$ has an upper bound in $\mathbb{Q}_\beta$, then $p_t^+(\beta)$ is such an upper bound.

(Note: $\bar{p}^+ \in N_{\delta+\xi+2}$.)

Next we may use Proposition 0.4 to pick a tree of conditions $\bar{p}^+ = \{p_t^{\xi, \alpha} : t \in T_a \& \ r_k(t) \leq \xi\}$ such that $\bar{p}^+ \subseteq \bar{p}^{\xi, \alpha}$ and

1. if $\xi < \gamma$, $t \in T_a$, $r_k(t) = \xi$, then either $p_t^{\xi, \alpha} \forces X_\gamma \xi = \emptyset$ or for some non-empty set $X_t^\xi \subseteq P$ we have $p_t^{\xi, \alpha} \forces X_\gamma \xi = X_t^\xi$.

(Again, by our rule of picking “the $<\check{\gamma}$ first”, $p_t^{\xi, \alpha} \in N_{\delta+\xi+2}$.) Then we choose a tree of conditions $\bar{p}^+ = \{p_t^{\xi, \alpha} : t \in T_a \& \ r_k(t) \leq \xi\}$ so that $\bar{p}^+ \leq \bar{p}^{\xi, \alpha}$ and for every $t \in T_a$ with $r_k(t) = \xi$ we have $\text{dom}(p_t^{\xi, \alpha}) = \text{dom}(p_t^{\xi, \alpha})$ and $\beta \in \text{dom}(p_t^{\xi, \alpha})$.

$p_t^{\xi, \alpha}(\beta)$ is the $<\check{\gamma}$ first $\mathbb{P}_\beta$-name for a condition in $\mathbb{Q}_\beta$ such that

$$p_t^{\xi, \alpha}(\beta)_{\check{\mathbb{P}}_\beta}. \text{ is given to Generic by } s_{\beta}^0 \text{ as the answer to } p_t^{\xi, \alpha}(\beta)_{\check{\mathbb{P}}_\beta} \text{ when } \forall \varepsilon \in \mathbb{P}_\xi \subseteq \check{\mathbb{P}}_\xi.$$

Note that, by the rule of picking “the $<\check{\gamma}$ first”, $\bar{p}^{\xi, \alpha} \in N_{\delta+\xi+2}$. It should be also clear that $p_t^{\xi, \alpha}, p^{\xi, \alpha}$ satisfy all the relevant demands stated in $(*)_{19}$.

**Case 2:** $\text{otp}(w_\delta \cap \xi) = \xi$ is a limit ordinal.

Suppose we have defined $\{t \in T_a : r_k(t) \leq \varepsilon\}$ and $p^{\varepsilon, \alpha}, p^{\varepsilon, \xi, \alpha} : \varepsilon \in w_\delta \cap \xi$.

By our rule of choosing “the $<\check{\gamma}$ first objects”, we know that the sequence $\{t \in T_a : r_k(t) \leq \varepsilon\} \setminus \text{dom}(p^{\varepsilon, \alpha}) \subseteq N_{\delta+\xi+1}$.

We also know that $\{t \in T_a : r_k(t) \leq \xi\} \subseteq N_{\delta+\xi+1}$. Let $T^+$ be the set of all limit branches in $\{t \in T_a : r_k(t) < \xi\}, s = (s_\varepsilon : \varepsilon \in w_\delta \cap \xi)$ such that $s|\xi = (s_\varepsilon : \varepsilon \in w_\delta \cap \xi) \in \{t \in T_a : r_k(t) \leq \varepsilon\}$ for $\varepsilon \in w_\delta \cap \xi$. (Of course, $T^+ \in N_{\delta+\xi+1}$.) We put

$$\{t \in T_a : r_k(t) \leq \xi\} = \{t \in T_a : r_k(t) < \xi\} \cup (T^+ \cap N_{\delta+\xi+1}) \subseteq N_{\delta+\xi+2}.$$ 

Due to $(*)_{19}$ at stages $\varepsilon \in w_\delta \cap \xi$, we may choose a tree of conditions $\bar{p}^+ = \{p_t^+ : t \in T_a \& \ r_k(t) \leq \xi\} \subseteq \mathbb{P}_\xi$ such that

1. $\text{dom}(p_t^+) \supseteq \left(\text{dom}(p) \cup \bigcup_{\varepsilon < \delta} \text{dom}(r_x) \cup w_0\right) \cap r_k(t)$ for $t \in T_a$, $r_k(t) \leq \xi$,
2. if $t \in T_a$, $r_k(t) = \xi_0 < \xi$ then $\text{dom}(p_t^+) \supseteq \text{dom}(p_t^{\xi_0, \alpha})$ and for each $\beta \in \text{dom}(p_t^+) \cap \xi_0$ we have $p_t^+ \forces \beta \in \check{\mathbb{P}}_\beta$, $p^{\xi_0, \alpha}(\beta) \leq p_t^+(\beta)$, and
Also like in that case we choose $\bar{p}$.

Claim 3.3.1.

Proof of the Claim.

Then, like in the successor case, we may find a tree of conditions $\bar{p}^{<\xi} = \langle p_i^{<\xi} : t \in T_\delta \rangle$ and $\bar{r}_k(t) \leq \xi$ such that $\bar{p}^* \leq \bar{p}^{<\xi}$ and

- if $\xi < \gamma$, $t \in T_\delta$, $\bar{r}_k(t) = \xi$, then either $p_i^{<\xi} \models \Gamma_{\delta,\xi} = \emptyset$ or for some non-empty set $X^i_{\delta,\xi} \subseteq P_\delta$ we have $p_i^{<\xi} \models \Gamma_{\delta,\xi} = X^i_{\delta,\xi}$.

Also like in that case we choose $p^{<\xi} = \langle p_i^{<\xi} : t \in T_\delta \& \bar{r}_k(t) \leq \xi \rangle$. Clearly, all relevant demands in $(\ast)_9^{\ast}(\ast)_7^{\ast}$ are satisfied.

The last stage of the above construction gives us a tree $T_\delta = \{ t \in T_\delta : \bar{r}_k(t) \leq \gamma \}$ and a tree of conditions $\bar{q}^{<\gamma} = \langle q_i^{<\gamma} : t \in T_\delta \rangle$. Since $T_\delta \subseteq N_{\delta + \sigma_p(w_\delta) + 1}$, we know that $|T_\delta| < \lambda$ so we may apply Proposition 0.4 to get a tree of conditions $q^\delta = \langle q_i^\delta : t \in T_\delta \rangle \geq \bar{q}^\delta$ such that $(\ast)_7^{\ast}\gamma$ is satisfied. Remembering $(\ast)_7^{\ast} + (\ast)_7^{\ast}$, we easily find $P_\xi$–names $\bar{g}_{\delta,\xi}$ for $\xi \in w_\delta$ such that

- $\models P_\xi \langle \bar{q}_{\delta,\xi} = \langle \bar{q}_{\delta,\xi}(\eta) : \eta \in P_{\delta,\xi} \rangle $ is a system of conditions in $Q_{\xi}$ ",
- $\models P_\xi \langle \forall \eta \in P_{\delta,\xi} \langle \bar{q}_{\delta,\xi}(\eta) \rangle \rangle $ , and
- if $t \in T_\delta$, $\bar{r}_k(t) > \xi$, then $q_i^\delta \models P_\xi \langle X_{\delta,\xi} \neq \emptyset \Rightarrow q_i^\delta(\xi) = \bar{g}_{\delta,\xi}(t) \rangle $.

So then $(\ast)_7^{\ast}$ is satisfied. Now we define $r^{\gamma}_{\delta,\gamma}, r_{\gamma} \in P_\gamma$ essentially by $(\ast)_7^{\ast}-(\ast)_7^{\ast}$ and $(\ast)_8^{\ast}$.

After completing all $\lambda$ stages of the construction, for each $\xi \in N \cap \gamma$ we look at the sequence $(\langle \alpha, \xi \rangle, X_{\alpha,\xi}, P_{\alpha,\xi}, g_{\alpha,\xi} : \alpha \leq \lambda)$. By $(\ast)_7^{\ast}\gamma$, it is a $P_\xi$–name for a play of $\beta^p_{\xi}(n_{\delta,\xi}, Q_{\xi})$ (where $\xi \in w_{\beta+1} \setminus w_\beta$) in which Generic uses her winning strategy $s_{\delta,\xi}$. Therefore, for every $\xi \in N \cap \gamma$ we may pick a $P_\xi$–name $q(\xi)$ for a condition in $Q_{\xi}$ such that

- if $\xi \in w_{\beta+1} \setminus w_\beta, \beta < \lambda$ (or $\xi \in w_\beta, \beta = 0$), then
- $\models P_\xi \langle q(\xi) \geq r_\beta(\xi) \rangle$ and $q(\xi) \models P_\xi \langle \forall \eta \in P_{\delta,\xi} \langle \bar{q}_{\delta,\xi}(\eta) \in \Gamma_{\delta,\xi} \rangle \in D^p[\xi_{\beta+1}] \rangle$. This determines a condition $q \in P_\xi$ (with dom($q$) = $N \cap \gamma$) and easily $(\forall \beta < \lambda)(p \leq r_\beta \leq q)$ (remember $(\ast)_2^{\ast}$). For each $\xi \in N \cap \gamma$ fix $P_{\xi_{\beta+1}}$–names $C_\xi^i, g_i^\xi$ (for $i < \lambda$) such that

\[ q(\xi + 1) \models P_{\xi_{\beta+1}} (\forall \eta \in \lambda) (C_\xi^i \subseteq D \cap V \& g_i^\xi \subseteq \lambda) \& (\forall \delta < \xi) (C_\xi^i \cap \text{set}^p(g_i^\xi))(\exists \eta \in P_{\delta,\xi}) (\bar{q}_{\delta,\xi}(\eta) \in \Gamma_{\delta,\xi}) \] .

Let $B$ be a $P_\gamma$–name for the set \{ $\delta < \lambda : C_{\delta \cap V} \subseteq N_{\delta} \in N_{\delta+1}$ \}. It follows from Lemma 1.6 that $\exists \models P_\gamma \langle B \in D^p[\xi_{\beta+1}] \rangle$.

Claim 3.3.1. If $\alpha_\delta = \delta$ is limit, then

\[ q \models P_\gamma \langle \text{if } \delta \in B \& (\forall \xi \in w_\delta) (\delta \in \Delta C_{i < \lambda}^{\xi} \cap \Delta \text{set}^p(g_i^\delta)) \Rightarrow (\exists t \in T_\delta) (\bar{r}_k(t) = \gamma \& q_i^\delta \in \Gamma_{\delta,\xi}) \rangle . \]

Proof of the Claim. Suppose that $\delta = \alpha_\delta$ is a limit ordinal and a condition $r \geq q$ forces (in $P_\gamma$) that

\begin{align*}
&\text{\scriptsize{(3.3.7) } \delta \in B, \quad \text{and}} \\
&\text{\scriptsize{(3.3.8) } (\forall \xi \in w_\delta) \left(\delta \in \Delta C_{i < \lambda}^{\xi} \cap \Delta \text{set}^p(g_i^\delta)\right).}
\end{align*}
Passing to a stronger condition if necessary, we may also assume that

\((\ast)^{\ast\ast}_{21}\) if \(r' \in \mathbb{P}_\gamma \cap N_\delta\), then either \(r' \leq r\) or \(r', r\) are incompatible in \(\mathbb{P}_\gamma\).

Let \(H^\delta = \{r' \in \mathbb{P}_\gamma \cap N_\delta : r' \leq r\}\). It follows from \((\ast)^{\ast\ast}_{20} + (\ast)^{\ast\ast}_{21}\) that

\((\ast)^{\ast\ast}_{21}\) \(r \Vdash \Gamma_{\mathbb{P}_\gamma} \cap N_\delta = H^\delta\) and \(H^\delta \in N_{3+1}\).

By 3.1(1)(a) + (a) + (7) we may choose a sequence \(\bar{r} = \langle r_\xi, \alpha \rangle : \xi \in w_\delta \land \alpha < \delta \rangle \in N_{3+1}\) such that

- \(\bar{r} \in \mathbb{P}_{\xi}\)-name for an element of \(\mathcal{Q}_\xi, \bar{r}(\xi, \alpha) \in N_\delta\)
- \(\bar{r} \in \mathbb{P}_{\xi} \bar{r}(\xi, \alpha) = f_\delta, \xi(\alpha)\)

Next we choose a sequence \(t^* = \langle (t^*_\xi) : \xi \in w_\delta \rangle \in \prod_{\xi \in w_\delta} (P_\delta \cup \{\ast\})\) so that for each \(\xi \in w_\delta, (t^*_\xi)\) is the \(< \chi\)-first member of \(P_\delta \cup \{\ast\}\) satisfying:

\((\ast)_{22}\) if \(t = t^*_\xi = \langle (t^*_\xi) : \xi \in w_\delta \cap \xi \rangle \in T_\delta\) and

\(\langle \forall \alpha \in \delta \rangle \bar{r}(\xi, \alpha) \in \langle r'(\xi) : r' \in H^\delta\rangle\),

then \((t^*_\xi) \in X_{\delta, \xi}\) (remember \((\ast)_{11}\))

Note that for every \(\xi \in w_\delta \cap \{\gamma\}\) the sequence \(t^*_\xi\) is definable (in \(L(\xi, r_\xi, H^\delta, w_\delta, \xi)\)) from \(p, \bar{r}, H^\delta, w_\delta, \xi\) and \(p^{\delta, \xi} : \bar{r}(\xi) \in w_\delta \cap \xi\). Consequently, if \(\xi \in w_\delta \cap \{\gamma\}\) and \(\xi = \text{otp}(w_\delta \cap \xi)\), then \(t^*_\xi \in N_{3+\xi+1}\). Now, by induction on \(\xi \in w_\delta \cup \{\gamma\}\) we are going to show that \(t^*_\xi \in T_\delta\) and choose conditions \(r^*_\xi, r^*_\xi\) \(r^*_\xi\) in \(\mathbb{P}_{\xi}\) such that

\((\ast)_{23}\) \(q^*_\delta, \xi| \leq r^*_\xi, r^*_\xi | \leq r^*_\xi\) and if \(e \in w_\delta \cap \xi\) then \(r^*_\xi \leq r^*_\xi\) and \(r^*_\xi \leq r^*_\xi\) and for every \(\beta \in \text{dom}(r^*_\xi)\)

\(r^*_\xi \in \mathbb{P}_{\xi} r^*_\xi, r^*_\xi | \leq r^*_\xi, r^*_\xi | \leq r^*_\xi\) and if \(e \in w_\delta \cap \xi\) then \(r^*_\xi \leq r^*_\xi\) and \(r^*_\xi \leq r^*_\xi\) and for every \(\beta \in \text{dom}(r^*_\xi)\)

Suppose that \(\text{otp}(w_\delta \cap \xi)\) is a limit ordinal and for \(e \in w_\delta \cap \xi\) we know that \(t^*_\xi \in T_\delta\) and we have defined \(r^*_\xi, r^*_\xi\). It follows from \((\ast)_3\) that \(t^*_\xi \in T_\delta\). Let \(\beta = \sup(w_\delta \cap \xi)\). It follows from \((\ast)^{\ast\ast}_{23}\) that we may find a condition \(r^*_\xi \in \mathbb{P}_{\xi}\) such that \(r^*_\xi \beta| \subseteq \gamma^*_\xi\) \(\beta \leq \gamma^*_\xi\) is stronger than all \(r^*_\xi\) \((e \in w_\delta \cap \xi)\) and \(r^*_\xi | \beta = r^*_\xi | \beta\).

Clearly \(r^*_\xi | \leq r^*_\xi\) and also \(q^*_\delta, \xi| \leq r^*_\xi\) \(\bar{r}(\xi, \alpha) \in \langle r'(\xi) : r' \in H^\delta\rangle\),

Note by induction on \(\alpha \in \beta, \xi\) we argue that \(\bar{r}(\xi, \alpha) \in \text{dom}(r^*_\xi)\).

So suppose that \(\beta \leq \alpha < \xi\) and we know already that \(q^*_\delta, \xi| \leq r^*_\xi\) \(\bar{r}(\xi, \alpha) \in \text{dom}(r^*_\xi)\) and for every \(\beta \in \text{dom}(r^*_\xi)\)

\(r^*_\xi \alpha | \leq q^*_\delta, \xi| \leq q^*_\delta, \xi| \leq \bar{r}(\xi, \alpha) \alpha| = r^*_\xi(\alpha)\)

as desired. Finally we define \(r^*_\xi \in \mathbb{P}_{\xi}\) essentially by \((\ast)^{\ast\ast}_{23}\).

Now suppose that \(\text{otp}(w_\delta \cap \xi)\) is a successor ordinal and let \(\xi_0 = \max(w_\delta \cap \xi)\). Assume we know that \(t^*_\xi \in T_\delta\) and that we have already defined \(r^*_\xi, r^*_\xi\) in \(\mathbb{P}_{\xi_0}\).

It follows from the choice of \(\xi\) and from \((\ast)^{\ast\ast}_{20}\) that

\(r^*_\xi | \xi_0 | \leq \bar{r}(\xi, \alpha) | \leq q^*_\delta, \xi| \leq r^*_\xi(\alpha) = r^*_\xi(\alpha),\)

Thus we may choose \(r^*_\xi \in \mathbb{P}_{\xi, \xi_0+1}\) and \(\eta \in P_\delta\) such that \(r^*_\xi \leq r^*_\xi, r^*_\xi | \leq r^*_\xi(\xi_0)\) and \(r^*_\xi | \xi_0 \in \mathbb{P}_{\xi_0} \eta \in \mathcal{X}_{\delta, \xi_0} \land \bar{q}^*_\delta, \xi_0| \leq r^*_\xi(\xi_0)\). Then \(r^*_\xi | (\xi_0 + 1) \leq r^*_\xi\).
and (by \((\ast)r, 3.1(1)(c,d)) r^* | \xi_0 \vDash P_{\xi_0} (\forall \alpha < \delta) (\tau(\xi_0, \eta(\alpha)) \leq r^*(\xi_0))\) and hence (by \((\ast)_{24}'\)) \(r | \xi_0 \vDash \tau(\xi_0, \eta(\alpha)) \leq r(\xi_0)\) for all \(\alpha < \delta\). Therefore
\[(\ast)_{24}'\mbox{ for all } \alpha < \delta, \tau(\xi_0, \eta(\alpha)) \in \{r(\xi_0) : r' \in H^d\}.
\]
Since \(p_{\xi_0}^{\delta, \xi_0} \leq r^* | \xi_0\) (remember \((\ast)_{24}'\) for \(\xi_0\) and \((\ast)_{s}\)) we may use \((\ast)_{11}\) to conclude that for some non-empty set \(X \subseteq P_\delta\) we get \(p_{\xi_0}^{\delta, \xi_0} \vDash X = X_{\delta, \xi_0}\) and \(\eta \in X\) satisfies (ii) of \((\ast)_{22}\). Hence \((t^*)_{\xi_0} \in X\) is such that
\[(\ast)_{24}'\mbox{ for all } \alpha < \delta, \tau(\xi_0, (t^*)_{\xi_0}(\alpha)) \in \{r(\xi_0) : r' \in H^d\},
\]
and in particular \(t^* | \xi \in T_\delta\) (remember \((\ast)_{11}\)). We claim that \((t^*)_{\xi_0} = \eta). If not, then by 3.1(1)(b) we have
\[r^* | \xi_0 \vDash (\exists \alpha < \delta) (\tau(\xi_0, (t^*)_{\xi_0}(\alpha)), \tau(\xi_0, \eta(\alpha))\) are incompatible in \(Q_{\xi_0}\),
\]
so we may pick \(\alpha < \delta\) and a condition \(r^+ \in P_{\xi_0}\) such that \(r^* | \xi_0 \leq r^+\) and
\[r^+ \vDash \tau(\xi_0, (t^*)_{\xi_0}(\alpha)), \tau(\xi_0, \eta(\alpha))\) are incompatible in \(Q_{\xi_0}\).
\]
However, \(r^+ \vDash \tau(\xi_0, (t^*)_{\xi_0}(\alpha)) \leq r(\xi_0) \& \tau(\xi_0, \eta(\alpha)) \leq r(\xi_0)\) ” (by \((\ast)_{24}'' + \ast)_{24}'\), a contradiction.

Now we define \(r_{\xi}^* \in \mathbb{P}_{\xi}\) so that \(r_{\xi}^*(\xi_0 + 1) = r^*\) and \(r_{\xi}^*(\xi_0, \xi) = r(\xi_0, \xi).\) By the above considerations and \((\ast)_{14}\) we know that \(q_{\xi}^d | (\xi, \xi) \leq r_{\xi}^* | (\xi_0 + 1)\).

Exactly like in the case of limit otp(\(w_2\cap \xi\)) we argue that \(q_{\xi}^d \leq r_{\xi}^*\). Finally, we choose \(r_{\xi}^* \in \mathbb{P}_{\xi}\) by \((\ast)_{23}\).

The last stage \(\gamma\) of the inductive process described above shows that \(t^* \in T_\delta\) and \(q_{\xi}^d \leq r_{\xi}^*, r \leq r_{\xi}^*.\) Now the claim readily follows.

We finish the proof of part (a) of the theorem exactly like in the proof of 2.7.

(b) Included in the proof of the first part.

4. EXAMPLES AND COUNTEREXAMPLES

Let us note that our canonical test forcing \(Q_{\pi}^E\) is \(B\)-noble:

**Proposition 4.1.** Assume that \(E, E\) are as in 1.11 and \(p = (P, S, D)\) is a \(\mathcal{D}\)-parameter on \(\lambda\) such that \(\lambda \\setminus S \in E\). Then the forcing \(Q_{\pi}^E\) is \(B\)-noble over \(p\).

**Proof.** The proof is a small modification of that of 1.12(2). First we fix an enumeration \(\langle \nu_\alpha : \alpha < \lambda \rangle = \langle \lambda \rangle (\alpha < \lambda, \lambda), \) and for \(\alpha < \lambda\) let \(f(\alpha) \in Q_{\pi}^E\) be a condition such that root(\(f(\alpha)\) = \(\nu_\alpha\) and
\[(\forall \nu \in f(\alpha)) ((\nu \triangleleft \nu \Rightarrow \text{succ}(f(\alpha)) = \lambda)).\]

Let \(p \in Q_{\pi}^E\). Consider the following strategy \(st\) of Generic in \(Q_{\pi}^E(p, Q_{\pi}^E)\). In the course of the play \(Generic\) is instructed to build aside a sequence \(\langle T_\xi : \xi < \lambda \rangle\) so that if \(\langle f_\xi, X_\xi, p_\xi, q_\xi : \xi < \lambda \rangle\) is the sequence of the innings of the two players, then the following conditions (a)–(d) are satisfied.

(a) \(T_\xi \subseteq Q_{\pi}^E\) and if \(\xi < \zeta < \lambda\) then \(p = T_0 \supseteq T_\xi \supseteq T_\zeta\) and \(T_\xi \cap \zeta = T_\zeta \cap \zeta\).

(b) If \(\xi < \lambda\) is limit, then \(T_\xi = \bigcap_{\zeta < \xi} T_\zeta\).

(c) If \(\xi \in S\), then
\[\bullet f_\xi = f | \xi\mbox{ and } X_\xi \subseteq P_\xi\mbox{ is a maximal set (possibly empty)}\] (a) for each \(\eta \in X_\xi\) the family \(\{f(\eta(\alpha)) : \alpha < \xi\} \cup \{T_\xi\}\) has an upper bound in \(Q_{\pi}^E\) and \(\text{lh}(\bigcup \{\nu_\eta(\alpha) : \alpha < \xi\}) = \xi\).
Definition 4.4. Let us recall the following definition.

The arguments of 4.1 can be repeated here with almost no changes (a slight modification is needed for the justification that $p$ is a $\ast$-parameter on $\lambda$).

Corollary 4.5. We may use the forcing $\bar{P}^E$ to substantially improve [RS05, Corollary 5.1]. First, let us recall the following definition.

Definition 4.2. Let $\bar{E}, E$ be as in 1.11. We define a forcing notion $\bar{P}^E$ as follows.

A condition $p$ in $\bar{P}^E$ is a complete $\lambda$-tree $p \subseteq \lambda^\lambda$ such that

- for every $\nu \in p$, either $|\text{succ}_p(\nu)| = 1$ or $\text{succ}_p(\nu) \in E$, and
- for some set $A \in E$ we have
  \[(\forall \nu \in p)(\text{lh}(\nu) \in A \Rightarrow \text{succ}_p(\nu) \in E).\]

The order $\leq \leq^E$ is the reverse inclusion: $p \leq q$ if and only if $(p, q \in \bar{P}^E$ and $q \leq p$).

Proposition 4.3. Assume that $\bar{E}, E$ are as in 1.11 and $p = (\bar{P}, S, D)$ is a $\Delta^\ell$-parameter on $\lambda$ such that $\lambda \setminus S \in E$. Then $\bar{P}^E$ is a ($<\lambda$)-complete forcing notion of size $2^\lambda$ which is $B$-noble over $p$.

Proof. The arguments of 4.1 can be repeated here with almost no changes (a slight modification is needed for the justification that $p^* \in \bar{P}^E$).

We may use the forcing $\bar{P}^E$ to substantially improve [RS05, Corollary 5.1]. First, let us recall the following definition.

Definition 4.4. Let $\mathcal{F}$ be a filter on $\lambda$ including all co-bounded subsets of $\lambda$, $\emptyset \notin \mathcal{F}$.

1. We say that a family $F \subseteq \lambda^\lambda$ is $\mathcal{F}$-dominating whenever
  \[(\forall g \in \lambda^\lambda)(\exists f \in F)(\{\alpha < \lambda : g(\alpha) < f(\alpha)\} \in \mathcal{F}).\]

2. We say that a family $F \subseteq \lambda^\lambda$ is $\mathcal{F}$-unbounded whenever
  \[(\forall g \in \lambda^\lambda)(\exists f \in F)(\{\alpha < \lambda : g(\alpha) < f(\alpha)\} \in (\mathcal{F})^+).\]

3. If $\mathcal{F}$ is the filter of co-bounded subsets of $\lambda$, then the corresponding dominating/unbounded numbers are also denoted by $\mathfrak{d}_\lambda, \mathfrak{b}_\lambda$. If $\mathcal{F}$ is the filter generated by club subsets of $\lambda$, then the corresponding numbers are called $\mathfrak{d}_{\text{cl}}, \mathfrak{b}_{\text{cl}}$.

Corollary 4.5. Assume $\lambda = \lambda^{<\lambda}$, $2^\lambda = \lambda^+$. Suppose that $p = (\bar{P}, S, D)$ is a $\Delta^\ell$-parameter on $\lambda$, and $E$ is a normal filter on $\lambda$ such that $\lambda \setminus S \in E$. Then there is a $\lambda^+\text{-cc} \lambda$-proper forcing notion $\mathbb{P}$ such that

$$\vdash\neg 2^\lambda = \lambda^{++} = b_{\mathcal{F}^E} = d_{\mathcal{F}^E} = d_\lambda \& b_\lambda = b_{\mathcal{D}[\bar{P}]} = d_{\mathcal{D}[\bar{P}]} = \lambda^+ \text{ "}. $$
Proof. For \( \nu < \lambda \) let \( E_\nu \) be the filter generated by clubs of \( \lambda \) and let \( \bar{E} = \langle E_\nu : \nu < \lambda \rangle \). Let \( Q = \langle P_\xi, Q_\xi : \xi < \lambda^+ \rangle \) be a \( \lambda \)-support iteration such that for every \( \xi < \lambda^+ \), \( \Vdash_{\bar{P}} \langle Q_\xi = \bar{P}_E \rangle \). (Remember, we use the convention that in \( \Vdash \) the normal filter generated by \( E \) is also denoted by \( E \) etc.) Let \( P = P_{\lambda^+} = \lim(Q) \).

It follows from 3.3(a)+4.3 that \( P \) is \( \lambda \)-proper. Using [RS, Theorem 2.2] (see also Eisworth [Eis03, §3]) we see that \( P \) satisfies the \( \lambda^+ \)-cc, \( \Vdash_{\bar{P}} 2^\lambda = \lambda^+ \) and \( P \) is \( (\lambda^+) \)-complete. Thus, the forcing with \( P \) does not collapse cardinals and it also follows from 3.3(b) that

\[ \Vdash_{\bar{P}} ^{\lambda \cap V} \lambda \cap V \text{ is } D[\bar{P}] \text{-dominating in } \lambda^+ \nabla. \]

It is also easy to check, that for each \( \xi < \lambda^+ \)

\[ \Vdash_{\bar{P}} ^{\lambda \cap V\bar{P}_\xi} \text{ is not } E^-\text{unbounded in } \lambda^+ \nabla \]

and hence we may easily conclude that \( \Vdash_{\bar{P}} 2^\lambda = \lambda^+ \nabla. \)

Definition 4.6. Assume that

- \( \lambda \) is weakly inaccessible, \( \lambda^+ = \lambda \),
- \( H : \lambda \rightarrow \lambda \) is such that \( |\alpha|^+ \leq |H(\alpha)| \) for each \( \alpha < \lambda \),
- \( F \) is a normal filter on \( \lambda \), \( \bar{F} = \langle F_\nu : \nu \in \bigcup_{\alpha < \lambda} H(\xi) \rangle \) where \( F_\nu \) is a \( (|\alpha|^+) \)-complete filter on \( H(\alpha) \) whenever \( \nu \in \bigcap_{\xi < \alpha} H(\xi), \alpha < \lambda \).

We define forcing notions \( \mathbb{Q}_{\bar{P}, F}^H \) and \( \mathbb{P}_{\bar{P}, F}^H \) as follows.

1. **A condition** \( p \) in \( \mathbb{Q}_{\bar{P}, F}^H \) is a complete \( \lambda \)-tree \( p \subseteq \bigcup_{\alpha < \lambda} \prod_{\xi < \alpha} H(\xi) \) such that
   - (a) for every \( \nu \in p \), either \( |\text{succ}_p(\nu)| = 1 \) or \( \text{succ}_p(\nu) \subseteq F_\nu \), and
   - (b) for every \( \eta \in \text{lim}_{\lambda}(p) \) the set \( \{ \alpha < \lambda : \text{succ}_p(\eta|\alpha) \subseteq F_{\eta|\alpha} \} \) belongs to \( \bar{F} \).
   
   **The order** of \( \mathbb{Q}_{\bar{P}, F}^H \) is the reverse inclusion.

2. **A condition** \( p \) in \( \mathbb{P}_{\bar{P}, F}^H \) is a complete \( \lambda \)-tree \( p \subseteq \bigcup_{\alpha < \lambda} \prod_{\xi < \alpha} H(\xi) \) satisfying
   - (a) above and
   - (b) \( \forall \nu \in p \left( \forall \eta \in A \in F \right) \text{succ}_p(\nu) \subseteq F_\nu \).

   **The order** of \( \mathbb{P}_{\bar{P}, F}^H \) is the reverse inclusion.

Proposition 4.7. Assume that \( \lambda, H, \bar{F}, F \) are as in 4.6. Let \( \mathbf{p} = (\bar{P}, S, D) \) be a \( \mathcal{D} \ell \)-parameter such that \( \lambda \setminus S \in F \). Then both \( \mathbb{Q}_{\bar{P}, F}^H \) and \( \mathbb{P}_{\bar{P}, F}^H \) are strategically \( (\lambda^+) \)-complete forcing notions of size \( 2^\lambda \) which are also \( B \)-noble over \( \mathbf{p} \).

Proof. Like 1.12(2), 4.1, 4.3.

The property of being \( B \)-noble seems to be a relative of properness for \( D \)-semi diamonds introduced in [RS01] and even more so of properness over \( D \)-diamonds studied in Eisworth [Eis03]. However, technical differences make it difficult to see what are possible dependencies between these notions (see Problem 7.2). In this context, let us note that there are forcing notions which are proper over semi diamonds, but are not \( B \)-noble over any \( \mathcal{D} \ell \)-parameter \( \mathbf{p} \). Let us consider, for example, a forcing notion \( \mathbb{P}^* \) defined as follows:

**a condition** in \( \mathbb{P}^* \) is a function \( p \) such that
Now suppose that $p \Vdash \text{Claim 4.9.1.}$. (above the condition

Proposition 4.9.

Proof. Suppose that $W$ is a $\mathcal{P}^*$-name such that

$$q \Vdash \text{"} W_0 \text{ is a function with } \text{dom}(W_0) \subseteq \lambda \text{ and } \text{rng}(W_0) \subseteq \lambda \text{".}$$

Let $W'$ be a $\mathcal{P}^*$-name for a member of $^\lambda \lambda$ such that

$$q \Vdash \text{"} W_0 \subseteq W' \text{ and } (\forall \alpha \in \lambda \setminus \text{dom}(W_0))(W(\alpha) = \alpha) \text{".}$$

Now suppose that $p = (\bar{P}, S, D)$ is a $\mathcal{D}\ell$-parameter and $\langle A_\alpha : \alpha < \lambda \rangle$ is a sequence of subsets of $\lambda$ such that $|A_\alpha| < \lambda$ for $\alpha < \lambda$. The following claim implies that $\mathcal{P}^*$ (above the condition $q$) is not $B$–noble over $p$ (remember 3.3(b)).

Claim 4.9.1. $\bar{P} \Vdash \text{"} \{\alpha < \lambda : W(\alpha) \in A_\alpha\} \notin \mathcal{D}[\mathcal{P}^*] \text{".}$

Proof of the Claim. Suppose that $p \geq q$ and $B_i, f_i$ (for $i < \lambda$) are $\mathcal{P}^*$–names for members of $D \cap V$ and members of $^\lambda \lambda$, respectively, such that

$$p \Vdash \text{"} \bigcap_{i<\lambda} B_i \cap \bigcap_{i<\lambda} \text{set}^\mathcal{P}(f_i) \subseteq \{\alpha < \lambda : W(\alpha) \in A_\alpha\} \text{".}$$

Build inductively a sequence $\langle p_i, B_i, f_i : i < \lambda \rangle$ such that for each $i < \lambda$:

(i) $p_i \in \mathcal{P}^*$, $p \leq p_0 \leq p_1 \leq p_i$ for $j < i$, 
(ii) $B_i \in D \cap V$, $f_i \in \lambda \lambda$ and 
(iii) $p_i \Vdash \text{"} B_i = B_i$ and $f_j[i] = f_j[i]$ for all $j \leq i$, and 
(iv) $i < \sup(\text{dom}(p_i) \cap \lambda)$.

Since $B = \bigcap_{i<\lambda} B_i \cap \bigcap_{i<\lambda} \text{set}^\mathcal{P}(f_i) \in D$, we may pick a limit ordinal $\delta < B$ such that 

$$\bigcap_{i<\delta} \text{dom}(p_i) \cup \text{rng}(p_i) \cap \lambda \leq \delta$$. 

Put $\alpha = \sup(A_\delta) + 888$ and $p^+ = \bigcup_{\delta} p_i \cup \{\delta, \alpha\}$. Then $p^+ \in \mathcal{P}^*$ is a condition stronger than all $p_i$ for $i < \delta$ and $p^+ \Vdash \text{"} \delta \in \bigcap_{i<\lambda} B_i \cap \bigcap_{i<\lambda} \text{set}^\mathcal{P}(f_i) \text{ and } W(\delta) = \alpha \notin A_\delta \text{"}$. A contradiction. $\square$

A similar construction can be carried out above any condition $q$ such that for some $\alpha \in \text{dom}(q)$ we have $\text{cf}(\alpha) = \lambda$ (the set of such conditions is dense in $\mathcal{P}^*$). $\square$

5. $Q_E^L$ vs $P_E^L$ and Cohen $\lambda$–reals

The forcing notions $Q_E^L$ and $P_E^L$ (introduced in 1.11 and 4.2, respectively) may appear to be almost the same. However, at least under some reasonable assumptions on $E, E$ they do have different properties.

Suppose that $V \subseteq V^*$ are transitive universes of ZFC (with the same ordinals) such that $\langle<^\lambda \lambda \cap V \rangle = \langle<^\lambda \lambda \cap V^* \rangle$. We say that a function $c \in \langle<^\lambda \lambda \cap V \rangle$ is a $\lambda$–Cohen over $V$ if for every open dense set $U \subseteq <^\lambda \lambda$ (where $<^\lambda \lambda$ is equipped with the partial order of the extension of sequences), $U \in V$, there is $\alpha < \lambda$ such that $c|\alpha \in U$.

Proposition 5.1. Assume that
(a) $\lambda$ is a strongly inaccessible cardinal,
(b) $S$ is the set of all strong limit cardinals $\kappa < \lambda$ of countable cofinality,
(c) $E$ is a normal filter on $\lambda$ such that $S \subseteq E$,
(d) $\bar{E} = \langle E_\nu : \nu \in < \lambda \lambda \rangle$ is a system of $(<\lambda)$-complete non-principal filters on $\lambda$.

Then the forcing notion $\mathbb{P}^E_\lambda$ adds a $\lambda$-Cohen over $V$.

Proof. Let $\kappa \in S$. We will say that a tree $T \subseteq <^\kappa \kappa$ is $\kappa$-interesting if

\[(\forall n < \omega) (\forall \nu \in T) (\text{lhd}(\nu) \leq \delta_n \Rightarrow \sup \{\text{rng}(\nu)\} < \delta_{n+1}).\]

Note that there are only $2^\kappa$ many $\kappa$-interesting trees (for $\kappa \in S$). Therefore we may fix a well ordering of the family of $\kappa$-interesting trees of length $2^\kappa$ and choose by induction a function $f_\kappa: <^\kappa \kappa \rightarrow <^\kappa 2 \setminus \{\emptyset\}$ such that

\[(\forall 1) \text{ if } T \subseteq <^\kappa \kappa \text{ is a } \kappa \text{-interesting tree, then}\]

\[(\forall \sigma \in <^\kappa 2 \setminus \{\emptyset\}) (\exists \eta \in \operatorname{lim}_\kappa(T))(f_\kappa(\eta) = \sigma).\]

Let $W$ be a $\mathbb{P}^E_\lambda$-name such that $\models_{\mathbb{P}^E_\lambda} W = \bigcup \{\text{root}(p) : p \in \Gamma_{\mathbb{P}^E_\lambda}\}$. Plainly, $\models W \in \lambda \lambda \lambda$. Next, let $C$ be a $\mathbb{P}^E_\lambda$-name such that $\models_{\mathbb{P}^E_\lambda} C = \{\kappa \in S : W'_\kappa \cap <^\kappa \kappa \}$ and let $\tau$ be a $\mathbb{P}^E_\lambda$-name such that

$\models_{\mathbb{P}^E_\lambda} " \tau \text{ is the concatenation of all elements of the sequence }$

$\langle f_\kappa(W'_\kappa) : \kappa \in C \rangle, \text{ i.e., } \tau = (\ldots \concat f_\kappa(W'_\kappa) \concat \ldots)_{\kappa \in C} "$.

Plainly, $\models \tau \in \lambda \lambda \lambda$ (remember, $\emptyset \notin \text{rng}(f_\kappa)$ for $\kappa \in S$).

We are going to argue that

$\models_{\mathbb{P}^E_\lambda} " \tau \text{ is } \lambda \text{-Cohen over } V "$.

To this end suppose that $U \subseteq <^\lambda 2$ is an open dense set and $p \in \mathbb{P}^E_\lambda$. Let

$B = \{\kappa \in S : (\forall \eta \in p \cap <^\kappa \kappa) (\text{succ}_p(\eta) \in E_\eta)\}$

(so $B \in E$). By induction on $n < \omega$ choose $\delta_n, T_n$ so that

\[(\forall 2) \delta_n \in B, \delta_n < \delta_{n+1}, T_n \subseteq <^\delta_n \lambda \text{ is a complete tree (thus every chain in } T_n \text{ has a } \leq^\delta \text{-bound in } T_n),\]

\[(\forall 3) T_n \subseteq T_{n+1} \subseteq p, T_{n+1} \cap <^\delta_n \lambda = T_n \cap <^\delta_n \lambda, T_0 \subseteq <^\delta_0 \lambda, |T_0 \cap <^\delta_0 \lambda| = 1,\]

\[(\forall 4) \text{ if } \nu \in T_n \cap <^\delta_n \lambda, \text{ then } \text{succ}_{T_n}(\nu) \neq \emptyset \text{ and}\]

\[2 \leq |\text{succ}_{T_n}(\nu)| \Rightarrow \text{lh}(\nu) \in \{\delta_0, \ldots, \delta_{n-1}\},\]

\[(\forall 5) \text{ if } \nu \in T_{n+1} \cap <^\delta_n \lambda, \text{ then } |\text{succ}_{T_{n+1}}(\nu)| = \delta_n \text{ and}\]

\[\sup (\text{rng}(\nu)) < \delta_{n+1} < \min (\text{succ}_{T_{n+1}}(\nu)).\]

Next put $\kappa = \sup(\delta_n : n < \omega)$ and $T = \bigcup T_n$. Clearly $\kappa \in S$ and $T \subseteq <^\kappa \kappa$ is a $\kappa$-interesting tree such that

\[(\forall n \in \operatorname{lim}_\kappa(T))(\kappa = \min(\delta \in S : \delta_0 < \delta \& \eta_0 \delta \in <^\delta \delta)).\]

Let $T_0 \cap <^\delta_0 \lambda = \{\eta_0\}$ and $C(\eta_0) = \{\delta \in S \cap (\delta_0 + 1) : \eta_0 \delta \in <^\delta \delta\}$, and let $\tau_0$ be the concatenation of all elements of the sequence $\langle f_\kappa(\eta_0) \delta : \delta \in C(\eta_0) \rangle$, i.e.,

$\tau_0 = (\ldots \concat f_\kappa(\eta_0) \delta \concat \ldots)_{\delta \in C(\eta_0)} \in <^\lambda 2$. Pick $\sigma \in <^\lambda 2$ such that $\tau_0 <_\sigma \in U$. It
follows from \((\circ)_1\) that we may find \(\eta \in \text{lim}_e(T)\) such that \(\sigma = \tau_0 \cdot f_e(\eta)\). Now note that \((p)_\eta \models p_E \sigma < \tau\).

**Proposition 5.2.** Assume that

(a) \(\lambda\) is a measurable cardinal,
(b) \(E = \langle E_\nu : \nu \in <^{< \lambda}\lambda \rangle\) is a system of normal ultrafilters on \(\lambda\),
(c) \(E\) is a normal filter on \(\lambda\),
(d) \(p \in Q_E^E\) and \(\tau\) is a \(Q_E^E\)-name such that \(p \models \tau \in \lambda 2\),
(e) \(\delta = \langle \delta_\alpha : \alpha < \lambda \rangle\) is an increasing continuous sequence of non-successor ordinals below \(\lambda\) such that \(\delta_0 = 0\) and \(2^{2^{\delta_0}} < |\delta_{\alpha + 1}|\) for all \(\alpha < \lambda\).

Then there is a condition \(q \in Q_E^E\) and a sequence \(\langle A_\alpha : \alpha < \lambda \rangle\) such that

(i) \(q \geq p\) and \(A_\alpha \subseteq [\delta_\alpha, \delta_{\alpha + 1})^2\), \(|A_\alpha| < |\delta_{\alpha + 1}|\) for \(\alpha < \lambda\), and
(ii) \(q \vDash_{Q_E^E} (\forall \alpha < \lambda) (\tau \in [\delta_\alpha, \delta_{\alpha + 1}) \in A_\alpha)\).

In particular, the forcing notion \(Q_E^E\) does not add any \(\lambda\)-Cohen over \(V\).

**Proof.** Let \(\langle \nu_\alpha : \alpha < \lambda \rangle\) be an enumeration of \(<^{< \lambda}\lambda\) such that \(\nu_\alpha < \nu_\beta\) implies \(\alpha < \beta\).

By induction on \(\alpha < \lambda\) we will construct a sequence \(\langle A_\alpha, p_\alpha, X_\alpha : \alpha < \lambda \rangle\) so that for each \(\alpha < \lambda\) we have:

\((\mathbb{2})_1\) \(A_\alpha \subseteq [\delta_\alpha, \delta_{\alpha + 1})^2\), \(|A_\alpha| < |\delta_{\alpha + 1}|\), \(p_\alpha \in Q_E^E\), \(X_\alpha \subseteq p_\alpha\), \(|X_\alpha| \leq \delta_\alpha\),
\((\mathbb{2})_2\) if \(\alpha < \beta < \lambda\), then \(p_\alpha \leq p_\beta\) and \(X_\alpha \subseteq X_\beta\),
\((\mathbb{2})_3\) \(p_0 = p\) and \(X_0 = \{\text{root}(p_0)\}\), and if \(\alpha\) is limit then \(p_\alpha = \bigcap_{\beta < \alpha} p_\beta\) and \(X_\alpha = \bigcup_{\beta < \alpha} X_\beta\),
\((\mathbb{2})_4\) if \(\nu \in X_{\alpha + 1}\), then \(\text{succ}_{p_\alpha}(\nu) \in E_{\nu}\),
\((\mathbb{2})_5\) if \(\alpha\) is limit, \(\nu \in X_\alpha\) and \(\alpha \in \bigcap_{\beta < \alpha} \text{succ}_{p_{\beta}}(\nu)\), then for some \(\eta \in X_{\alpha + 1}\) we have \(\nu^{-1}(\alpha) \subseteq \eta\),
\((\mathbb{2})_6\) if \(\nu_\alpha \in p_{\alpha}\), then there is \(\eta \in X_{\alpha + 1}\) such that \(\nu_\alpha \subseteq \eta\) and if (additionally) \(\nu \subseteq p_\alpha\), then \(\eta = \nu_\alpha\),
\((\mathbb{2})_7\) if \(\alpha\) is limit and \(\nu \in \alpha \cap p_\alpha\) and \(\text{succ}_{p_\alpha}(\nu) \in E_{\nu}\), then \(\nu \in X_{\alpha + 1}\),
\((\mathbb{2})_8\) \(p_{\alpha + 1} \vDash_{Q_E^E} (\forall \tau) (\tau \in [\delta_\alpha, \delta_{\alpha + 1}) \in A_\alpha)\).

Suppose that we have determined \(p_\beta, X_\beta\) for \(\beta < \alpha\) and \(A_\beta\) for \(\beta + 1 < \alpha\) so that the relevant instances of \((\mathbb{2})_1 - (\mathbb{2})_8\) are satisfied. If \(\alpha\) is limit or \(0\), then \(p_\alpha, X_\alpha\) are defined by \((\mathbb{2})_3\) (and \(A_\alpha\) will be chosen at the next step). One easily verifies that \(p_\alpha, X_\alpha\) satisfy the requirements in \((\mathbb{2})_1 - (\mathbb{2})_4\).

So suppose now that \(\alpha = \gamma + 1\) (and we have defined \(p_\gamma, X_\gamma\) and \(A_\beta\) for \(\beta < \gamma\)).

We may easily choose a set \(X_\alpha \subseteq p_\gamma\) such that

\((\mathbb{2})_9\) \(X_\gamma \subseteq X_\alpha\), \(|X_\alpha| < \delta_\alpha\) and \(X_\alpha\) satisfies \((\mathbb{2})_4 - (\mathbb{2})_7\) (with \(\alpha\) there corresponding to \(\gamma\) here),
\((\mathbb{2})_{10}\) if \(\eta_0, \eta_1 \in X_\alpha\), \(\nu = \eta_0 \cap \eta_1\), then \(\nu \in X_\alpha\), and
\((\mathbb{2})_{11}\) if \(\eta_0 : \xi < \zeta \subseteq X_\alpha\) is \(<\cdot\cdot\cdot\)-increasing, then there is \(\eta \in X_\alpha\) such that \((\forall \xi < \zeta)(\eta_\xi \subseteq \eta)\).

Next, for each \(\eta \in X_\alpha\) choose a function \(\sigma_\eta : [\delta_\gamma, \delta_{\gamma + 1}) \rightarrow 2\) and a condition \(q_\eta \in Q_E^E\) so that

\((\mathbb{2})_{12}\) \(\text{root}(q_\eta) = \eta\), \((p_\gamma)_\eta \leq q_\eta\), \((\forall \nu \in X_\alpha)(\eta \nless \nu \Rightarrow \nu(\text{lh}(\eta)) \notin \text{succ}_{q_\eta}(\eta))\) and \(q_\eta \vDash_{Q_E^E} (\forall \tau) (\tau \in [\delta_\gamma, \delta_{\gamma + 1}) \in \sigma_\eta)\).
(Possible by assumption (b) and 2.4.) Put $p_\alpha = \bigcup_{\eta \in X_\alpha} q_\eta$ and $A_\gamma = \{\sigma_\eta : \eta \in X_\alpha\}$. Plainly, $|A_\gamma| \leq |X_\alpha| < \delta_{\gamma+1}$ and $p_\alpha \in Q_E^E$ (to verify that $p_\alpha$ is a complete $\lambda$-tree use $(\mathfrak{S})_{10} + (\mathfrak{S})_{11}$; the other requirements easily follow from the fact that $|X_\alpha| < \lambda$). One also easily checks that $\Vdash_{E} Q_\eta \tau[\delta_\gamma, \delta_{\gamma+1}] \in A_\gamma$.

After the inductive construction is carried out, we put $q = \bigcap_{\alpha < \lambda} p_\alpha$. It follows from $(\mathfrak{S})_2 + (\mathfrak{S})_5 + (\mathfrak{S})_6$ that $q$ is a complete $\lambda$–tree, $|\Vdash_{E} q(\nu)| = 1$ or $\Vdash_{E} q(\nu) \in E_\nu$ for each $\nu \in q$, and

$$\bigcup_{\alpha < \lambda} X_\alpha = \{\nu \in q : \Vdash_{E} q(\nu) \in E_\nu\}.$$ 

Suppose now that $\eta \in \lim_{\lambda}(q)$. Then for each $\alpha < \lambda$ we have $\eta \in \lim_{\alpha}(p_\alpha)$ and hence $B_\alpha \overset{\text{def}}{=} \{\xi < \lambda : \Vdash_{E} p_\alpha(\eta[\xi]) \in E_{\eta[\xi]}\} \in E$. Let $C = \{\delta < \lambda : \delta \text{ is limit and } \eta[\delta] \in \mathcal{F}\}$ (it is a club of $\lambda$). Since $E$ is a normal filter, $C \cap \triangle B_\alpha \in E$. Suppose $\delta \in C \cap \triangle B_\alpha$. Then by $(\mathfrak{S})_3 + (\mathfrak{S})_7$ we have $\eta[\delta] \in X_{\delta+1}$ and thus $\Vdash_{E} \eta[\delta] \in E_{\eta[\delta]}$. Consequently, $q \in Q_E^E$.

Finally, it follows from $(\mathfrak{S})_8$ that $q \Vdash (\forall \alpha < \lambda) (\tau[\delta_\gamma, \delta_{\gamma+1}] \in A_\gamma)$.

Let us note that forcing notions of the form $Q_E^E$ may add $\lambda$–Cohens if the filters $E_\nu$ are far from being ultrafilters.

**Proposition 5.3.** Assume that

(a) $E$ is a normal filter on $\lambda = \lambda^<\lambda$,
(b) $\check{E} = (E_\nu : \nu \in \check{\lambda})$ is a system of $\check{\lambda}$–complete filters on $\lambda$,
(c) for every $\nu \in \check{\lambda}$ there is a family $\{A_\nu^\alpha : \alpha < \lambda\}$ of pairwise disjoint sets from $(E_\nu)^+$.

Then both the forcing notions $Q_E^E$ and $P_E^E$ add $\lambda$–Cohens over $V$.

**Proof.** We will sketch the argument for $Q_E^E$ only (no changes are needed for the case of $P_E^E$).

For each $\nu \in \check{\lambda}$ choose a function $h_\nu : \lambda \rightarrow \check{\lambda} \setminus \{\emptyset\}$ such that

$$(\forall \sigma \in \check{\lambda} \setminus \{\emptyset\}) (\exists \alpha < \lambda) (\forall \xi \in A_\nu^\alpha) (h_\nu(\xi) = \sigma).$$

Let $W$ be a $Q_E^E$–name such that $\Vdash_{Q_E^E} W = \bigcup \{\text{root}(p) : p \in \Gamma_{Q_E^E}\}$ and let $\tau$ be a $Q_E^E$–name such that

$\Vdash_{Q_E^E} \tau$ is the concatenation of all elements of the sequence

$$(h_{W^\xi}(W(\xi)) : \xi \in \check{\lambda}), \text{ i.e.,}$$

$$\tau = \langle h_{W(0)}(W(0)) \downarrow h_{W(0)}(W(1)) \downarrow \cdots h_{W^\xi}(W(\xi)) \downarrow \cdots \downarrow \xi < \lambda \rangle.$$ 

One easily verifies that $\Vdash_{Q_E^E} \tau \in \lambda^2$ is a $\lambda$–Cohen over $V$.

The result in 5.2 would be specially interesting if we only knew that it is preserved in $\lambda$–support iterations. Unfortunately, at the moment we do not know if this is true (see Problem 7.3(1)). However, we may consider properties stronger than adding $\lambda$–Cohens and then our earlier results give some input.
Definition 5.4. Suppose that $V \subseteq V^*$ are transitive universes of ZFC (with the same ordinals) such that $\langle \lambda, V \rangle = \langle \lambda, V^* \rangle$. We say that a function $c \in \lambda^2 \cap V^*$ is a

(1) strongly $\lambda$–Cohen over $V$ if it is a $\lambda$–Cohen (i.e., for every open dense set $U \subseteq \lambda^2$ from $V$ there is $\alpha < \lambda$ such that $c[\alpha] \in U$) and

(2) strongly $\lambda$–Cohen over $V$ if it is a $\lambda$–Cohen and

(3) More generally, if $D$ is a normal filter on $\lambda$, $D \in V$ then we say that $c \in \lambda^2 \cap V^*$ is $D$–strongly $\lambda$–Cohen over $V$ if in (2) we replace “stationary” by “$\in (D^V)^+\cap$” (where $D^V$ is the normal filter generated by $D$ in $V^*$). Similarly for strongly $\lambda$–Cohen.

Remark 5.5. (1) To explain our motivation for 5.4, let us recall that if $c \in \lambda^2$ is $\lambda$–Cohen over $V$ and $\langle \eta_\alpha, \beta_\alpha : \alpha < \lambda \rangle \in V$ is such that $\alpha < \beta_\alpha < \lambda$ and $\eta_\alpha \in [\alpha, \beta_\alpha)^2$ for $\alpha < \lambda$, then

$V^* \models \{ \alpha < \lambda : \eta_\alpha \not\subseteq c \}$ is stationary;

(2) Let $\langle \eta_\alpha, \beta_\alpha : \alpha < \lambda \rangle \in V$ be such that $\alpha < \beta_\alpha < \lambda$ and $\eta_\alpha \in [\alpha, \beta_\alpha)^2$ for $\alpha < \lambda$. Let $C = (\lambda^2, <)$ (so this is the $\lambda$–Cohen forcing notion) and let $c$ be the canonical $C$–name for the generic $\lambda$–real (i.e., $\Vdash_C c = \bigcup \Gamma_c$). Let $Q$ be a $C$–name for a forcing notion in which conditions are closed bounded sets $d \subseteq \lambda$ such that $\langle \forall \alpha \in d \rangle (\eta_\alpha \not\subseteq c)$ ordered by end extension. Then $C \ast Q$ is essentially the $\lambda$–Cohen forcing and

$\Vdash_{C \ast Q} \exists c \not\subseteq \lambda$–Cohen over $V$.

Hence, if we add a $\lambda$–Cohen then we also add a non-strong $\lambda$–Cohen.

(3) Note that strongly $\lambda$–Cohen implies strongly $\lambda$–Cohen. (Simply, for a sequence $\langle \eta_\alpha, \beta_\alpha : \alpha < \lambda \rangle$ consider $\{ 1 - \eta_\alpha, \beta_\alpha : \alpha < \lambda \}$.)

Proposition 5.6. Assume that $\lambda$ is a strongly inaccessible cardinal and $p = (\mathcal{P}, \lambda, D_\lambda)$ is a $\mathcal{D}_\lambda$–parameter such that $P_\delta = \delta_\delta$ and $D_\lambda$ is the filter generated by club subsets of $\lambda$.

(1) If a forcing notion $Q$ is reasonably $\mathcal{B}$–bounding over $p$, then

$\Vdash_Q \exists \lambda$–Cohen over $V$.

(2) If $\hat{Q} = \langle \mathcal{P}_\alpha, Q_\alpha : \alpha < \gamma \rangle$ is a $\alpha$–support iteration such that for every $\alpha < \lambda$,

$\Vdash_{\mathcal{P}_\alpha} Q_\alpha$ is reasonably $\mathcal{B}$–bounding over $p[\mathcal{P}_\alpha]$,

then

$\Vdash_{\mathcal{P}_\alpha} \exists \lambda$–Cohen over $V$. 
Proof. (1) Note that, in $V^Q, D_\lambda[Q]$ is the filter generated by clubs of $\lambda$.

Let $p \in Q$ and let $\eta$ be a $Q$–name such that $p \Vdash \eta \in \lambda^2$. Let $st$ be a winning strategy of Generic in the game $\mathcal{D}_p^B + (p, Q)$.

Let us consider a play of $\mathcal{D}_p^B + (p, Q)$ in which Generic follows the instructions of $st$ while Antigeneric plays as follows. In the course of the play, in addition to his innings $q_\alpha^\ast$, Antigeneric constructs a sequence $\langle \kappa_\alpha, \langle \eta_\alpha^t : t \in I_\alpha \rangle : \alpha < \lambda \rangle$ such that if $\langle I_\alpha, \langle p_\alpha^t, q_\alpha^t : t \in I_\alpha \rangle : \alpha < \lambda \rangle$ is the sequence of the innings of the two players then the following two demands are satisfied.

(1) $\kappa_\alpha$ is a cardinal such that $2^{I_\alpha}|+\alpha+\kappa_\alpha < \kappa_\alpha$ and $\eta_\alpha^t \in \kappa_\alpha^2$ (for $t \in I_\alpha$).

(2) $q_\alpha^t \Vdash Q \eta_\alpha^t = \eta_\alpha^t$ for each $t \in I_\alpha$.

Since the play is won by Generic, there is a condition $q \geq p$ such that

$q \Vdash Q \{ \alpha < \lambda : (\exists t \in I_\alpha)(q_\alpha^t \in \Gamma_Q) \}$ contains a club of $\lambda$.

It follows from (1) that for each $\alpha < \lambda$ we may choose $\varepsilon_\alpha^0 < \varepsilon_\alpha^1$ from the interval $(\alpha, \kappa_\alpha)$ such that $(\forall t \in I_\alpha)(\eta_\alpha^t(\varepsilon_\alpha^0) = \eta_\alpha^t(\varepsilon_\alpha^1))$. For each $\alpha < \lambda$ choose $\nu_\alpha : [\alpha, \kappa_\alpha] \rightarrow 2$ so that $\nu_\alpha(\varepsilon_\alpha^0) = 0$ and $\nu_\alpha(\varepsilon_\alpha^1) = 1$. Then

$q \Vdash Q \{ \alpha < \lambda : \nu_\alpha \notin \eta \}$ contains a club of $\lambda$.

(2) Similar, but we have to work with trees of conditions as in the proof of 1.10. □

6. Marrying $B$–bounding with fuzzy proper

In this section we introduce a property of forcing notions which, in a sense, marries the $B$–bounding forcing notions of [RS05, Definition 3.1(5)] with the fuzzy proper forcings introduced in [RS07, §A.3]. This property, defined in the language of games, is based on two games: the servant game $\mathcal{D}_S^\text{servant}$ which is the part coming from the fuzzy properness and the master game $\mathcal{D}_S^\text{master}$ which is related to the reasonable boundedness property. Later in this section we will even formulate a true preservation theorem for a slightly modified game.

In this section we assume the following:

Context 6.1. (1) $\lambda$ is a strongly inaccessible cardinal,
(2) $D$ is a normal filter on $\lambda$,
(3) $S \in D$, $0 \notin S$, all successor ordinals below $\lambda$ belong to $S$, $\lambda \backslash S$ is unbounded.

Definition 6.2. Let $Q$ be a forcing notion.

(1) A $Q$–servant over $S$ is a sequence $\bar{q} = \langle q_\delta^S : \delta \in S \& t \in I_\delta \rangle$ such that $|I_\delta| < \lambda$ (for $\delta \in S$) and $q_\delta^t \in Q$ (for $\delta \in S$, $t \in I_\delta$).

(2) Let $\bar{q}$ be a $Q$–servant over $S$ and $q \in Q$. We define a game $\mathcal{D}_{S,D}^\text{servant}(\bar{q}, q, Q)$ as follows. A play of $\mathcal{D}_{S,D}^\text{servant}(\bar{q}, q, Q)$ lasts at most $\lambda$ steps during which the players, COM and INC, attempt to construct a sequence $\langle r_\alpha, A_\alpha : \alpha < \lambda \rangle$ such that

- if $r_\alpha \in Q$, $q \leq r_\alpha$, $A_\alpha \in D$ and $\alpha < \beta < \lambda \Rightarrow r_\alpha \leq r_\beta$.

The terms $r_\alpha, A_\alpha$ are chosen successively by the two players so that

- if $\alpha \notin S$, then INC picks $r_\alpha$, $A_\alpha$, and
- if $\alpha \in S$, then COM chooses $r_\alpha$, $A_\alpha$.

If at some moment of the play one of the players has no legal move, then INC wins; otherwise, if both players always had legal moves and the sequence $\langle r_\alpha, A_\alpha : \alpha < \lambda \rangle$ has been constructed, then COM wins.
(c) \( (\forall \delta \in S)(\delta \in \bigcap_{\alpha<\delta} \Lambda_{\alpha} \& \delta \text{ is limit } \Rightarrow (\exists t \in I_\delta)(q_\delta^t \leq r_\delta)). \)

(3) If COM has a winning strategy in the game \( \mathcal{O}_{S,D}^{\text{servant}}(q, q, Q) \), then we will say that \( q \) is an \( (S, D) \)-knitting condition for the servant \( q \).

**Definition 6.3.** Let \( Q \) be a strategically \( (<\lambda) \)–complete forcing notion.

(1) For a condition \( p \in Q \) we define a game \( \mathcal{O}_{S,D}^{\text{master}}(q, Q) \) between two players, Generic and Antigeneric. The game is a small modification of \( \mathcal{O}_{S,D}^{\text{rb}}(p, Q) \) (see 1.8) — the main difference is in the winning condition. A play of \( \mathcal{O}_{S,D}^{\text{master}}(p, Q) \) lasts \( \lambda \) steps and during a play a sequence

\[ \{I_\alpha, \langle q_\alpha^t, q_\alpha^w : t \in I_\alpha \rangle : \alpha < \lambda \} \]

is constructed. Suppose that the players have arrived to a stage \( \alpha < \lambda \) of the game. Now,

\[ \langle N \rangle_\alpha \]

first Generic chooses a non-empty set \( I_\alpha \) of cardinality \( < \lambda \) and a system \( \langle p_\alpha^t : t \in I_\alpha \rangle \) of conditions from \( Q \).

\[ \langle \mathcal{Z} \rangle_\alpha \]

then Antigeneric answers by picking a system \( \langle q_\alpha^t : t \in I_\alpha \rangle \) of conditions from \( Q \) such that \( (\forall t \in I_\alpha)(p_\alpha^t \leq q_\alpha^t) \). At the end, Generic wins the play \( \{I_\alpha, \langle q_\alpha^t, q_\alpha^w : t \in I_\alpha \rangle : \alpha < \lambda \} \) of \( \mathcal{O}_{S,D}^{\text{master}}(p, Q) \) if and only if letting \( q = \langle q_\alpha^t : \alpha \in S \& t \in I_\alpha \rangle \) (it is a \( Q \)-servant over \( S \)) we have

\[ (\otimes)_{D,S}^{\text{master}} \]

there exists an \( (S, D) \)–knitting condition \( p \geq p \) for the servant \( q \).

(2) A forcing notion \( Q \) is reasonably merry over \( (S, D) \) if (it is strategically \( (<\lambda) \)–complete and) Generic has a winning strategy in the game \( \mathcal{O}_{S,D}^{\text{master}}(p, Q) \) for any \( p \in Q \).

**Theorem 6.4.** Assume that \( \lambda, S, D \) are as in 6.1. Let \( \hat{Q} = (P_\alpha, Q_\alpha : \alpha < \gamma) \) be a \( \lambda \)-support iteration such that for each \( \alpha < \gamma \):

\[ \models_{P_\alpha} \text{"Q}_\alpha \text{ is reasonably merry over (S, D)"}. \]

Then

(a) \( P_\gamma = \text{lim}(\hat{Q}) \) is \( \lambda \)-proper, and

(b) for every \( \mathcal{P}_\gamma \)-name \( \bar{z} \) for a function from \( \lambda \) to \( V \) and a condition \( p \in \mathcal{P}_\gamma \), there are \( q \geq p \) and \( \langle A_\xi : \xi < \lambda \rangle \) such that \( (\forall \xi < \lambda)(|A_\xi| < \lambda) \) and \( q \models \text{"}\{\xi < \lambda : \bar{z}(\xi) \in A_\xi\} \in (D^\mathcal{P}_\gamma)^+ \text{"}. \)

**Proof.** (a) The proof starts with arguments very much like those in [RS05, Theorems 3.1, 3.2], so we will state only what should be done (without actually describing how the construction can be carried out). The major difference comes later, in arguments that the chosen condition is suitably generic.

Suppose that \( N \prec (H(\chi), \in, <^\gamma) \) is such that

\[ <^\lambda N \subseteq N, \ |N| = \lambda \text{ and } Q, S, D, \ldots \in N. \]

Let \( p \in N \cap \mathcal{P}_\gamma \) and \( \langle \tau_\alpha : \alpha < \lambda \rangle \) list all \( \mathcal{P}_\gamma \)-names for ordinals from \( N \). For each \( \xi \in N \cap \gamma \) fix a \( \mathcal{P}_\xi \)-name \( St_\xi^0 \in N \) for a winning strategy of Complete in \( \mathcal{O}_{S,D}^0(Q_\xi, \emptyset Q_\xi) \) such that it instructs Complete to play \( \emptyset \bar{q}_\xi \) as long as her opponent plays \( \emptyset \bar{q}_\xi \).

Let us pick an increasing continuous sequence \( \langle w_3 : \delta < \lambda \rangle \) of subsets of \( \gamma \) such that \( \bigcup_{\delta<\lambda} w_3 = N \cap \gamma, w_3 = \{0\} \) and \( |w_3| < \lambda. \)

By induction on \( \delta < \lambda \) choose
so that if the following conditions $(*_0)^{-(*_1)}_{11}$ are satisfied (for each $\delta < \lambda$).

1. All objects listed in $([2]_0)$ belong to $N$ and they are known after stage $\delta$ of the construction.

2. $r_\alpha$, $r_\delta \in P_{\gamma}$, $r_\delta(0) = r_\alpha(0) = p(0)$, and for each $\alpha < \delta < \lambda$ we have $(\forall \xi \in w_{\alpha+1})(r_\alpha(\xi) = r_\delta(\xi) = r_\delta(\xi))$ and $p \leq r_\alpha \leq r_\delta \leq r_\delta$.

3. If $\xi \in \text{dom}(r_\alpha) \setminus w_\delta$, then $r_\delta(\xi) \models$ "the sequence $\langle r_\alpha(\xi), r_\alpha(\xi) : \alpha < \delta \rangle$ is a legal partial play of $\cal{D}_\lambda(\xi, \emptyset_{\xi})$ in which Complete follows $st_\xi$ ", and if $\xi \in w_{\delta+1} \setminus w_\delta$, then $st_\xi$ is a $P_\xi$-name for a winning strategy of Generic in $\cal{D}_{S,D}^{\text{master}}(r_\delta(\xi), \xi)$ such that if $p_\xi^\theta = (p_\delta^\theta : t \in I_\alpha)$ is given by that strategy to Generic at stage $\alpha$, then $I_\alpha$ is an ordinal below $\lambda$. Also $st_0$ is a suitable winning strategy of Generic in $\cal{D}_{S,D}^{\text{master}}(p(0), \emptyset)$.

4. $T_\delta = (T_\delta, r_\delta)$ is a standard $(w_\delta, 1)$-$\gamma$-tree, $|T_\delta| < \lambda$.

5. $p_\delta^\theta = (p_\delta^\theta : t \in T_\delta)$ and $q_\delta^\theta = (q_\delta^\theta : t \in T_\delta)$ are standard trees of conditions in $\cal{Q}_\delta$, $p_\delta^\theta \leq q_\delta^\theta$.

6. If $t \in T_\delta$, $r_\delta(t) = \gamma$, then the condition $q_\delta^\theta$ decides the values of all names $\langle r_\alpha : \alpha < \delta \rangle$.

7. For $t \in T_\delta$ we have $(\text{dom}(p) \cup \bigcup \text{dom}(r_\alpha) \cup w_\delta) \cap r_\delta(t) \subseteq \text{dom}(p_\delta^\theta)$ and for each $\xi \in \text{dom}(p_\delta^\theta) \setminus w_\delta$:

$$p_\delta^\theta(\xi) \models P_\xi \text{ "if the set } \{r_\alpha(\xi) : \alpha < \delta \} \cup \{p(\xi)\} \text{ has an upper bound in } \cal{Q}_\xi, \text{ then } p_\delta^\theta(\xi) \text{ is such an upper bound ",}$$

8. If $\xi \in w_\delta$, then $\xi_{\delta,\xi}$ is a $P_\xi$-name for an ordinal below $\lambda$, $\xi_{\delta,\xi}, \xi_{\delta,\xi}$ are $P_\xi$-names for $\xi_{\delta,\xi}$-sequences of conditions in $\cal{Q}_\xi$.

9. If $\xi \in w_{\delta+1} \setminus w_\delta$, then $\models P_\xi \text{ " } (\xi_{\alpha,\xi}, \xi_{\alpha,\xi}, \xi_{\alpha,\xi} : \alpha < \lambda) \text{ is a play of } \cal{D}_{S,D}^{\text{master}}(r_\delta(\xi), \cal{Q}_\xi) \text{ in which Generic uses } st_\xi \text { ",} "$

10. If $t \in T_\delta$, $r_\delta(t) = \xi < \gamma$, then the condition $q_\delta^\theta$ decides the value of $\xi_{\delta,\xi}$, say $p_\delta^\theta \models " \xi_{\delta,\xi} = \xi_{\delta,\xi}^\theta \text{", and } \{\xi : t < s \in T_\delta\} = \xi_{\delta,\xi}^\theta$ and $q_\delta^\theta \models P_\xi \text{ " } (\xi_{\delta,\xi}(\xi) \leq p_{\delta,\xi}(\xi) \text{ and } \xi_{\delta,\xi}(\xi) = q_{\delta,\xi}(\xi) \text{ for } \varepsilon < \varepsilon_{\delta,\xi}^\theta \text{ "}."

11. If $t_0, t_1 \in T_\delta$, $r_\delta(t_0) = r_\delta(t_1)$ and $\xi \in w_\delta \cap r_\delta(t_0), t_0(\xi) = t_1(\xi)$ but $t_0(\xi) \neq t_1(\xi)$, then $p_{t_0, t_1}^\delta \models P_\xi \text{ " the conditions } p_{t_0, t_1}^\delta(\xi), p_{t_1, t_1}^\delta(\xi) \text{ are incompatible ",} "$

12. If $t_0 \in T_\delta$, $r_\delta(t_0) = r_\delta(t)$ and if $t \in T_\delta, \xi \in \text{dom}(r_\delta) \cap r_\delta(t) \setminus w_\delta$, and $q_\delta^\theta(\xi) \leq q \in P_\xi, r_\delta(\xi) \leq q$, then $q \models P_\xi \text{ " if the set } \{r_\alpha(\xi) : \alpha < \delta \} \cup \{q_\delta^\theta(\xi), p(\xi)\} \text{ has an upper bound in } \cal{Q}_\xi, \text{ then } r_\delta(\xi) \text{ is such an upper bound ",} "$

After the construction is carried out we define a condition $r \in P_\gamma$ as follows. We let $\text{dom}(r) = N \cap \gamma$ and for $\xi \in \text{dom}(r)$ we let $r(\xi)$ be a $P_\xi$-name for a condition.
in $Q_\xi$ such that if $\xi \in w_{\alpha+1} \setminus w_\alpha$, $\alpha < \lambda$ (or $\xi = 0 = \alpha$), then

$\|q_\xi\| \geq r_\alpha(\xi)$ is an $(S, D)$–knighting condition for the servant $\overline{q_\xi} \overset{\text{def}}{=} (\overline{q_{\delta, \xi}(\varepsilon)} : \delta \in S \& \varepsilon < \varepsilon_{\delta, \xi})$ ”.

Clearly $r$ is well defined (remember $(*)_8$). Note also that $r_3 \leq r$ for all $\delta < \lambda$ and $p \leq r$. We will argue that $r$ is an $(N, \mathbb{P}_\rho)$–generic condition. To this end suppose towards contradiction that $r^* \geq r$, $\alpha^* < \lambda$ and $r^* \Vdash \tau_\alpha \notin N$.

For each $\xi \in N \cap \gamma$ fix a $\mathbb{P}_\xi$–name $st_\xi^S$ for a winning strategy of COM in the game $\mathcal{O}_{\text{servant}}(\overline{q_\xi}, r(\xi), \xi)$. Moreover, for each $\xi < \gamma$ fix a $\mathbb{P}_\xi$–name $st_\xi^Q$ for a winning strategy of Complete in $\mathcal{O}_\xi^Q(Q_\xi, \emptyset_{Q_\xi})$ such that it instructs Complete to play $\emptyset_{Q_\xi}$ as long as her opponent plays $\emptyset_{Q_\xi}$.

By induction on $\delta < \lambda$ we will build a sequence

$$
(r_\delta^*, r_\delta^+, \langle A_{\alpha,i}^\xi, A_{\alpha,i}^\lambda : i < \lambda \& \xi \in N \cap \gamma \rangle : \delta < \lambda)
$$

such that the following demands $(*)_{12} - (*)_{16}$ are satisfied:

$(*)_{12}$ $r_\delta^* \subseteq \mathbb{P}_\gamma$, $r^* \leq r_\alpha^* \leq r_\alpha^+ \leq r_\delta^*$ for $\alpha < \delta < \lambda$,

$(*)_{13}$ $A_{\alpha,i}^\xi$ is a $\mathbb{P}_\xi$–name for an element of $D \cap V$ for $\xi \in N \cap \gamma$, $i < \lambda$,

$(*)_{14}$ if $\delta \in \lambda \setminus S$ and $\xi \in w_\lambda$, then $r_\delta^*|\xi \Vdash (\forall \alpha < \delta)(\forall i < \delta)(A_{\alpha,i}^\xi = A_{\alpha,i}^\lambda)$,

$(*)_{15}$ if $\beta < \delta < \lambda$ and $\xi \in w_{\beta+1} \setminus w_\beta$, then for some $\mathbb{P}_\xi$–names $\langle g^\xi_\alpha : \alpha \leq \beta \rangle$ we have

$r_\delta^*|\xi \Vdash " the sequence $\langle g^\xi_\alpha, \Delta A_{\alpha,i}^{\xi, \lambda} : i < \lambda \rangle$ is a legal partial play of $\mathcal{O}^S_{\text{servant}}(\overline{q_\xi}, r(\xi), \xi)$ in which Generic follows $st_\xi^S",$

$(*)_{16}$ dom($r_\delta^*$) = dom($r_\delta^+$), $r_\delta^* \restriction w_\delta = r_\delta^+ \restriction w_\delta$ and for each $\xi \in \mathcal{O}_{\text{servant}}(\overline{q_\xi}, r(\xi), \xi)$.

So suppose that we have arrived to a stage $\delta < \lambda$ of the construction and

- $r_\delta^*, r_\delta^+$ for $\alpha < \delta$,
- $A_{\alpha,i}^\xi$ for $i < \lambda$, $\alpha < \delta$ and $\xi \in \bigcup_{\beta<\delta} w_\beta$,
- $A_{\alpha,i}^\lambda$ for $\alpha, i < \sup(\delta \setminus S)$ and $\xi \in \bigcup_{\beta<\sup(\delta \setminus S)} w_\beta$

have been determined.

**Case 1**: $\delta \notin S$.

Note that by our assumption on $S$ (in 6.1), $\delta$ is not a successor ordinal, so $w_\delta = \bigcup w_\alpha$ (or $\delta = 0$ and $w_0 = \{0\}$). By $(*)_{16} + (\ast)_{15}$ we may choose a condition $r_\delta^+$ stronger than all $r_\alpha^+$ (for $\alpha < \delta$) and stronger than $r^*$ and such that for each $\xi \in w_\delta$

- if $\alpha < \delta$ and $i < \delta$, then $r_\delta^*|\xi$ forces a value to $A_{\alpha,i}^\xi$, say $r_\delta^*|\xi \Vdash \overline{q_\xi}_{\delta, \xi}(\varepsilon) = A_{\alpha,i}^\xi$.

For $\xi \in w_\delta$ and $i < \lambda$ we also let $A_{\alpha,i}^\xi$ be a $\mathbb{P}_\xi$–name for the interval $(\delta, \lambda)$. The condition $r_\delta^+$ is fully determined by $(*)_{16}$. 

Case 2: \( \delta \in S \) is a successor ordinal, say \( \delta = \beta + 1 \).

First, for each \( \xi \in w_{\delta} \setminus w_{\beta} \) we pick \( \mathbb{P}_\xi \)-names \( s^\xi_\alpha \) and \( A^\xi_{\alpha,i} \) (for \( \alpha < \delta, i < \lambda \)) such that

\[
r^\delta_{\beta} \models \{ \xi \suchthat r^\delta_{\beta} (\xi) = s^\xi_\alpha \}
\]

so that

\[
r^\delta_{\beta} \models \{ \xi \suchthat r^\delta_{\beta} (\xi) = s^\xi_\alpha \}
\]

that and such that we may find a condition \( r^\delta_{\beta} (\xi) \) and \( A^\delta_{\delta,i} \) (for \( i < \lambda \)) such that

\[
r^\delta_{\beta} \models \{ \xi \suchthat r^\delta_{\beta} (\xi) = s^\xi_\alpha \}
\]

Next, we let \( \text{dom}(r^\delta_{\beta}) = \text{dom}(r^\delta_{\beta}) \) and for each \( \xi \in w_{\beta} \) we choose \( \mathbb{P}_\xi \)-names \( r^\delta_{\beta} (\xi) \) and \( A^\delta_{\delta,i} \) (for \( i < \lambda \)) such that

\[
r^\delta_{\beta} \models \{ \xi \suchthat r^\delta_{\beta} (\xi) = s^\xi_\alpha \}
\]

For \( \xi \in \text{dom}(r^\delta_{\beta}) \setminus w_{\beta} \) we put \( r^\delta_{\beta} (\xi) = r^\delta_{\beta} (\xi) \). Then we define condition \( r^\delta_{\beta} \) by \( (\ast_{15}) \).

Case 3: \( \delta \in S \) is a limit ordinal.

We let \( \text{dom}(r^\delta_{\beta}) = \bigcup_{\alpha < \delta} \text{dom}(r^\delta_{\beta}) \) and by induction on \( \xi \in \text{dom}(r^\delta_{\beta}) \) we define \( r^\delta_{\beta} (\xi) \) so that

\[
\begin{align*}
\text{if } & \xi \notin w_{\beta} , \text{ then } r^\delta_{\beta} (\xi) \models (\forall \alpha < \delta) (r^\delta_{\beta} (\xi) \leq r^\delta_{\beta} (\xi)) \text{ (exists by } (\ast_{16})) , \\
\text{if } & \xi \in w_{\beta} , \text{ then for some } \mathbb{P}_\xi \text{-names } A^\delta_{\delta,i} \text{ for members of } D \cap V
\end{align*}
\]

\[
r^\delta_{\beta} \models \{ \xi \suchthat r^\delta_{\beta} (\xi) = s^\xi_\alpha \}
\]

The condition \( r^\delta_{\beta} \) is given by \( (\ast_{16}) \).

After the above construction is carried out we note that

\[
\{ \delta < \lambda : (\forall \xi \in w_{\beta})(\forall i < \delta)(\delta \in A^\delta_{\delta,i}) \} \in D ,
\]

so we may choose an ordinal \( \delta \in S \setminus (\alpha + 1) \) which is a limit of points from \( \lambda \setminus S \) and such that \( \delta \in \bigcap_{\alpha < \delta} \bigcup_{\xi < \lambda} A^\delta_{\delta,i} \) for all \( \xi \in w_{\beta} \). The following claim provides the desired contradiction (remember \( (\ast_0) + (\ast_5) \)).

Claim 6.4.1. For some \( t \in T \) such that \( \text{rk}_t (\beta) = \gamma \) the conditions \( q^\xi_\alpha \) and \( r^\delta_{\beta} \) are compatible

Proof of the Claim. The proof is very much like that of Claim 2.7.1. Let \( (\varepsilon_\beta : \beta \leq \beta^*) = w_{\beta} \cup \{ \gamma \} \) be the increasing enumeration. For each \( \xi < \gamma \) fix a \( \mathbb{P}_\xi \)-name \( st^\xi_\xi \) for a winning strategy of Complete in \( \mathbb{D}_0^\alpha (\mathcal{Q}_\xi, \mathcal{Y}_\xi) \) such that it instructs Complete to play \( \mathcal{Y}_\xi \) as long as her opponent plays \( \mathcal{Y}_\xi \).

By induction on \( \beta \leq \beta^* \) we will choose conditions \( s_{\beta}, s^*_{\beta} \in \mathbb{P}_\varepsilon_{\beta} \) and \( t = (t)_{\varepsilon_{\beta}} : \beta < \beta^* \in T \) such that letting \( t^\beta_{\varepsilon_{\beta}} = (t)_{\varepsilon_{\beta}} : \beta^* < \beta \in T \), we have

\[
\begin{align*}
(\Box)_a \quad & q^\delta_{\beta} \leq s_{\beta} \text{ and } r^\delta_{\beta} \varepsilon_{\beta} \leq s_{\beta} , \\
(\Box)_b \quad & \text{dom}(s_{\beta}) = \text{dom}(s^*_{\beta}) \text{ and for every } \xi < \varepsilon_{\beta} ,
\end{align*}
\]

\[
s^*_{\beta} \models \{ \xi \suchthat (s^*_{\beta} (\xi)) \text{ is a partial legal play of } \mathbb{D}_0^\alpha (\mathcal{Q}_\xi, \mathcal{Y}_\xi) \text{ in which Complete uses her winning strategy } st^\xi_\xi \}.
\]

Suppose that \( \beta \leq \beta^* \) is a limit ordinal and we have already defined \( t^\beta_{\varepsilon_{\beta}} = (t)_{\varepsilon_{\beta}} : \beta^* < \beta \) and \( (s^*_{\beta}, s^*_{\beta'}) : \beta' < \beta \). Let \( \xi = \sup(\varepsilon_{\beta'} : \beta' < \beta) \). It follows from \( (\Box)_b \) that we may find a condition \( s_{\beta} \in \mathbb{P}_\varepsilon_{\beta} \) such that \( s_{\beta} (\xi) \) is stronger than all \( s^*_{\beta'} \) (for \( \beta' < \beta \)) and \( s_{\beta} (\xi) < s_{\beta} (\xi) \). Clearly \( r^\delta_{\beta} \varepsilon_{\beta} \leq s_{\beta} \) and also \( q^\delta_{\beta} (\xi) \leq s_{\beta} (\xi) \).
Then the condition $s$ and $t$ (remember ((\ref{ref}))).

Suppose that $\xi < \zeta < \epsilon_\beta$ and we know $q^s_{\zeta}(\xi) \leq s_\beta(\zeta)$. By $(\ast)_2 + (\ast)_4 + (\ast)_6$ we know that $s_\beta(\zeta) \vdash (\forall \theta < \delta)(r_\zeta(\theta) \leq p^s_{\zeta}(\theta)) \leq q^s_{\zeta}((\zeta))$ and therefore we may use $(\ast)_11$ to conclude that

$$s_\beta(\zeta) \vdash q^s_{\zeta}(\zeta) \leq r_\beta(\zeta) \leq r^*_\beta(\zeta) = s_\beta(\zeta).$$

Then the condition $s_\beta \in \mathcal{P}_{\text{d}}$ is determined.

Now suppose that $\beta = \beta' + 1 \leq \beta^*$ and we have already defined $s_{\beta'}, s^*_{\beta'} \in \mathcal{F}_{\epsilon_{\beta'}}$ and $t_{\beta'} \in T_3$. It follows from the choice of $\delta$ and $(\ast)$_14 that $r^*_\beta(\epsilon_{\beta'}) \vdash \delta \in \bigcap_{\alpha < \delta, \lambda < \epsilon_{\beta'}} A^\delta_{\alpha, i}$ and hence, by the choice of $r$ and $(\ast)_15$ we have (remember (\ref{ref}) of 6.2(2))

$$s_{\beta'} \vdash_{\mathcal{P}_{\epsilon_{\beta'}}} \left(\exists \epsilon < \epsilon_{\delta_{\epsilon_{\beta'}}} \right) (t_{\beta'}(\epsilon_{\beta'}) \leq r^*_\beta(\epsilon_{\beta'})).$$

Therefore we may use $(\ast)_9$ to choose $\epsilon = (t)_{\epsilon_{\beta'}}$ and a condition $s_{\beta} \in \mathcal{F}_{\epsilon_{\beta}}$ such that

- $t^\beta \stackrel{\text{def}}{=} t^\beta_{\epsilon_{\beta}} \cup \{(\epsilon_{\beta'}, \epsilon)\} \in T_{\beta}$, $s_{\beta'} \leq s_{\beta}[\epsilon_{\beta'}]$ and $s_{\beta}[\epsilon_{\beta'}] \vdash_{\mathcal{P}_{\epsilon_{\beta'}}} \left(\exists \epsilon < \epsilon_{\delta_{\epsilon_{\beta'}}} \right) (t_{\beta'}(\epsilon_{\beta'}) \leq r^*_\beta(\epsilon_{\beta'}) = s_\beta(\epsilon_{\beta'})).$
- $r^*_\beta(\epsilon_{\beta'}, \epsilon_{\beta}) = s_{\beta}[\epsilon_{\beta'}]$.

We finish exactly like in the limit case.

After the inductive construction is completed, look at $t = t^\beta_{\epsilon_{\beta}}$ and $s_{\beta'}$. \hfill \Box

\begin{definition}
(See [RS05, Def. 3.1]). Let $\mathcal{Q}$ be a strategically ($<\lambda$)-complete forcing notion.

1. Let $p \in \mathcal{Q}$. A game $\mathcal{G}^B(p, \mathcal{Q})$ is defined similarly to $\mathcal{G}^B(p, \mathcal{Q})$ (see 1.8) except that the winning criterion $(\oplus)^B_{\text{p}}$ is weakened to

\begin{itemize}
  \item [(\oplus)^B_{\text{p}}] there is a condition $p^* \in \mathcal{Q}$ stronger than $p$ and such that
  \begin{align*}
  p^* \vdash_{\mathcal{Q}} & \left(\exists \alpha < \lambda : (\exists t \in I_\alpha) \left(q^t \in \Gamma_\alpha\right)\right) \in D^\mathcal{Q}.
  \end{align*}
\end{itemize}

2. A forcing notion $\mathcal{Q}$ is reasonably $B$-bounding over $D$ if for any $p \in \mathcal{Q}$, Generic has a winning strategy in the game $\mathcal{G}^B(p, \mathcal{Q})$.

\end{definition}

\begin{observation}
If $\mathcal{Q}$ is reasonably $B$-bounding over $D$, then it is reasonably merry over $(S, D)$.

It is not clear though, if forcing notions which are reasonably $B$-bounding over a $\mathcal{D}_T$-parameter $p$ are also reasonably merry (see Problem 7.4). Also, we do not know if fuzzy properness introduced in [RS07, §A.3] implies that the considered forcing notion is reasonably merry (see Problem 7.5), even though the former property seems to be almost built into the latter one.

One may ask if being reasonably merry implies being $B$-bounding. There are examples that this is not the case. The forcing notion $\mathcal{Q}^{\mathcal{D}}_{\eta}$ (see 6.8 below) was introduced in [RS05, Section 6] and by [RS05, Proposition 6.4] we know that it is not reasonably $B$-bounding over $D$. However we will see in 6.12 that it is reasonably merry over $(S, D)$.

\begin{definition}
(See [RS05, Def. 5.1]). (1) Let $\alpha < \beta < \lambda$. An $(\alpha, \beta)$-extending function is a mapping $e : \mathcal{P}(\alpha) \rightarrow \mathcal{P}(\beta) \setminus \mathcal{P}(\alpha)$ such that $e(u) \cap \alpha = u$ for all $u \in \mathcal{P}(\alpha)$.

\end{definition}
(2) Let $C$ be an unbounded subset of $\lambda$. A \textit{$C$-extending sequence} is a sequence $c = \langle c_\alpha : \alpha \in C \rangle$ such that each $c_\alpha$ is an $(\alpha, \min(C \setminus (\alpha + 1)))$-extending function.

(3) Let $C \subseteq \lambda$, $|C| = \lambda$, $\beta \in C$, $w \subseteq \beta$ and let $\epsilon = \langle c_\alpha : \alpha \in C \rangle$ be a $C$-extending sequence. We define $pos^+(w, \epsilon, \beta)$ as the family of all subsets $u$ of $\beta$ such that

(i) if $\alpha_0 = \min(\{\alpha \in C : (\forall \xi < w)((\xi < \alpha)\} )$, then $u \cap \alpha_0 = w$ (so if $\alpha_0 = \beta$, then $u = w$), and

(ii) if $\alpha_0, \alpha_1 \in C$, $w \subseteq \alpha_0 < \alpha_1 = \min(C \setminus (\alpha_0 + 1)) \leq \beta$, then either $c_{\alpha_0}(u \cap \alpha_0) = u \cap \alpha_1$ or $u \cap \alpha_0 = u \cap \alpha_1$,

(iii) if $\sup(w) < \alpha_0 = \sup(C \cap \alpha_0) \notin C$, $\alpha_1 = \min \left(C \setminus (\alpha_0 + 1)\right) \leq \beta$, then $u \cap \alpha_1 = u \cap \alpha_0$.

For $\alpha_0 \in \beta \cap C$ such that $w \subseteq \alpha_0$, the family $\text{pos}(w, \epsilon, \alpha_0, \beta)$ consists of all elements $u$ of $\text{pos}^+(w, \epsilon, \beta)$ which satisfy also the following condition:

(iv) if $\alpha_1 = \min \left(C \setminus (\alpha_0 + 1)\right) \leq \beta$, then $u \cap \alpha_1 = c_{\alpha_0}(w)$.

Definition 6.8 (See [RS05, Def. 6.2]). We define a forcing notion $Q^2_D$ as follows.

A \textbf{condition in} $Q^2_D$ is a triple $p = (w^p, C^p, \epsilon^p)$ such that

(i) $C^p \in D$, $w^p \subseteq \text{min}(C^p)$,

(ii) $\epsilon^p = \langle c_\alpha^p : \alpha \in C^p \rangle$ is a $C^p$-extending sequence.

The order $\leq Q^2_D \leq$ of $Q^2_D$ is given by $p \leq Q^2_D q$ if and only if

(a) $C^q \subseteq C^p$ and $w^q \in \text{pos}^+(w^p, C^p, \min(C^q))$ and

(b) if $\alpha_0, \alpha_1 \in C^q$, $\alpha_0 < \alpha_1 = \min(C^q \setminus (\alpha_0 + 1))$ and $u \in \text{pos}^+(w^q, C^q, \alpha_0)$, then $c_{\alpha_0}^p(u) \in \text{pos}(u, C^p, \alpha_0, \alpha_1)$.

For $p \in Q^2_D$, $\alpha \in C^p$ and $u \in \text{pos}^+(w^p, C^p, \alpha)$ we let $p \upharpoonright_{\alpha, u} \overset{\text{def}}{=} (u, C^p \setminus \alpha, C^p \setminus \alpha)$.

In [RS05, Problem 6.1] we asked if $\lambda$-support iterations of forcing notions $Q^2_D$ are $\lambda$-proper. Now we may answer this question positively (assuming that $\lambda$ is strongly inaccessible). First, let us state some auxiliary definitions and facts.

Proposition 6.9.

(1) $Q^2_D$ is a $(\langle \lambda \rangle)$-complete forcing notion of cardinality $2^\lambda$.

(2) If $p \in Q^2_D$ and $\alpha \in C^p$, then

• for each $u \in \text{pos}^+(w^p, C^p, \alpha)$, $p \upharpoonright_{\alpha, u} \in Q^2_D$ is a condition stronger than $p$, and

• the family $\{p \upharpoonright_{\alpha, u} : u \in \text{pos}^+(w^p, C^p, \alpha)\}$ is pre-dense above $p$.

(3) Let $p \in Q^2_D$ and $\alpha < \beta$ be two successive members of $C^p$. Suppose that for each $u \in \text{pos}^+(w^p, C^p, \alpha)$ we are given a condition $q_u \in Q^2_D$ such that $p \upharpoonright_{\beta, C^p} \overset{\text{def}}{=} (u, C^p \setminus \alpha, \alpha')$ is a condition stronger than $p$, and $\alpha'$ is minimal in $C^p \setminus \beta$ we have

(a) $p \leq q$, $w^q = w^p$, $C^q \cap \beta = C^p \cap \beta$ and $c_\delta^q = c_\delta^p$ for $\delta \in C^q \setminus \alpha$, and

(b) $\bigcup \{w^\delta : u \in \text{pos}^+(w^p, C^p, \alpha)\} \subseteq \alpha'$, and

(c) $q_u \leq p \upharpoonright_{\alpha, C^p} \overset{\text{def}}{=} (u, C^q \setminus \alpha, \alpha)$ for every $u \in \text{pos}^+(w^p, C^p, \alpha)$.

(4) Assume that $p \in Q^2_D$, $\alpha \in C^p$ and $\tau$ is a $Q^2_D$-name such that $p \models \langle \tau \in V \rangle$. Then there is a condition $q \in Q^2_D$ stronger than $p$ and such that

(a) $w^q = w^p$, $\alpha \in C^q$ and $C^q \cap \alpha = C^p \cap \alpha$, and

(b) if $u \in \text{pos}^+(w^q, C^q, \alpha)$ and $\gamma = \min(C^q \setminus (\alpha + 1))$, then the condition $q \upharpoonright_{\alpha, C^q} \overset{\text{def}}{=} (u, C^q \setminus \alpha, \gamma)$ forces a value to $\tau$. 
Proof. Fully parallel to [RS05, Proposition 5.1].

Definition 6.10. The natural limit of an $\leq_{\mathbb{Q}_D^2}$-increasing sequence $\bar{p} = \langle p_\xi : \xi < \gamma \rangle \subseteq \mathbb{Q}_D^2$ (where $\gamma < \lambda$ is a limit ordinal) is the condition $q = (w^q, C^q, c^q)$ defined as follows:

- $w^q = \bigcup \limits_{\xi < \gamma} w^q_\xi$, $C^q = \bigcap \limits_{\xi < \gamma} C^q_\xi$ and $c^q = \langle c^q_\xi : \delta \in C^q \rangle$ is such that for $\delta \in C^q$ and $u \subseteq \delta$ we have $c^q_\delta(u) = \bigcup \limits_{\xi < \gamma} c^q_\xi(u)$.

Proposition 6.11. (1) Suppose $\bar{p} = \langle p_\xi : \xi < \lambda \rangle$ is an $\leq_{\mathbb{Q}_D^2}$-increasing sequence of conditions from $\mathbb{Q}_D^2$ such that

(a) $w^p_\xi = w^{p_\xi}$ for all $\xi < \lambda$, and

(b) if $\gamma < \lambda$ is limit, then $p_\gamma$ is the natural limit of $\bar{p}|_{\gamma}$, and

(c) for each $\xi < \lambda$, if $\delta \in C^p_\xi$, $\text{otp}(C^p_\xi \cap \delta) = \xi$, then $C^p_{\xi+1} \cap (\delta + 1) = C^p_\xi \cap (\delta + 1)$ and for every $\alpha \in C^p_{\xi+1} \cap \delta$ we have $c^p_{\xi+1} = c^p_\xi$.

Then the sequence $\bar{p}$ has an upper bound in $\mathbb{Q}_D^2$.

(2) Suppose that $p \in \mathbb{Q}_D^2$ and $h$ is a $\mathbb{Q}_D^2$-name such that $p \models \langle h : \lambda \rightarrow V^\prime \rangle$. Then there is a condition $q \in \mathbb{Q}_D^2$ stronger than $p$ and such that

- if $\delta < \delta'$ are two successive points of $C^q$, $u \in \text{pos}(w^q, c^q, \delta)$, then the condition $q|_{\delta < \delta'} c^q_{\delta'}(u)$ decides the value of $b|_{\delta + 1}$.

Proof. Fully parallel to [RS05, Proposition 5.2].

Proposition 6.12. Assume that $\lambda, S, D$ are as in 6.1. The forcing notion $\mathbb{Q}_D^2$ is reasonably merry over $(S, D)$.

Proof. Let $p \in \mathbb{Q}_D^2$. We will describe a strategy $\textbf{st}$ for Generic in the game $\mathcal{G}^\text{master}_{S, D}(p, \mathbb{Q}_D^2)$ - this strategy is essentially the same as the one in the proof of [RS05, Proposition 5.4], only the argument that it is a winning strategy is different.

In the course of a play the strategy $\textbf{st}$ instructs Generic to build aside an increasing sequence of conditions $\bar{p}' = \langle p'_\alpha : \alpha < \lambda \rangle \subseteq \mathbb{Q}_D^2$ such that for each $\alpha < \lambda$:

- (a) $p'_5 = p$ and $w^{p'_\alpha} = w^p$, and

- (b) if $\alpha < \lambda$ is limit, then $p'_\alpha$ is the natural limit of $\bar{p}'|_{\alpha}$, and

- (c) if $\delta \in C^p_\alpha$, $\text{otp}(C^p_\alpha \cap \delta) = \alpha$, then $C^{p'_\alpha+1} \cap (\delta + 1) = C^p_\alpha \cap (\delta + 1)$ and for every $\xi \in C^{p'_\alpha+1} \cap \delta$ we have $c^{p'_\alpha+1}_\delta = c^{p'_\alpha}_\delta$, and

- (d) after stage $\alpha$ of the play of $\mathcal{G}^\text{master}_{S, D}(p, \mathbb{Q}_D^2)$, the condition $p'_\alpha$ is determined.

After arriving to the stage $\alpha$, Generic is instructed to pick $\delta \in C^{p'_\alpha}$ such that $\text{otp}(C^{p'_\alpha} \cap \delta) = \alpha$, put $\gamma = \text{min}(C^{p'_\alpha} \setminus (\delta + 1))$ and play as her innings of this stage:

$$I_\alpha = \text{pos}^+(w^{p'_\alpha}, c^{p'_\alpha}, \delta) \quad \text{and} \quad p^*_\alpha = p'_\alpha \upharpoonright \gamma c^{p'_\alpha}(u) \quad \text{for} \quad u \in I_\alpha.$$ 

Then Antigeneric answers with $\langle q^*_u : u \in I_\alpha \rangle \subseteq \mathbb{Q}_D^2$. Since $p^*_\alpha \downharpoonright \gamma c^{p'_\alpha}(u) \leq q^*_u$ for each $u \in \text{pos}^+(w^{p'_\alpha}, c^{p'_\alpha}, \delta)$, Generic may use 6.9(3) (with $\delta, \gamma, p^*_\alpha, q^*_u$ here standing for $\alpha, \beta, p, q_\alpha$ there) to pick a condition $p^*_\alpha+1$ such that, letting $\alpha' = \text{min}(C^{p^*_\alpha+1} \setminus \gamma)$, we have

- (e) $p^*_\alpha \leq p^*_\alpha+1$, $w^{p^*_\alpha+1} = w^p$, $C^{p^*_\alpha+1} \cap \gamma = C^p_\alpha \cap \gamma$ and $c^{p^*_\alpha+1}_\xi = c^{p^*_\alpha}_\xi$ for $\xi \in C^{p^*_\alpha+1} \cap \delta$, and

- (f) $\bigcup \{ w^{q^*_u} : u \in I_\alpha \} \subseteq \alpha'$, and
(g) \( q^\alpha_\delta \leq p^*_{\alpha+1} |_{\alpha, \gamma} C^\delta_{\alpha+1}(u) \) for every \( u \in I_\alpha \).

This completes the description of \( st \). Suppose that \( \langle I_\alpha, (p^\alpha_\delta, q^\alpha_\delta : u \in I_\alpha) : \alpha < \lambda \rangle \) is the result of a play of \( \mathcal{D}_{S,D}^{\text{master}}(p, Q_{\Delta}^D) \) in which Generic followed \( st \) and constructed aside \( \bar{p}^* = (p^\alpha_\delta : \alpha < \lambda) \). By 6.11, there is a condition \( p^* \in Q_{\Delta}^D \) stronger than all \( p^\alpha_\delta \) (for \( \alpha < \lambda \)). We claim that \( p^* \) is an \((S,D)\)-knighting condition for the servant \( \bar{q} = \langle q^\alpha_\delta : \alpha \in S \land u \in I_\alpha \rangle \). To this end consider the following strategy \( st^* \) of COM in \( \mathcal{D}_{S,D}^{\text{servant}}(\bar{q}, p^*, Q_{\Delta}^D) \). After arriving to a stage \( \alpha \in S \) of a play of \( \mathcal{D}_{S,D}^{\text{servant}}(\bar{q}, p^*, Q_{\Delta}^D) \), when \( (r_\beta, A_\beta : \beta < \alpha) \) has been already constructed, COM plays as follows.

If \( \alpha \) is a successor or \( \alpha \notin \bigcap_{\beta < \alpha} A_\beta \), then she just puts \( r_\alpha, A_\alpha \) such that:

(h) \( r_\beta \leq r_\alpha \) for all \( \beta < \alpha \) and if \( \delta \in C^\alpha \) is such that \( \text{otp}(\delta \cap C^\alpha) = \alpha \), then \( w^{\alpha} \setminus (\delta + 1) \neq \emptyset \), and

(i) \( A_\alpha = \bigcap_{\beta < \alpha} A_\beta \cap \bigcap_{\beta < \alpha} C^{\rho^\beta} \setminus (\sup(w^{\rho^\beta}) + 1) \).

If \( \alpha \in \bigcap_{\beta < \alpha} A_\beta \) is a limit ordinal, then COM first lets \( u = \bigcup_{\beta < \alpha} w^{\rho^\beta} \). It follows from (h)+(i) from earlier stages that \( u \subseteq C^\alpha \) and \( \text{otp}(\alpha \cap C^\alpha) = \alpha \). Note that \( \alpha \in \bigcap_{\beta < \alpha} C^{\rho^\beta} \), \( u \in \bigcap_{\beta < \alpha} \text{pos}^{\rho^\beta}(C^{\rho^\beta}, C^{\rho^\alpha}, \alpha) \) and \( u \in I_\alpha \). Let \( \alpha' = \min(C^{\rho^\alpha} \setminus (\alpha + 1)) \) and \( \alpha'' \in C^{\rho^\alpha}_\alpha \setminus (\alpha + 1) \). It follows from (c)+(g) that for each \( \beta < \alpha \) we have

\[
\begin{align*}
q^\alpha_\delta & \leq p^*_{\alpha+1} |_{\alpha, \gamma} C^\delta_{\alpha+1}(u) \\
& \leq p^*_{\alpha, \gamma} C^\delta_{\alpha}(u) \\
& \leq \bar{r}_\beta |_{\alpha''} C^\delta_{\alpha}(u).
\end{align*}
\]

Hence COM may choose a condition \( r_\alpha \geq q^\alpha_\delta \) stronger than all \( r_\beta \) (for \( \beta < \alpha \)) and satisfying (h). Then \( A_\alpha \) is given by (i).

It follows directly from the description of \( st^* \) that it is a winning strategy of COM in \( \mathcal{D}_{S,D}^{\text{servant}}(\bar{q}, p^*, Q_{\Delta}^D) \).

The master game \( \mathcal{D}_{S,D}^{\text{master}} \) used to define the property of being reasonably merry is essentially a variant of the A–reasonable boundedness game \( \mathcal{D}_{A}^{\text{master}} \) of [RS05, Def. 3.1]. The related bounding property was weakened in [RS, Def. 2.9] by introducing double A–reasonably completeness game \( \mathcal{D}_{A}^{\text{2master}} \). We may use these ideas to introduce a property much weaker than “reasonably merry”, though the description of the resulting notions becomes somewhat more complicated. As an award for additional complication we get a true preservation theorem, however. In the rest of this section, in addition to 6.1 we assume also

**Context 6.13.** \( \bar{\mu} = \langle \mu_\alpha : \alpha < \lambda \rangle \) is a sequence of cardinals below \( \lambda \) such that \( (\forall \alpha < \lambda) (\forall 0 \leq \mu_\alpha \leq \mu_\alpha^{\text{lh}}) \).

**Definition 6.14.** Let \( Q \) be a forcing notion.

1. A double \( Q \)-servant over \( S, \bar{\mu} \) is a sequence

\[
\bar{q} = (q^\delta : \delta \in S \land \gamma < \mu_\delta \cdot \xi_\delta)
\]

such that for \( \delta \in S \),

- \( 0 < \xi_\delta < \lambda \) and \( q^\delta_0 \in Q \) (for \( \gamma < \mu_\delta \cdot \xi_\delta \)),
- \( (\forall i, i' < \xi_\delta) (\forall j < \mu_\delta) (i' < i \Rightarrow q^\delta_{(i,j)} \leq q^\delta_{(i'+j,j)}) \).

(Here \( \mu_\delta \) is treated as an ordinal and \( \mu_\delta \cdot \xi_\delta \) is the ordinal product of \( \mu_\delta \) and \( \xi_\delta \).)
(2) Let $\bar{q}$ be a double $Q$–servant over $S, \bar{\mu}$ and let $q \in Q$. We define a game $\mathcal{D}^{2\text{ser}}_{S,D,\bar{\mu}}(\bar{q}, q, Q)$ as follows. A play of $\mathcal{D}^{2\text{ser}}_{S,D,\bar{\mu}}(\bar{q}, q, Q)$ lasts at most $\lambda$ steps during which the players, COM and INC, attempt to construct a sequence $\langle r_\alpha, A_\alpha : \alpha < \lambda \rangle$ such that

- $r_\alpha \in Q$, $q \leq r_\alpha$, $A_\alpha \in D$ and $\alpha < \beta < \lambda \Rightarrow r_\alpha \leq r_\beta$.

The terms $r_\alpha, A_\alpha$ are chosen successively by the two players so that

- if $\alpha \notin S$, then INC picks $r_\alpha, A_\alpha$, and
- if $\alpha \in S$, then COM chooses $r_\alpha, A_\alpha$.

If at some moment of the play one of the players has no legal move, then INC wins; otherwise, if both players always had legal moves and the sequence $\langle r_\alpha, A_\alpha : \alpha < \lambda \rangle$ has been constructed, then COM wins if and only if

$$\bigwedge (\forall \delta \in S) [(\exists \alpha < \delta) (A_\alpha \& \delta \text{ is limit}) \Rightarrow (\exists j \in \mu_\delta)(\forall i < \xi_\delta(q^\delta_{\mu_\delta + j} : q_\delta))].$$

(3) If COM has a winning strategy in the game $\mathcal{D}^{2\text{ser}}_{S,D,\bar{\mu}}(\bar{q}, q, Q)$, then we will say that $q$ is an knighting condition for the double servant $\bar{q}$.

Definition 6.15. Let $Q$ be a strategically ($<\lambda$)–complete forcing notion.

(1) For a condition $p \in Q$ we define a game $\mathcal{D}^{2\text{mas}}_{S,D,\bar{\mu}}(p, Q)$ between two players, Generic and Antigeneric. A play of $\mathcal{D}^{2\text{mas}}_{S,D,\bar{\mu}}(p, Q)$ lasts $\lambda$ steps and during a play a sequence

$$\langle \xi_\alpha, q^\alpha_i, q^\alpha_\delta : \gamma < \mu_\alpha \cdot \xi_\alpha : \alpha < \lambda \rangle.$$ 

is constructed. (Again, $\mu_\alpha \cdot \xi_\alpha$ is the ordinal product of $\mu_\alpha$ and $\xi_\alpha$.) Suppose that the players have arrived to a stage $\alpha < \lambda$ of the game. First, Antigeneric picks a non-zero ordinal $\xi_\alpha < \lambda$. Then the two players start a subgame of length $\mu_\alpha \cdot \xi_\alpha$ alternately choosing the terms of the sequence $\langle q^\alpha_i, q^\alpha_\delta : \gamma < \mu_\alpha \cdot \xi_\alpha \rangle$. At a stage $\gamma = \mu_\alpha \cdot i + j$ (where $i < \xi_\alpha, j < \mu_\alpha$) of the subgame, first Generic picks a condition $p^\alpha_i \in Q$ stronger than $p$ and stronger than all conditions $q^\alpha_\delta$ for $\delta < \gamma$ of the form $\delta = \mu_\alpha \cdot i' + j$ (where $i' < i$), and then Antigeneric answers with a condition $q^\alpha_\delta$ stronger than $p^\alpha_i$.

At the end, Generic wins the play if

(a) there were always legal moves for both players (so a sequence

$$\langle \xi_\alpha, q^\alpha_i, q^\alpha_\delta : \gamma < \mu_\alpha \cdot \xi_\alpha : \alpha < \lambda \rangle$$

has been constructed) and

(b) for each $\alpha \in S$, the conditions in $\{ p^\alpha_i : j < \mu_\alpha \}$ are pairwise incompatible, and

(c) letting

$$\bar{q} = \langle \xi_\delta, q^\delta_\delta : \delta \in S \& \gamma < \mu_\delta \cdot \xi_\delta \rangle$$

(it is a double $Q$–servant over $S$) we may find a knighting condition $q \geq p$ for the double servant $\bar{q}$.

(2) A forcing notion $Q$ is reasonably double merry over $(S, D, \bar{\mu})$ if (it is strategically ($<\lambda$)–complete and) Generic has a winning strategy in the game $\mathcal{D}^{2\text{mas}}_{S,D,\bar{\mu}}(p, Q)$ for any $p \in Q$.

Theorem 6.16. Assume that $\lambda, S, D, \bar{\mu}$ are as in 6.1+6.13. Let $\bar{Q} = \langle \bar{p}_\alpha, \bar{q}_\alpha : \alpha < \gamma \rangle$ be a $\lambda$–support iteration such that for each $\alpha < \gamma$:

$$\bar{p}_\alpha \ 	ext{"} \bar{Q}_\alpha \text{ is reasonably double merry over } (S, D, \bar{\mu}) \text{"}.$$
Then $P_\gamma = \lim(\bar{Q})$ is reasonably double merry over $(S, D, \bar{\mu})$ (so also $\lambda$–proper).

**Proof.** Combine the proof of [RS, Thm 2.12] (the description of the strategy here is the same as the one there) with the end of the proof of 6.4(a). □

7. Open problems

**Problem 7.1.** Let $p = (P, S, D)$ be a $\mathcal{D}\ell$–parameter. Does “reasonably $B$-bounding over $p$” (see 1.8) imply “reasonably $B$-bounding over $D$” (of [RS05, Def. 3.1])? Does “reasonably $B$-bounding over $p$” imply “$B$–noble over $p$”? (Note 6.6.)

**Problem 7.2.** Are there any relations among the notions of properness over $D$–semi diamonds (of [RS01]), properness over $D$–diamonds (of [Eis03]) and $B$–nobleness (of 3.1)?

**Problem 7.3.** Does $\lambda$–support iterations of forcing notions of the form $Q_{\lambda}E$ add $\lambda$–Cohen’s? Here we may look at iterations as in 3.3 or 2.7.

**Problem 7.4.** Does “reasonably $B$–bounding over $p$” (for a $\mathcal{D}\ell$–parameter $p$) imply “reasonably merry over $(S, D)$” (for some $S, D$ as in 6.1)?

**Problem 7.5.** Does “fuzzy proper over quasi $D$–diamonds for $W$” (see [RS07, Def. A.3.6]) imply “reasonably merry”? (Any result in this direction may require additional assumptions on $W, \bar{Q}$ in [RS07, A.3.1, A.3.3].)

**References**


Department of Mathematics, University of Nebraska at Omaha, Omaha, NE 68182-0243, USA  
E-mail address: roslanow@member.ams.org  
URL: http://www.unomaha.edu/logic

Einstein Institute of Mathematics, Edmond J. Safra Campus, Givat Ram, The Hebrew University of Jerusalem, Jerusalem, 91904, Israel, and Department of Mathematics, Rutgers University, New Brunswick, NJ 08854, USA  
E-mail address: shelah@math.huji.ac.il  
URL: http://shelah.logic.at