

## Model Companions of $T_{\text{Aut}}$ for Stable $T$

John T. Baldwin and Saharon Shelah

**Abstract** We introduce the notion  $T$  does not omit obstructions. If a stable theory does not admit obstructions then it does not have the finite cover property (nfcp). For any theory  $T$ , form a new theory  $T_{\text{Aut}}$  by adding a new unary function symbol and axioms asserting it is an automorphism. The main result of the paper asserts the following: If  $T$  is a stable theory,  $T$  does not admit obstructions if and only if  $T_{\text{Aut}}$  has a model companion. The proof involves some interesting new consequences of the nfcp.

### 1 Introduction

Let  $T$  be a complete first-order theory in a countable relational language  $L$ . We assume relation symbols have been added to make each formula equivalent to a predicate. Adjoin a new unary function symbol  $\sigma$  to obtain the language  $L_\sigma$ ;  $T_{\text{Aut}}$  is obtained by adding axioms asserting that  $\sigma$  is an  $L$ -automorphism.

The modern study of the model companion of theories with an automorphism has two aspects. One line, stemming from Lascar [7], deals with “generic” automorphisms of arbitrary structures. A second, beginning with Chatzidakis and Hrushovski [3] and questions of Macintyre about the Frobenius automorphism, is more concerned with specific algebraic theories. This paper is more in the first tradition: we find general necessary and sufficient conditions for a stable first-order theory with automorphism to have a model companion.

Kikyo in [4] investigates the existence of model companions of  $T_{\text{Aut}}$  when  $T$  is unstable. He also includes an argument of Kudaibergenov showing that if  $T$  is stable with the finite cover property then  $T_{\text{Aut}}$  has no model companion. This argument was implicit in Chatzidakis and Pillay [2] and is a rediscovery of a theorem of Winkler [11] in the 70s. We provide necessary and sufficient conditions for  $T_{\text{Aut}}$  to have a model companion when  $T$  is stable. Namely, we introduce a new condition,  $T$  admits obstructions, and show that  $T_{\text{Aut}}$  has a model companion if and only if  $T$  does

Received June 5, 2001; accepted March 18, 2003; printed August 29, 2003

2001 Mathematics Subject Classification: Primary, 03C45

Keywords: stability, expansion by automorphism

©2003 University of Notre Dame

not admit obstructions. This condition is a weakening of the finite cover property: if a stable theory  $T$  has the finite cover property then  $T$  admits obstructions.

Kikyo also proved that if  $T$  is an unstable theory without the independence property,  $T_{\text{Aut}}$  does not have a model companion. Kikyo and Shelah [6] have improved this by weakening the hypothesis to  $T$  has the strict order property.

For  $p$  a type over  $A$  and  $\sigma$  an automorphism,  $\sigma(p)$  denotes

$$\{\varphi(\mathbf{x}, \sigma(\mathbf{a})) : \varphi(\mathbf{x}, \mathbf{a}) \in p\}.$$

(References of the form ‘II.4.13’ are to Shelah [9].) Further related work is contained in Shelah [10] which investigates when  $T_{\text{Aut}}$  has a stable model completion.

## 2 Example

In the following example we examine exactly why a particular  $T_{\text{Aut}}$  does not have a model companion. Eventually we will show that the obstruction illustrated here represents the reason  $T_{\text{Aut}}$  (for stable  $T$ ) can fail to have a model companion. Let  $L$  contain two binary relation symbols  $E$  and  $R$  and unary predicates  $P_i$  for  $i < \omega$ . The theory  $T$  asserts that  $E$  and  $R$  are equivalence relations and that  $E$  has infinitely many infinite classes which are refined by  $R$  into two-element classes. Moreover, each  $P_i$  holds only elements from one  $E$ -class and contains exactly one element from each  $R$ -class of that  $E$ -class. Thus,  $x \neq y \wedge P_i(x) \wedge P_j(y)$  implies  $\neg R(x, y)$  if  $i \neq j$ .

Now  $T_{\text{Aut}}$  does not have a model companion. To see this, let  $\psi(x, y, z)$  be the formula,  $E(x, z) \wedge E(y, z) \wedge R(x, y) \wedge x \neq y$ . Let  $\Gamma$  be the  $L_\sigma$ -type in the variables  $\{z\} \cup \{x_i y_i : i < \omega\}$  which asserts  $\psi(x_i, y_i, z)$  holds for each  $i$ , the sequence  $\langle x_i y_i : i < \omega \rangle$  is  $L$ -indiscernible, the  $x_i$  are distinct and the  $y_i$  are distinct, and for every  $\varphi(x, \mathbf{w}) \in L(T)$ ,

$$(\forall \mathbf{w}) \bigvee_{i \in U} \{\bigwedge \varphi(x_i, \mathbf{w}) \leftrightarrow \varphi(y_i, \sigma(\mathbf{w})) : U \subseteq \lg(\mathbf{w}) + 3, |U| > (\lg(\mathbf{w}) + 3)/2\}.$$

Thus if  $\langle b_i c_i : i < \omega \rangle a$  realize  $\Gamma$  in a model  $M$ ,

$$\sigma(\text{avg}(\langle b_i : i < \omega \rangle / M)) = \text{avg}(\langle c_i : i < \omega \rangle / M).$$

For any finite  $\Delta \subset L(T)$ , let  $\chi_{\Delta, k}(\mathbf{x}, \mathbf{y}, z)$  be the conjunction of the  $\Delta$ -formulas satisfied by  $\langle b_i c_i : i < k \rangle a$  where  $\langle b_i c_i : i < k \rangle a$  are an initial segment of a realization of  $\Gamma$ . Let  $\theta_{\Delta, k}$  be the sentence

$$(\forall x_0, \dots, x_{k-1}, y_0, \dots, y_{k-1}, z) \chi_{\Delta, k}(x_0, \dots, x_{k-1}, y_0, \dots, y_{k-1}, z) \rightarrow (\exists x_0, y_0, x_1, y_1) [\psi(x_0, y_0, z) \wedge \psi(x_1, y_1, z) \wedge \sigma(x_1) = y_1].$$

We claim that if  $T_{\text{Aut}}$  has a model companion  $T_{\text{Aut}}^*$ , then for some  $k$  and  $\Delta$ ,

$$T_{\text{Aut}}^* \vdash \theta_{\Delta, k}.$$

For this, let  $M \models T_{\text{Aut}}^*$  such that  $\langle b_i c_i : i < k \rangle a$  satisfy  $\Gamma$  in  $M$ . Suppose  $M \upharpoonright L < N$  and  $N$  is an  $|M|^+$ -saturated model of  $T$ . In  $N$  we can find  $b, c$  realizing the average of  $\langle b_i : i < \omega \rangle$  and  $\langle c_i : i < \omega \rangle$  over  $M$ , respectively. Then

$$\sigma(\text{avg}(\langle b_i : i < \omega \rangle / M)) = \text{avg}(\langle c_i : i < \omega \rangle / M),$$

and so there is an automorphism  $\sigma^*$  of  $N$  extending  $\sigma$  and taking  $b$  to  $c$ . Since  $(M, \sigma)$  is existentially closed ( $T_{\text{Aut}}^*$  is model complete), we can pull  $b, c$  down to  $M$ . By compactness, some finite subset  $\Gamma_0$  of  $\Gamma$  suffices and letting  $\Delta$  be the formulas mentioned in  $\Gamma_0$  and  $k$  the number of  $x_i, y_i$  appearing in  $\Gamma_0$  we have the claim.

But now we show that if  $(M, \sigma)$  is any model of  $T_{\text{Aut}}$ , then for any finite  $\Delta$  and any  $k$ ,  $(M, \sigma) \models \neg\theta_{\Delta,k}$ . For this, choose  $b_i, c_i$  for  $i < k$  which are  $E$ -equivalent to each other and to an element  $a$  in a class  $P_j$  where  $P_j$  does not occur in  $\Delta$  and with  $R(b_i, c_i)$  and  $b_i \neq c_i$ . Then  $\mathbf{b}, \mathbf{c}, a$  satisfy  $\chi_{\Delta,k}$  but there are no  $b_k, c_k$  and automorphism  $\sigma$  which make  $\theta_{\Delta,k}$  true. So, for each  $j$ ,

$$T \vdash (\forall x, y, z)(\psi(x, y, z) \wedge P_j(z) \rightarrow [P_j(x) \leftrightarrow \neg P_j(y)]).$$

To put this situation in a more general framework, recall some notation from [9].  $\Delta$  will denote a finite set of formulas:  $\{\varphi_i(\mathbf{x}, \mathbf{y}_i) : \text{lg}(\mathbf{x}) = m, i < |\Delta|\}$ ;  $p$  is a  $\Delta$ - $m$ -type over  $A$  if  $p$  is a set of formulas  $\varphi_i(\mathbf{x}, \mathbf{a})$  where  $\mathbf{x} = \langle x_1, \dots, x_{m-1} \rangle$  (these specific variables), and  $\mathbf{a}$  from  $A$  is substituted for  $\mathbf{y}_i$ . Thus, if  $A$  is finite there are only finitely many  $\Delta$ - $m$ -types over  $A$ .

Now let  $\Delta_1$  contain Boolean combinations of  $x = y, R(x, y), E(x, y)$ . Let  $\Delta_2$  expand  $\Delta_1$  by adding a finite number of the  $P_j(z)$  and let  $\Delta_3$  contain  $P_k(x)$  where  $P_k$  does not occur in  $\Delta_2$ .

Now we have the following situation: there exists a set  $X = \{b_0, b_1, c_0, c_1, a\}$ ,  $P_j(a)$  holds, all five are  $E$ -equivalent, and  $R(b_i, c_i)$  for  $i = 0, 1$  such that

1.  $\langle \mathbf{b}_i \mathbf{c}_i : i \leq 12 \rangle$  is  $\Delta_2$ -indiscernible over  $a$ ;
2.  $\langle b_0 c_0, b_1 c_1 \rangle$  can be extended to an infinite set of indiscernibles  $\overline{\mathbf{bc}}$  which satisfy the following:
  - (a)  $\psi(b_i, c_i, a)$ ;
  - (b)  $\sigma(\text{avg}_{\Delta_2}(\overline{\mathbf{b}}/M)) = \text{avg}_{\Delta_2}(\overline{\mathbf{c}}/M)$ ;
3.  $\text{tp}_{\Delta_1}(\mathbf{b}_2 \mathbf{c}_2/X) \vdash \sigma(\text{tp}_{\Delta_3}(\mathbf{b}_2/X)) \neq \text{tp}_{\Delta_3}(\mathbf{c}_2/X)$ .

We call a sequence such as  $\langle \mathbf{b}_i \mathbf{c}_i : i \leq 2 \rangle \mathbf{a}$  a  $(\sigma, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over the empty set. In order to “finitize” the notions we will give below more technical formulations of the last two conditions; we will have to discuss obstructions over a finite set  $A$ . In the example, the identity was the only automorphism of the prime model. We will have to introduce a third sequence  $\mathbf{b}'$  to deal with arbitrary  $\sigma$ . But this example demonstrates the key aspects of obstruction that are the second reason for  $T_{\text{Aut}}$  to lack a model companion.

### 3 Preliminaries

In order to express the notions described in the example, we need several notions from basic stability theory. By working with finite sets of formulas in a stable theory without the finite cover property, we are able to refine arguments about infinite sets of indiscernibles to arguments about sufficiently long finite sequences.

We may speak recklessly of indiscernible sequence but in this paper we deal exclusively with  $\Delta$ -indiscernible sets which are defined just below. For infinite sequences in a stable theory such recklessness is without penalty (since infinite indiscernible sequences are indiscernible sets); since we are speaking of finite sequences, it is essential that we really mean indiscernible sets.

**Definition 3.1** Let  $\Delta$  be a finite set of formulas which we will assume to be closed under permutation or identification of variables and under negation;  $\neg\neg\varphi$  is identified with  $\varphi$ .

1. We say that  $E = \langle \mathbf{a}_i : i \in I \rangle$ , where all  $\mathbf{a}_i$  have the same length  $m$ , is  $(\Delta, r)$ -indiscernible if it satisfies the following conditions. Suppose  $i_0, \dots, i_{r-1}$  and

$j_0, \dots, j_{r-1}$  are distinct elements of  $I$  and for  $i < t < r$ ,  $u_i$  is a subset of  $m$  with  $\sum_{i < t} |u_i| = r$ ,  $\lg(\mathbf{x}_i) = u_i$ , and  $\varphi(\mathbf{x}_0, \dots, \mathbf{x}_{t-1}) \in \Delta$ . Then

$$\begin{aligned} \varphi(\mathbf{a}_{i_0} \upharpoonright u_0, \mathbf{a}_{i_1} \upharpoonright u_1, \dots, \mathbf{a}_{i_{t-1}} \upharpoonright u_{t-1}) &\leftrightarrow \\ \varphi(\mathbf{a}_{j_0} \upharpoonright u_0, \mathbf{a}_{j_1} \upharpoonright u_1, \dots, \mathbf{a}_{j_{t-1}} \upharpoonright u_{t-1}). \end{aligned}$$

2.  $E$  is  $\Delta$ -indiscernible if it is  $(\Delta, r)$ -indiscernible for all  $r$ , or equivalently for all  $r'$  with  $r'$  at most the maximum number of variables in a formula in  $\Delta$ .
3. For any sequence  $E = \langle \mathbf{a}_i : i \in I \rangle$  and  $j \in I$  we write  $E_j$  for  $\langle \mathbf{a}_i : i < j \rangle$ .

**Remark 3.2** Pedantically the formulas in  $\Delta$  contain variables only among  $x_0, \dots, x_n$  for some  $n$ , but we will freely write  $\varphi(x)$ ,  $\varphi(y)$  to increase intelligibility. We do not distinguish strictly between an arbitrary finite set of formulas  $\Delta$  and its closure described in Definition 3.1.

We will rely on the following facts/definitions from [9] to introduce two crucial functions for this paper:  $F(\Delta, n)$  and  $f(\Delta, n)$ .

**Fact 3.3** Recall that if  $T$  is stable, then for every finite  $\Delta \subset L(T)$  and  $n < \omega$  there is a finite  $\Delta' = F(\Delta, n)$  with  $\Delta \subseteq \Delta' \subset L(T)$  and a  $k^* = f(\Delta, n)$  with the following properties.

1. Assume we have finite set  $A$  and a set  $\mathbf{E} = \langle \mathbf{e}_i : i \in I \rangle$  of  $n$ -tuples such that for  $i < j$ ,

$$\text{tp}_{\Delta'}(\mathbf{e}_j/\mathbf{E}_i A) = \text{tp}_{\Delta'}(\mathbf{e}_i/\mathbf{E}_i A)$$

and

$$R_{(\Delta', 2)}(\mathbf{e}_j/\mathbf{E}_j A) = R_{(\Delta', 2)}(\mathbf{e}_i/\mathbf{E}_i A),$$

(whence,  $\text{tp}_{\Delta'}(\mathbf{e}_j/\mathbf{E}_i A)$  is definable over  $A$ ). Then  $\mathbf{E}$  is a set of  $\Delta$ -indiscernibles over  $A$ .

2. For any set of  $\Delta'$ -indiscernibles over the empty set,  $\mathbf{E} = \langle \mathbf{e}_i : i < k \rangle$  with  $\lg(\mathbf{e}_i) = n$  and  $k \geq k^*$  for any  $\varphi(\mathbf{u}, \mathbf{v}) \in \Delta$  and any  $\mathbf{d}$  with  $\lg(\mathbf{d}) = \lg(\mathbf{v}) = m$  either  $\{\mathbf{e}_i : \varphi(\mathbf{e}_i, \mathbf{d})\}$  or  $\{\mathbf{e}_i : \neg\varphi(\mathbf{e}_i, \mathbf{d})\}$  has strictly less than  $k^*/10$  elements. (II.4.13, II.2.20)
3. This implies that, for appropriate choice of  $k^*$ ,
  - (a) there is an integer  $m = m(\Delta, n) \geq n$  such that for any set of  $\Delta'$ -indiscernibles  $\langle \mathbf{e}_i : i < k \rangle$  over  $A$  with  $\lg(\mathbf{e}_i) = n$  and  $k \geq k^*$  and any  $\mathbf{a}$  with  $\lg(\mathbf{a}) \leq m$  there is a  $U \subseteq k$  with  $|U| < k^*/2$  such that  $\langle \mathbf{e}_i : i \in k - U \rangle$  is  $\Delta$ -indiscernible over  $\mathbf{Aa}$ ;
  - (b) moreover, if  $k \geq k^*$ , for any set  $A$ ,  $\text{avg}_{\Delta}(\langle \mathbf{e}_i : i < k \rangle/A)$  is well defined: namely,  $\text{avg}_{\Delta}(\langle \mathbf{e}_i : i < k \rangle/A) =$

$$\{\varphi(\mathbf{x}, \mathbf{a}) : |\{\mathbf{e}_i : i < k, \varphi(\mathbf{e}_i, \mathbf{a})\}| \geq \frac{k^*}{10}, \mathbf{a} \in A, \varphi(\mathbf{x}, \mathbf{y}) \in \Delta\}.$$

In (3a),  $m$  is the least  $k \geq n$  such that all  $\varphi \in \Delta$  have at most  $k$  free variables. But in (3b),  $\text{avg}_{\Delta}(\langle \mathbf{e}_i : i < k \rangle/A)$  need not be consistent. (Let  $A$  be all the members of one finite class in the standard fcp example.)

The closure conditions on  $\Delta$  given in Definition 3.1 guarantee the following.

**Fact 3.4** If  $\langle \mathbf{a}_i \mathbf{b}_i : i < \alpha \rangle$  is a set of  $\Delta$ -indiscernibles over  $\mathbf{a}$ , with the length of the  $\mathbf{a}_i$  equal to  $n$ , then  $\langle \mathbf{a}_i \mathbf{a}_i \mathbf{b}_i : i < \alpha \rangle$  and  $\langle \mathbf{b}_i : i < \alpha \rangle$  are  $\Delta$ -indiscernible sets as well.

**Definition 3.5** The theory  $T$  does not have the finite cover property if for every finite  $\Delta$  (considered with  $n$ -variable parameters) there is a  $k$  such that for any  $A$  and any  $\Delta$ -type  $q$  over  $A$ , if  $q$  is  $k$ -consistent then  $q$  is consistent. We require  $f(\Delta, n)$  to be greater than this  $k$ .

**Fact 3.6** If  $T$  does not have the finite cover property, in addition to Fact 3.3, we can choose  $k^* = f(\Delta, n)$  to satisfy the following conditions.

1. If  $\mathbf{E} = \langle \mathbf{e}_i : i < k^* \rangle$  is a set of  $n$ -tuples, which is  $\Delta'$ -indiscernible over the empty set, for any  $A$ ,  $\text{avg}_\Delta(\mathbf{E}/A)$  is a *consistent* complete  $\Delta$ -type over  $A$ .
2. Any set of  $\Delta'$ -indiscernibles (of  $n$ -tuples) with length at least  $k^*$  can be extended to one of infinite length (II.4.6).
3. For any pair of  $F(F(\Delta, n), n)$ -indiscernible sets  $\mathbf{E}^1 = \langle \mathbf{e}_i^1 : i < k \rangle$  and  $\mathbf{E}^2 = \langle \mathbf{e}_i^2 : i < k \rangle$  over  $\mathbf{a}$  with  $\text{lg}(\mathbf{e}_i^j) = n$  (for  $j = 1, 2$ ) and  $k \geq k^*$  such that

$$\text{avg}_{F(\Delta, n)}(\mathbf{E}^1/\mathbf{a}\mathbf{E}^1\mathbf{E}^2) = \text{avg}_{F(\Delta, n)}(\mathbf{E}^2/\mathbf{a}\mathbf{E}^1\mathbf{E}^2),$$

there exists  $\mathbf{J} = \langle \mathbf{e}_j : k < j < \omega \rangle$  such that both  $\mathbf{E}^1\mathbf{J}$  and  $\mathbf{E}^2\mathbf{J}$  are  $F(\Delta, n)$ -indiscernible over  $\mathbf{a}$ .

4. We express the displayed condition in (3) on  $\mathbf{E}^1, \mathbf{E}^2$  by the formula,  $\lambda_\Delta(\bar{\mathbf{e}}^1, \bar{\mathbf{e}}^2, \mathbf{a})$ , where  $\bar{\mathbf{e}}^i$  enumerates  $\mathbf{E}^i$ .
5. If  $\mathbf{E}^1$  and  $\mathbf{E}^2$  contained in a model  $M$  are  $F(\Delta, n)$ -indiscernible over  $\mathbf{a} \in M$  and each have length at least  $k^*$ , then  $M \models \lambda_\Delta(\bar{\mathbf{e}}^1, \bar{\mathbf{e}}^2, \mathbf{a})$  if and only if  $\text{avg}_\Delta(\mathbf{E}^2/M) = \text{avg}_\Delta(\mathbf{E}^1/M)$ .

**Proof** For (1), make sure that  $k^*$  is large enough that every  $\Delta$ -type which is  $k^*$ -consistent is consistent (II.4.4(3)). Now (3) follows by extending the common  $F(\Delta, n)$ -average of  $\mathbf{E}^1$  and  $\mathbf{E}^2$  over  $\mathbf{a}\mathbf{E}^1\mathbf{E}^2$  by (2). Finally, condition 5 holds by adapting the argument for III.1.8 from the set of all  $L$ -formulas to  $\Delta$ .  $\square$

Note that both  $F$  and  $f$  can be chosen to be increasing in  $\Delta$  and  $n$ .

The following observations culminate in a new consequence of nfcpc that will be used to reduce from “obstruction” to “simple obstruction” in Section 4. The first is III.3.4 of [9] or V.1.23 of Baldwin [1]. Just choose  $u$  such that  $\text{tp}(\mathbf{b}_*/X)$  does not fork over  $u$  and  $\text{tp}(\mathbf{b}_*/u)$  is stationary. Note that we do not assume the existence of an  $\mathbf{a}_*$ .

**Fact 3.7** Suppose  $T$  is stable, and further suppose that  $\langle \mathbf{b}_i : i < |T|^+ \rangle \cup \{\mathbf{b}_*\}$  and  $\mathbf{X} = \langle \mathbf{a}_i\mathbf{b}_i : i < |T|^+ \rangle$  are sequences of  $L$ -indiscernibles over  $\mathbf{a}$ . Then there is a  $\mathbf{U} \subseteq X$ ,  $\mathbf{U}$  indexed by an initial segment of cardinality  $\leq |T|$ , such that if  $\mathbf{X}_{\mathbf{U}, j}$  denotes  $\mathbf{X} - (\mathbf{U} \cup \{\mathbf{a}_j\mathbf{b}_j\})$ , for every  $\mathbf{b}_j \notin \mathbf{U}$ ,

$$\text{tp}(\mathbf{b}_*/\mathbf{X}_{\mathbf{U}, j}) = \text{tp}(\mathbf{b}_j/\mathbf{X}_{\mathbf{U}, j}).$$

Now we “finitize” this fact.

**Lemma 3.8** Fix  $n$  and  $m$ . For every finite set of formulas  $\Delta$  there exist  $k_0, k_1 < \omega$  and a finite set of formulas  $\Delta^+$  such that if  $k_2 \geq k_1$ ,  $\mathbf{a}_i, \mathbf{b}_i$  are sequences of length  $n$ ,  $\mathbf{a}$  is set of parameters of size  $m$ , while  $\mathbf{X} = \langle \mathbf{a}_i\mathbf{b}_i : i < k_2 \rangle$  and  $\langle \mathbf{b}_i : i < k_2 \rangle \cup \mathbf{b}_*$  are sets of  $\Delta^+$ -indiscernibles over  $\mathbf{a}$ , then there is a  $\mathbf{U} \subseteq \mathbf{X}$ ,  $|\mathbf{U}| < k_0$  such that if  $\mathbf{X}_{\mathbf{U}, j}$  denotes  $\mathbf{X} - (\mathbf{U} \cup \{\mathbf{a}_j\mathbf{b}_j\})$ , for every  $\mathbf{a}_j, \mathbf{b}_j \notin \mathbf{U}$ ,

$$\text{tp}_\Delta(\mathbf{b}_*/\mathbf{X}_{\mathbf{U}, j}) = \text{tp}_\Delta(\mathbf{b}_j/\mathbf{X}_{\mathbf{U}, j}).$$

**Proof** If not, for every  $t = \langle k_0, k_1, \Delta_1 \rangle$  with  $\Delta_1$  a finite set of formulas containing  $\Delta$  there is a  $k_2^t \geq k_1$ , a structure  $M^t$ , and a sequence  $X^t = \langle \mathbf{a}_i^t \mathbf{b}_i^t : i < k_2 \rangle$  and  $\mathbf{b}_*^t$  such that  $\mathbf{X}$  and  $\langle \mathbf{b}_i^t : i < k_2 \rangle \cup \{\mathbf{b}_*\}$  are sequences of  $\Delta^+$ -indiscernibles over  $\mathbf{a}$  contained in  $M^t$  but for every  $\mathbf{U} \subseteq X$ , with  $|\mathbf{U}| < k_0^t$ , for some  $\mathbf{b}_j \notin \mathbf{U}$ ,

$$\text{tp}_\Delta(\mathbf{b}_*/X_{\mathbf{U},j}) \neq \text{tp}_\Delta(\mathbf{b}_j/X_{\mathbf{U},j}).$$

Now expand  $L$  by adding constants for  $\mathbf{a}$ ,  $\mathbf{b}_*$  and a new  $2n$ -ary predicate symbol  $P$ . Then  $(M^t, P)$ , where  $P$  holds of each tuple  $\langle \mathbf{a}_i^t \mathbf{b}_i^t \rangle$ , satisfies the first-order sentence expressing this failure. By compactness and saturation we obtain a contradiction to Fact 3.7.  $\square$

**Lemma 3.9** *Suppose  $T$  is stable without the finite cover property. For every finite  $\Delta$ ,  $n$  there are  $k_0, k_1$  and  $\Delta_1$  such that if  $\mathbf{X} = \langle \mathbf{a}_i \mathbf{b}_i : i < k_1 \rangle$ ,  $\mathbf{X}' = \langle \mathbf{b}_i : i < k_1 \rangle$ , and  $X' \cup \mathbf{b}_*$  are sequences of  $\Delta_1$ -indiscernibles (of sequences of length  $n$ ) there is  $V \subseteq k_1$  with  $|V| < k_0$  and there is an  $\mathbf{a}_*$  such that  $\langle \mathbf{a}_i \mathbf{b}_i : i \in k_1 - V \rangle \cup \{\mathbf{a}_* \mathbf{b}_*\}$  is a sequence of  $\Delta$ -indiscernibles.*

**Proof** Choose  $\Delta^+, k_0, k_1$  according to Lemma 3.8. Choose  $\mathbf{a}_{k_1}, \mathbf{b}_{k_1}$  to realize the average of  $\mathbf{X}$  and let  $q(\mathbf{x}, \mathbf{y})$  denote the  $\Delta$ -type of  $\mathbf{a}_{k_1}, \mathbf{b}_{k_1}$  over  $\mathbf{X}$ . Since the finite cover property fails there is an  $r$  such that the consistency of  $q$  is determined by its  $r$ -element subsets. Let  $\Delta_1$  contain  $\Delta^+$  and all formulas of the form  $\bigwedge_{i < r} (\exists \mathbf{x}) \varphi_i(\mathbf{x}, \mathbf{y}, \mathbf{u}_1, \mathbf{v}_1, \dots, \mathbf{u}_s, \mathbf{v}_s)$  where  $\varphi_i(\mathbf{a}_{k_1}, \mathbf{b}_{k_1}, \mathbf{a}_{i_1}, \mathbf{b}_{i_1}, \dots, \mathbf{a}_{i_s}, \mathbf{b}_{i_s})$  holds and  $\varphi_i \in \Delta$ . Now if  $\mathbf{X}$  and  $\mathbf{X}'$  satisfy the hypotheses for this choice of  $\Delta_1$ , by Lemma 3.8,  $q(\mathbf{x}, \mathbf{b}_*)$  is consistent as required.  $\square$

#### 4 Obstructions

In this section we introduce the main new notion of this paper: obstruction. We are concerned with a formula  $\psi(\mathbf{x}, \mathbf{y}, \mathbf{z})$  where  $\text{lg}(\mathbf{x}) = \text{lg}(\mathbf{y}) = n$  and  $\text{lg}(\mathbf{z}) = m$ . We will apply Facts 3.3 and 3.6 with  $\mathbf{e}_i = \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i$  where each of  $\mathbf{b}_i, \mathbf{b}'_i$ , and  $\mathbf{c}_i$  has length  $n$ . Thus, our exposition will depend on functions  $F(\Delta, 3n), f(\Delta, 3n)$ . In several cases, we apply Fact 3.3 with  $\varphi(\mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_3, \mathbf{v})$  as  $\theta(\mathbf{u}_2, \mathbf{v}) \leftrightarrow \theta(\mathbf{u}_3, \mathbf{v})$  for various  $\theta$ . The following notation is crucial to state the definition.

**Notation 4.1** If  $\bar{\mathbf{d}} = \langle \mathbf{d}_i : i < r \rangle$  is a sequence of  $3n$ -tuples, which is  $\Delta$ -indiscernible over a finite sequence  $\mathbf{f}$ , and  $r \geq k^* = f(\Delta, 3n)$ , then  $\tau_\Delta(\mathbf{z}, \bar{\mathbf{d}}\mathbf{f})$  is the formula with free variable  $\mathbf{z}$  of length  $3n$  and parameters  $\bar{\mathbf{d}}\mathbf{f}$  which asserts that there is a subsequence  $\bar{\mathbf{d}}'$  of  $\bar{\mathbf{d}}$  with length  $f(\Delta, 3n)$  so that  $\bar{\mathbf{d}}'\mathbf{z}$  forms a set of  $\Delta$ -indiscernibles over  $\mathbf{f}$ .

Now we come to the main notion. Intuitively,  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i < k \rangle \mathbf{a}$  is a  $(\Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A$  if  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i < k \rangle$  is an indefinitely extendible sequence of  $\Delta_2$ -indiscernibles over  $\mathbf{a}$  such that the  $\mathbf{b}_i$ s,  $\mathbf{b}'_i$ s, and  $\mathbf{c}_i$ s each have length  $n$ , and the  $\Delta_2$ -average of the  $\mathbf{b}'_i$ s and the  $\mathbf{c}_i$ s is the same (over any set) but any realizations of the  $\Delta_1$ -type of the  $\mathbf{b}'_i$  and the  $\Delta_1$ -type of the  $\mathbf{c}_i$  over  $\mathbf{a}$  and the sequence  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i < k \rangle$  have different  $\Delta_3$ -types over  $A$ . We state the definition more formally.

**Definition 4.2** Fix a finite  $\Delta_1 \subseteq \Delta_2 \subseteq L(T)$  and  $\Delta_3 \subseteq L(T)$ , finite  $\mathbf{a} \subseteq A \subset M \models T$  with  $\text{lg}(\mathbf{a}) \leq m(\Delta_2, n)$  (as in Fact 3.3),  $\sigma$  an automorphism of  $M$ , and a natural number  $n$ . We say  $\langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A$  if the following conditions hold.

1.  $\langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle$  is  $F(\Delta_2, 3n)$ -indiscernible over  $\mathbf{a}$ .
2.  $k \geq f(\Delta_2, 3n)$ ;  $\lg(\mathbf{b}_i) = \lg(\mathbf{b}'_i) = \lg(\mathbf{c}_i) = n$ .
3.  $\text{avg}_{\Delta_2}(\bar{\mathbf{e}}^1/M) = \text{avg}_{\Delta_2}(\bar{\mathbf{e}}^2/M)$   
where  $\bar{\mathbf{e}}^1 = \langle \sigma(\mathbf{b}_i) : i < k \rangle$  and  $\bar{\mathbf{e}}^2 = \langle \mathbf{c}_i : i < k \rangle$ .
4. Using the formula  $\tau_{\Delta_1}$  from Notation 4.1 with  $\mathbf{x}, \mathbf{x}', \mathbf{y}$  representing the free variable  $\mathbf{z}$  there, we have

$$M \models (\forall \mathbf{x}, \mathbf{x}', \mathbf{y}) [\tau_{\Delta_1}(\mathbf{x}, \mathbf{x}', \mathbf{y}, \langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle \mathbf{a}) \rightarrow \bigvee \{ \varphi(\mathbf{x}', \mathbf{f}) \wedge \neg \varphi(\mathbf{y}, \mathbf{f}) : \mathbf{f} \in A, \varphi \in \Delta_3 \}].$$

By Fact 3.6, condition 3 is expressed by a formula of  $\mathbf{e}^1, \mathbf{e}^2$ , and  $\mathbf{a}$ . Crucially, the hypothesis of condition 4 in Definition 4.2 is an  $L$ -formula with parameters  $\langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle \mathbf{a}$ ; the conclusion is an  $L$ -formula with parameters from  $A$  as well.  $\Delta_1$  and  $\Delta_2$  have  $3n$  type-variables;  $\Delta_3$  has  $n$  type-variables.

**Fact 4.3** Note that (a) if  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \leq k \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A$  and  $\Delta_1 \subseteq \Delta'_2 \subseteq \Delta_2$ , then  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \leq k \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta'_2, \Delta_3, n)$ -obstruction over  $A$ . Further, (b) if  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \leq k \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A$  and  $A \subseteq A'$ , where  $A'$  is finite,  $\Delta_1 \subseteq \Delta'_2 \subseteq \Delta_2$ , then  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \leq k \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta'_2, \Delta_3, n)$ -obstruction over  $A'$ . Finally, (c) if  $W \subseteq k$  is large enough then  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \in W \rangle \mathbf{a}$  is a  $(\sigma, \Delta_1, \Delta'_2, \Delta_3, n)$ -obstruction over  $A$ .

#### Definition 4.4

1. We say  $(M, \sigma) \models T_{\text{Aut}}$  has no  $\sigma$ -obstructions when there is a function  $G(\Delta_1, n)$  with  $F(\Delta_1, 3n) \subseteq G(\Delta_1, n) \subset_{<\omega} L(T)$  such that if  $\Delta_1$  is a finite subset of  $L(T)$  and  $G(\Delta_1, n)$  is contained in the finite  $\Delta_3 \subset L(T)$ , then for every finite subset  $A$  of  $M$ , there is no  $(\sigma, \Delta_1, G(\Delta_1, n), \Delta_3, n)$ -obstruction over  $A$ .
2. We say  $T$  has no obstructions when there is a function  $G(\Delta_1, n)$ —which does not depend on  $(M, \sigma)$ —such that for each  $(M, \sigma) \models T_{\text{Aut}}$ , if  $\Delta_1$  is a finite subset of  $L(T)$ ,  $A$  is finite subset of  $M$ , and  $\Delta_3$  is a finite subset of  $L(T)$ , there is no  $(\sigma, \Delta_1, G(\Delta_1, n), \Delta_3, n)$ -obstruction over  $A$ .

**Definition 4.5** A simple obstruction is an obstruction where the automorphism  $\sigma$  is the identity. The notions of a theory or model having a simple obstruction are the obvious modifications of the previous definition.

**Lemma 4.6**  $T$  has obstructions if and only if  $T$  has simple obstructions.

**Proof** Suppose  $T$  has obstructions; we must find simple obstructions; the other direction is obvious. So, suppose for some  $\Delta_1$ , and  $n$ , and for every finite  $\Delta_2 \supseteq F(\Delta_1, 3n)$ , there is a finite  $\Delta_3$  and a tuple  $(M^{\Delta_2}, \sigma^{\Delta_2}, A^{\Delta_2}, k^{\Delta_2})$  such that  $(M^{\Delta_2}, \sigma^{\Delta_2}) \models T_{\text{Aut}}$ ,  $A^{\Delta_2}$  is a finite subset of  $M^{\Delta_2}$  and  $\mathbf{b}^{\Delta_2}, \sigma(\mathbf{b}^{\Delta_2}), \mathbf{c}^{\Delta_2}, \mathbf{a}^{\Delta_2}$  contained in  $M^{\Delta_2}$  are a  $(\sigma^{\Delta_2}, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction of length  $k^{\Delta_2}$  over  $A^{\Delta_2}$ . Choose  $\Delta_2$  so that it contains the set of formulas  $\Delta_1^+$  associated to  $\Delta_1$  as  $\Delta_1$  is associated to  $\Delta$  in Lemma 3.9. Without loss of generality  $\lg(\mathbf{a}) = m = m(\Delta_1, 3n)$ , and we can choose an appropriate  $\Delta_3$  depending on  $\Delta_2$ . Now define a family of simple obstructions by replacing each component of the given sequence of obstructions by an appropriate object with left prefix *sim*.

$$\begin{aligned}
\text{sim}_A \Delta_2 &= A^{\Delta_2}. \\
\text{sim}(\mathbf{b}_i^{\Delta_2}) &= \text{sim}(\mathbf{b}'_i{}^{\Delta_2}) = \sigma(\mathbf{b}_i^{\Delta_2}). \\
\text{sim}_{\mathbf{c}_i} \Delta_2 &= \mathbf{c}_i^{\Delta_2}. \\
\text{sim}_{\mathbf{a}} \Delta_2 &= \mathbf{a}^{\Delta_2}. \\
\text{sim}_{\Delta_1} &= \Delta_1^+.
\end{aligned}$$

We use the same sets of formulas for the  $\Delta_2$  and  $\Delta_3$ .

We now have an obstruction with respect to the identity. For condition 1 this follows from Fact 3.4; conditions 2 and 3 are immediate; we check condition 4. Since the following argument is uniform  $\Delta_2$ , we omit the superscript  $\Delta_2$ . Examining the implication in condition 1 for both the original obstruction and the simple obstruction we see it suffices to show, if we write

$$\tau^1(\mathbf{x}', \mathbf{x}', \mathbf{y}) \quad \text{for} \quad \tau_{\Delta_1^+}(\mathbf{x}', \mathbf{x}', \mathbf{y}, \langle \sigma(\mathbf{b}_i)\sigma(\mathbf{b}_i)\mathbf{c}_i : i < k \rangle \mathbf{a})$$

and

$$\tau^2(\mathbf{x}, \mathbf{x}', \mathbf{y}) \quad \text{for} \quad \tau_{\Delta_1^+}(\mathbf{x}, \mathbf{x}', \mathbf{y}, \langle \mathbf{b}_i\sigma(\mathbf{b}_i)\mathbf{c}_i : i < k \rangle \mathbf{a}),$$

that

$$(\forall \mathbf{x}')(\forall \mathbf{y})[\tau^1(\mathbf{x}', \mathbf{x}', \mathbf{y}) \rightarrow (\exists \mathbf{x})\tau^2(\mathbf{x}, \mathbf{x}', \mathbf{y})].$$

We verify the implication. Lemma 3.9 (taking the  $\mathbf{b}_i$  there as  $\sigma(\mathbf{b}_i)$ ,  $\mathbf{c}_i$  here) implies for any  $\mathbf{b}'$ ,  $\mathbf{c}$  such that  $\langle \sigma(\mathbf{b}_i), \mathbf{c}_i : i < k \rangle \cup \{\mathbf{b}', \mathbf{c}\}$  is a sequence of  $\Delta_1^+$ -indiscernibles, there is a  $\mathbf{b}$  so that  $\langle \mathbf{b}_i\sigma(\mathbf{b}_i), \mathbf{c}_i : i < k \rangle \cup \{\mathbf{b}, \mathbf{b}', \mathbf{c}\}$  is a sequence of  $\Delta_1$ -indiscernibles. Thus  $[\tau^1(\mathbf{b}', \mathbf{b}', \mathbf{c}) \rightarrow \tau^2(\mathbf{b}, \mathbf{b}', \mathbf{c})]$  and we finish.  $\square$

**Lemma 4.7** *If  $T$  is a stable theory with the finite cover property then  $T$  has a simple obstruction.*

**Proof** By II.4.4 of [9], there is a formula  $E(\mathbf{x}, \mathbf{y}, \mathbf{z})$  such that for each  $\mathbf{d}$ ,  $E(\mathbf{x}, \mathbf{y}, \mathbf{d})$  is an equivalence relation and for arbitrarily large  $n$  there is a  $\mathbf{d}_n$  such that  $E(\mathbf{x}, \mathbf{y}, \mathbf{d}_n)$  has exactly  $n$  classes. Let  $\Delta_1$  be  $\{E(\mathbf{x}, \mathbf{y}, \mathbf{z}), \neg E(\mathbf{x}, \mathbf{y}, \mathbf{z})\}$  and consider any  $\Delta_2$ . Fix  $\text{lg}(\mathbf{x}) = \text{lg}(\mathbf{y}) = r$ . There are arbitrarily long sequences  $\mathbf{b}_n = \langle \mathbf{b}'_j : j < n \rangle$  such that for some  $\mathbf{d}_n$ ,  $\mathbf{b}_n$  is a set of representatives for distinct classes of  $E(\mathbf{x}, \mathbf{y}, \mathbf{d}_n)$ . So by Ramsey, for any  $\Delta_2$  we can find such  $\mathbf{b}_k$  where  $k = f(\Delta_2, 3r)$  and  $\mathbf{b}_k$  is a sequence of length  $k$  of  $F(\Delta_2, 3n)$ -indiscernibles over the empty set and  $\mathbf{b}_k$  is a set of representatives for distinct classes of  $E(\mathbf{x}, \mathbf{y}, \mathbf{d}_n)$ . Now, if  $\Delta_3$  contains formulas which express that the number of equivalence classes of  $E(\mathbf{x}, \mathbf{y}, \mathbf{z})$  is greater than  $n$  and  $A$  contains representatives of all equivalence classes of  $E(\mathbf{x}, \mathbf{y}, \mathbf{d}_n)$  we have an (identity,  $\Delta_1, \Delta_2, \Delta_3, r$ )-obstruction over  $A$ . (Let  $\mathbf{b} = \mathbf{b}' = \mathbf{c} = \mathbf{b}_k$ ; the hypothesis of clause 4 of Definition 4.2 is trivially false so the condition is satisfied.)  $\square$

## 5 Model Companions of $T_{\text{Aut}}$

In this section we establish necessary and sufficient conditions on stable  $T$  for  $T_{\text{Aut}}$  to have a model companion. First, we notice when the model companion, if it exists, is complete.

Note that  $\text{acl}(\emptyset) = \text{dcl}(\emptyset)$  in  $C_T^{eq}$  means every finite equivalence relation  $E(\mathbf{x}, \mathbf{y})$  of  $T$  is defined by a finite conjunction:  $\bigwedge_{i < n} \varphi_i(\mathbf{x}) \leftrightarrow \varphi_i(\mathbf{y})$ .

**Fact 5.1**

1. If  $T$  is stable,  $T_{\text{Aut}}$  has the amalgamation property.



2. If, in addition,  $\text{acl}(\emptyset) = \text{dcl}(\emptyset)$  in  $C_T^{eq}$  then  $T_{\text{Aut}}$  has the joint embedding property.

**Proof** The first part of this lemma was proved by Theorem 3.3 of [7] using the definability of types. For the second part, the hypothesis implies that types over the empty set are stationary and the result follows by similar arguments.  $\square$

**Lemma 5.2** *Suppose  $T$  is stable and  $T_{\text{Aut}}$  has a model companion  $T_\sigma^*$ .*

1. *Then  $T_\sigma^*$  is complete if and only if  $\text{acl}(\emptyset) = \text{dcl}(\emptyset)$  in  $C_T^{eq}$ .*
2. *If  $(M, \sigma) \models T_{\text{Aut}}$  then the union of the complete diagram of  $M$  (in  $L$ ) with the diagram of  $(M, \sigma)$  and  $T_\sigma^*$  is complete.*

**Proof** (1) We have just seen that if  $\text{acl}(\emptyset) = \text{dcl}(\emptyset)$  in  $C_T^{eq}$ , then  $T_{\text{Aut}}$  has the joint embedding property; this implies in general that the model companion is complete. If  $\text{acl}(\emptyset) \neq \text{dcl}(\emptyset)$  in  $C_T^{eq}$ , let  $E(\mathbf{x}, \mathbf{y})$  be a finite equivalence relation witnessing  $\text{acl}(\emptyset) \neq \text{dcl}(\emptyset)$ . Because  $E$  is a finite equivalence relation,

$$T_1 = T_{\text{Aut}} \cup \{(\forall \mathbf{x})E(\mathbf{x}, \sigma(\mathbf{x}))\}$$

is a consistent extension of  $T_{\text{Aut}}$ . But since

$$T_{\text{Aut}} \cup \{\neg E(\mathbf{x}, \mathbf{y})\} \cup \{\varphi(\mathbf{x}) \leftrightarrow \varphi(\mathbf{y}) : \varphi \in L(T)\}$$

is consistent, so is

$$T_2 = T_{\text{Aut}} \cup \{(\exists \mathbf{x})\neg E(\mathbf{x}, \sigma(\mathbf{x}))\}.$$

But  $T_1$  and  $T_2$  are contradictory, so  $T_\sigma^*$  is not complete.

(2) Since we have joint embedding (from amalgamation over any model) the result follows as in Fact 5.1.  $\square$

We now prove the equivalence of three conditions. The first is a condition on a pair of models. The second is given by an infinite set of  $L_\sigma$ -sentences (take the union over all finite  $\Delta_2$ ) and the average requires names for infinitely many elements of  $M$ . The third is expressed by a single first-order sentence in  $L_\sigma$ . The equivalence of the first and third suffices (Theorem 5.8) to show the existence of a model companion. In fact, (1) implies (2) implies (3) requires only stability; the nfcp is used to prove (3) implies (1).

**Lemma 5.3** *Suppose  $T$  is stable without the fcp. Let  $(M, \sigma) \models T_{\text{Aut}}$ ,  $\mathbf{a} \in M$  and suppose that  $(M, \sigma)$  has no  $\sigma$ -obstructions. Fix  $\psi(\mathbf{x}, \mathbf{y}, \mathbf{z}) \in L_T$  with  $\text{lg}(\mathbf{x}) = \text{lg}(\mathbf{y}) = n$  and  $\text{lg}(\mathbf{z}) = \text{lg}(\mathbf{a}) = m$ . The following three assertions are equivalent.*

1. *There exists  $(N, \sigma)$ ,  $(M, \sigma) \subseteq (N, \sigma) \models T_{\text{Aut}}$  and*

$$N \models (\exists \mathbf{x}\mathbf{y})[\psi(\mathbf{x}, \mathbf{y}, \mathbf{a}) \wedge \sigma(\mathbf{x}) = \mathbf{y}].$$

2. *Fix  $\Delta_1 = \{\psi(\mathbf{x}, \mathbf{y}, \mathbf{z})\}$  and without loss of generality  $\text{lg}(\mathbf{z}) \leq m(\Delta_1, n)$ . For  $k \geq 5 \cdot f(\Delta_1, 3n)$  and any finite  $\Delta_2 \supseteq F(\Delta_1, 3n)$  (Fact 3.3), there are  $\mathbf{b}_i\sigma(\mathbf{b}_i)\mathbf{c}_i \in {}^{3n}M$  for  $i < k$  such that*

(a)  $\langle \mathbf{b}_i\sigma(\mathbf{b}_i)\mathbf{c}_i : i < k \rangle$  is  $F(\Delta_2, 3n)$ -indiscernible over  $\mathbf{a}$ ,

(b) for each  $i < k$ ,  $\psi(\mathbf{b}_i, \mathbf{c}_i, \mathbf{a})$  holds,

(c) for every  $\mathbf{d} \in {}^m M$  and  $\varphi(\mathbf{u}, \mathbf{v}) \in \Delta_2$  we have

$$|\{i < k : \varphi(\sigma(\mathbf{b}_i), \mathbf{d}) \leftrightarrow \varphi(\mathbf{c}_i, \mathbf{d})\}| \geq f(\Delta_2, 3n)/2.$$

3. Let  $\Delta_2 = G(\Delta_1, n)$ . Then there are  $\mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i \in {}^{3n}M$  for  $i < k = 5 \cdot f(\Delta_2, 3n)$  such that (for  $\lambda_{\Delta_2}$  from Fact 3.6 (4)),
- (a)  $\langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle$  is  $F(G(\Delta_1, n), 3n)$ -indiscernible over  $\mathbf{a}$ ,
  - (b) for each  $i < k$ ,  $\psi(\mathbf{b}_i, \mathbf{c}_i, \mathbf{a})$ ,
  - (c)  $\lambda_{\Delta_2}(\langle \sigma(\mathbf{b}_i) : i < k \rangle, \langle \mathbf{c}_i : i < k \rangle, \mathbf{a})$ .

**Proof** First we show (1) implies (2). Fix  $\mathbf{b}, \mathbf{c} \in N$  with  $N \models \psi(\mathbf{b}, \mathbf{c}, \mathbf{a}) \wedge \sigma(\mathbf{b}) = \mathbf{c}$ . For  $\Delta_2$ , let  $\Delta_2^+ = F(F(\Delta_2, 3n), 3n)$ . For each  $\Delta_2$ , choose a finite  $p \subseteq \text{tp}_{L(T)}(\mathbf{b}, \mathbf{c}/M)$  with the same  $(\Delta_2^+, 2)$  rank as  $\text{tp}_{\Delta_2}(\mathbf{b}, \mathbf{c}/M)$  (so  $\text{tp}_{\Delta_2}(\mathbf{b}, \mathbf{c}/M)$  is definable over  $\text{dom } p$ ). Now inductively construct (by Fact 3.3) an  $F(\Delta_2, 3n)$ -indiscernible sequence  $\langle \mathbf{b}_i, \mathbf{c}_i : i < \omega \rangle$  by choosing  $\mathbf{b}_i, \mathbf{c}_i$  in  $M$  realizing the restriction of  $\text{tp}_{\Delta_2^+}(\mathbf{b}, \mathbf{c}/M)$  to  $\text{dom } p$  along with the points already chosen. Let  $\mathbf{b}'_i = \sigma(\mathbf{b}_i)$ . For some infinite  $U \subseteq \omega$ ,  $\langle \mathbf{b}_i, \mathbf{b}'_i, \mathbf{c}_i : i \in U \rangle$  is  $F(\Delta_2, 3n)$ -indiscernible over  $\mathbf{a}$ ; renumbering let  $U = \omega$ . Now conditions (2a) and (2b) of assertion 2 are clear. For clause (2c),

$$\begin{aligned} \text{avg}_{\Delta_2}(\langle \mathbf{c}_i : i < \omega \rangle/M) &= \text{tp}_{\Delta_2}(\mathbf{c}, M) \\ &= \sigma(\text{tp}_{\Delta_2}(\mathbf{b}, M)) \\ &= \sigma(\text{avg}_{\Delta_2}(\langle \mathbf{b}_i : i < \omega \rangle/M)). \end{aligned}$$

The first and last equalities hold by the choice of the  $\mathbf{b}_i, \mathbf{c}_i$  and the middle since  $\sigma(\mathbf{b}) = \mathbf{c}$ . So, for each  $\varphi \in \Delta_2$  and each  $\mathbf{d} \in M$  of appropriate length,

$$\varphi(\mathbf{x}, \mathbf{d}) \in \text{avg}_{\Delta_2}(\langle \mathbf{c}_i : i < \omega \rangle/M)$$

if and only if

$$\varphi(\mathbf{x}, \sigma^{-1}(\mathbf{d})) \in \text{avg}_{\Delta_2}(\langle \mathbf{b}_i : i < \omega \rangle/M).$$

So for some  $S_1, S_2 \subset \omega$  with  $|S_1|, |S_2| < f(\Delta_2, 3n)/2$ , we have for all  $i \in \omega - (S_1 \cup S_2)$ ,  $\varphi(\mathbf{c}_i, \mathbf{d})$  if and only if  $\varphi(\mathbf{b}_i, \sigma^{-1}(\mathbf{d}))$ . Since  $\sigma$  is an automorphism of  $M$  this implies for  $i \in \omega - (S_1 \cup S_2)$ ,  $\varphi(\mathbf{c}_i, \mathbf{d})$  if and only if  $\varphi(\sigma(\mathbf{b}_i), \mathbf{d})$  which gives condition (2c) by using the first  $k$  elements of  $\langle \mathbf{b}_i, \sigma(\mathbf{b}_i) \mathbf{c}_i : i \in \omega - S_1 \cup S_2 \rangle$ .

(3) is a special case of (2). To see this, note that (3c) is easily implied by the form analogous to (2c): For every  $m \leq m(\Delta_1, n)$  and  $\mathbf{d} \in {}^m M$  and  $\varphi(\mathbf{u}, \mathbf{v}) \in G(\Delta_1, n)$  we have

$$\{|i < k : \varphi(\sigma(\mathbf{b}_i), \mathbf{d}) \leftrightarrow \varphi(\mathbf{c}_i, \mathbf{d})\} \geq f(\Delta_2, 3n)/2.$$

If  $T$  does not have fcp (3c) implies (2c) holds and we use that fact implicitly in the following argument. It remains only to show that (3) implies (1) with  $\Delta_1 = \{\psi\}$  and  $\Delta_2 = G(\Delta_1, n)$ . Without loss of generality we may assume  $N$  is  $\aleph_1$ -saturated. We claim the type

$$\Gamma = \{\psi(\mathbf{x}, \mathbf{y}, \mathbf{a})\} \cup \{\varphi(\mathbf{x}, \mathbf{d}) \leftrightarrow \varphi(\mathbf{y}, \sigma(\mathbf{d})) : \mathbf{d} \in M, \varphi \in L(T)\} \cup \text{diag}(M)$$

is consistent. This clearly suffices.

Let  $k = f(\Delta_2, 3n)$ . Suppose  $\langle \mathbf{b}_i \sigma(\mathbf{b}_i) \mathbf{c}_i : i < k \rangle \mathbf{a}$  satisfy (3). Let  $\Gamma_0$  be a finite subset of  $\Gamma$  and suppose only formulas from the finite set  $\Delta_3$  and only parameters from the finite set  $A$  appear in  $\Gamma_0$ . Write  $\mathbf{b}'_i$  for  $\sigma(\mathbf{b}_i)$ .

Now  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \leq f(\Delta_2, 3n) \rangle \mathbf{a}$  easily satisfy conditions 1 and 2 of Definition 4.2 for being a  $(\Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A$  and, in view of Fact 3.6 (3), (4), the third is given by condition (3c). Since there is no obstruction, condition 4 must fail. So there exist  $\mathbf{b}^*, (\mathbf{b}^*)', \mathbf{c}^*$  so that

$$M \models \tau_{\Delta_1}(\mathbf{b}^*, (\mathbf{b}^*)', \mathbf{c}^*, \langle \mathbf{b}_i, \mathbf{b}'_i, \mathbf{c}_i : i < k \rangle \mathbf{a}),$$

and  $\text{tp}_{\Delta_3}((\mathbf{b}^*)'/A) = \text{tp}_{\Delta_3}(\mathbf{c}^*/A)$  so  $\Gamma_0$  is satisfiable.  $\square$

As we will note in Theorem 5.8, we have established a sufficient condition for  $T_{\text{Aut}}$  to have a model companion. The next argument shows it is also necessary.

**Lemma 5.4** *Suppose  $T$  is stable; if  $T$  has an obstruction then  $T_{\text{Aut}}$  does not have a model companion.*

**Proof** We may assume  $T$  does not have fcp since if it does we know by Winkler [11] and Kudaibergenov [4] that  $T_{\text{Aut}}$  does not have a model companion. By Lemma 4.6, we may assume  $T$  has a simple obstruction. So it suffices to prove the following:

Suppose for some  $\Delta_1$ , and  $n$ , and for every finite  $\Delta_2 \supseteq F(\Delta_1, 3n)$ , there is a finite  $\Delta_3$  and a tuple  $(M^{\Delta_2}, \text{id}_{M^{\Delta_2}}, A^{\Delta_2}, k^{\Delta_2})$  such that  $(M^{\Delta_2}, \text{id}_{M^{\Delta_2}}) \models T_{\text{Aut}}$ ,  $A^{\Delta_2}$  is a finite subset of  $M^{\Delta_2}$ ,  $\mathbf{b}^{\Delta_2}, \sigma^{\Delta_2}(\mathbf{b}^{\Delta_2}) = \text{id}_{M^{\Delta_2}}(\mathbf{b}^{\Delta_2}) = \mathbf{b}^{\Delta_2}, \mathbf{c}^{\Delta_2}, \mathbf{a}^{\Delta_2}$  contained in  $M^{\Delta_2}$  are an  $(\text{id}_{M^{\Delta_2}}, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction of length  $k^{\Delta_2}$  over  $A^{\Delta_2}$ .

Then the collection  $\mathbf{K}_\sigma$  of existentially closed models of  $T_{\text{Aut}}$  is not an elementary class.

Without loss of generality  $\text{lg}(\mathbf{a}) = m = m(\Delta_1, 3n)$  and we can write  $\Delta_3 = \Delta_3(\Delta_2)$ . By the usual coding we may assume  $\Delta_1 = \{\psi(\mathbf{x}, \mathbf{y}, \mathbf{z})\}$  with  $\text{lg}(\mathbf{x}) = \text{lg}(\mathbf{y}) = n$ ,  $\text{lg}(\mathbf{z}) = m$ ,  $k = f(\Delta_1, 3n)$ . Again, without loss of generality each  $(M^{\Delta_2}, \text{id}_{M^{\Delta_2}})$  can be expanded to an existentially closed model  $(M^{\Delta_2}, \sigma_{M^{\Delta_2}})$  such that  $\sigma_{M^{\Delta_2}}$  fixes  $A^{\Delta_2}, \mathbf{a}^{\Delta_2}, \mathbf{b}^{\Delta_2}$ , and  $\mathbf{c}^{\Delta_2}$  pointwise. Let  $\mathcal{D}$  be a nonprincipal ultrafilter on  $Y = \{\Delta_2 : F(\Delta_1, 3n) \subseteq \Delta_2 \subset_\omega L(T)\}$  such that for any  $\Delta \in Y$  the family of supersets in  $Y$  of  $\Delta$  is in  $\mathcal{D}$ . Expand the language  $L$  to  $L^+$  by adding a unary function symbol  $\sigma$ , a new unary predicate symbol  $P$ , a  $3n$ -ary relation symbol  $Q$  and constants  $\mathbf{a}$ . Expand each of the  $M^{\Delta_2}$  to an  $L^+$ -structure  $N_{\Delta_2}$  by interpreting  $P$  as  $A^{\Delta_2}$ ,  $\mathbf{a}$  as  $\mathbf{a}^{\Delta_2}$ ,  $\sigma$  as  $\sigma^{\Delta_2} = \text{id}_{M^{\Delta_2}}$  and  $Q$  as the set  $\{\mathbf{b}_i^{\Delta_2} \mathbf{b}_i^{\Delta_2} \mathbf{c}_i^{\Delta_2} : i < k^{\Delta_2}\}$  of  $3n$ -tuples. Let  $N^*$  be the ultraproduct of the  $N_{\Delta_2}$  modulo  $\mathcal{D}$ . Let  $A$  denote  $P(N^*)$ ,  $\mathbf{a}^*$  denote the ultraproduct of the  $\mathbf{a}^{\Delta_2}$ , and  $\langle \mathbf{b}_i \mathbf{b}_i \mathbf{c}_i : i \in I \rangle$  enumerate  $Q(N^*)$ .

**Claim 5.5**

1.  $\text{lg}(\mathbf{b}_i) = \text{lg}(\sigma(\mathbf{b}_i)) = \text{lg}(\mathbf{c}) = n$ ;  $\text{lg}(\mathbf{a}) = m$ .
2.  $\langle \mathbf{b}_i \mathbf{b}_i \mathbf{c}_i : i \in I \rangle$  is a sequence of  $L(T)$ -indiscernibles over  $\mathbf{a}^*$ .
3. For each finite  $\Delta_2 \subseteq L(T)$  with  $\Delta_2 \supseteq F(\Delta_1, 3n)$  and each finite subsequence from  $\langle \mathbf{b}_i \mathbf{b}_i \mathbf{c}_i : i \in I \rangle$  indexed by  $J$  of length at least  $k = f(\Delta_2, 3n)$  the  $\Delta_2$ -type of  $\langle \mathbf{b}_i \mathbf{b}_i \mathbf{c}_i : i \in J \rangle \mathbf{a}$  is the  $\Delta_2$ -type of some  $(\text{id}, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction  $(\mathbf{b}^{\Delta_2}, \mathbf{b}^{\Delta_2}, \mathbf{c}^{\Delta_2}) \mathbf{a}^{\Delta_2}$  in  $M^{\Delta_2}$  over  $A^{\Delta_2}$ .
4.  $\text{av}_{\text{GL}}(\langle \mathbf{b}_i : i \in I \rangle / N^*) = \text{av}_{\text{GL}}(\langle \mathbf{c}_i : i \in I \rangle / N^*)$ .

**Proof** This claim follows directly from the properties of ultraproducts. (For item 3, apply Fact 4.3 and the definition of the ultrafilter  $\mathcal{D}$ .)  $\square$

Let  $\Gamma$  be the  $L$ -type in the variables  $\langle \mathbf{x}_i \mathbf{x}'_i \mathbf{y}_i : i \in I \rangle \cup \{\mathbf{z}\}$  over the empty set of  $\langle \mathbf{b}_i, \mathbf{b}_i, \mathbf{c}_i : i \in I \rangle \mathbf{a}^*$ . For any finite  $\Delta \subset L(T)$ , let  $\chi_{\Delta, k}(\overline{\mathbf{x}\mathbf{x}'\mathbf{y}}, \mathbf{z})$  be the conjunction of the  $\Delta$ -type over the empty set of a subsequence of  $k$  elements from  $\langle \mathbf{b}_i \mathbf{b}_i \mathbf{c}_i : i \in I \rangle$  and  $\mathbf{a}^*$  from a realization of  $\Gamma$  with  $\mathbf{z}$  for  $\mathbf{a}^*$ .

**Notation 5.6** Recall the definition of  $\tau_{\Delta_1}$  from Notation 4.1. Let  $r = f(\Delta_1, n)$  and let  $\theta_{\Delta_1}(\mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z})$  be the formula

$$\begin{aligned} & (\exists \mathbf{x}, \mathbf{x}', \mathbf{y}) [\chi_{\Delta_1, r}(\mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z}) \\ & \wedge \tau_{\Delta_1}(\mathbf{x}, \mathbf{x}', \mathbf{y}, \mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z}) \wedge \sigma(\mathbf{x}') = \mathbf{y}]. \end{aligned}$$

Without loss of generality we assume  $0, 1, \dots, r-1$  index disjoint sequences.

**Claim 5.7** If  $\mathbf{K}_\sigma$ , the family of existentially closed models of  $T_{\text{Aut}}$ , is axiomatized by  $T_{\text{Aut}}^*$ , then

$$T_{\text{Aut}}^* \cup \Gamma \cup \{\sigma(\mathbf{x}_i) = \mathbf{x}'_i : i \in I\} \vdash \theta_{\Delta_1}(\mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z}).$$

(Abusing notation we write this with the  $\mathbf{x}_i, \mathbf{x}'_i, \mathbf{y}_i$  for  $i \in I$  free.)

**Proof** Note for each  $i$ ,  $\mathbf{x}_i = \mathbf{x}'_i$  is in  $\Gamma$ . For this, let  $(M', \sigma') \models T_{\text{Aut}}^*$  such that  $\langle \mathbf{b}_i, \mathbf{c}_i : i \in I \rangle \mathbf{a}$  satisfy  $\Gamma$  in  $M'$ . Suppose  $M' \prec M''$  and  $M''$  is an  $|M'|^+$ -saturated model of  $T$ . In  $M''$  we can find  $\mathbf{b}, \mathbf{b}', \mathbf{c}$  realizing the average of  $\langle \mathbf{b}_i \mathbf{b}'_i \mathbf{c}_i : i \in I \rangle$  over  $M'$ . Then

$$\begin{aligned} \sigma'(\text{tp}(\mathbf{b}/M')) &= \sigma'(\text{avg}(\langle \mathbf{b}_i : i \in I \rangle / M')) \\ &= \text{avg}(\langle \sigma'(\mathbf{b}_i) : i \in I \rangle / M') \\ &= \text{avg}(\langle \mathbf{c}_i : i \in I \rangle / M') \\ &= (\text{tp}(\mathbf{c}/M')). \end{aligned}$$

(The first and last equalities are by the choice of  $\mathbf{b}, \mathbf{c}$ ; the second holds as  $\sigma'$  is an automorphism, and the third follows from clause 4 in the description of the ultra-product, Claim 5.5.) Now since  $M''$  is  $|M'|^+$ -saturated there is an automorphism  $\sigma''$  of  $M''$  extending  $\sigma'$  and taking  $\mathbf{b}$  to  $\mathbf{c}$ .

As  $(M', \sigma') \models T_{\text{Aut}}^*$ , it is existentially closed. So we can pull  $\mathbf{b}, \mathbf{c}$  down to  $M'$ . Thus,  $(M', \sigma') \models \theta_{\Delta_1}(\mathbf{b}_0, \dots, \mathbf{b}_{r-1}, \mathbf{b}_0, \dots, \mathbf{b}_{r-1}, \mathbf{c}_0, \dots, \mathbf{c}_{r-1}, \mathbf{a})$ . But  $(M', \sigma')$  was an arbitrary model of  $T_{\text{Aut}}^* \cup \Gamma \cup \{\sigma(\mathbf{x}_i) = \mathbf{x}'_i : i \in I\}$ ; so

$$\begin{aligned} T_{\text{Aut}}^* \cup \Gamma \cup \{\sigma(\mathbf{x}_i) = \mathbf{x}'_i : i \in I\} \vdash \\ \theta_{\Delta_1}(\mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z}). \quad \square \end{aligned}$$

By compactness, some finite subset  $\Gamma_0$  of  $\Gamma$  and a finite number of the specifications of  $\sigma$  suffice; let  $\Delta^*$  be the formulas mentioned in  $\Gamma_0$  along with those in  $F(\Delta_1, 3n)$  and  $k$  the number of  $x_i, y_i$  appearing in  $\Gamma_0$  and let  $\Delta_2 = F(\Delta^*, n)$ . Without loss of generality,  $k \geq f(\Delta_1, 3n)$ . Then,  $T_{\text{Aut}}^* \vdash$

$$\begin{aligned} & (\forall \mathbf{x}_0, \dots, \mathbf{x}_{k-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{k-1}, \mathbf{y}_0, \dots, \mathbf{y}_{k-1}) \\ & [(\chi_{\Delta_2, k}(\mathbf{x}_0, \dots, \mathbf{x}_{k-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{k-1}, \mathbf{y}_0, \dots, \mathbf{y}_{k-1}, \mathbf{z}) \\ & \wedge \bigwedge_{i < k} \sigma(\mathbf{x}_i) = \mathbf{x}'_i) \rightarrow \theta_{\Delta_1}(\mathbf{x}_0, \dots, \mathbf{x}_{r-1}, \mathbf{x}'_0, \dots, \mathbf{x}'_{r-1}, \mathbf{y}_0, \dots, \mathbf{y}_{r-1}, \mathbf{z})]. \end{aligned}$$

By item 3 in Claim 5.5, fix a  $\Delta_2$  and  $\Delta_3 = \Delta_3(\Delta_2, m)$  containing  $\Delta_2$  and  $\langle \mathbf{b}_i^{\Delta_2}, \mathbf{b}'_i^{\Delta_2}, \mathbf{c}_i^{\Delta_2} : i < k \rangle \mathbf{a}^{\Delta_2}$  which form an  $(\text{id}, \Delta_1, \Delta_2, \Delta_3, n)$ -obstruction over  $A^{\Delta_2}$  and so that

$$M^{\Delta_2} \models \chi_{\Delta_2, k}(\langle \mathbf{b}_i^{\Delta_2}, \mathbf{b}'_i^{\Delta_2}, \mathbf{c}_i^{\Delta_2} : i < k \rangle \mathbf{a}^{\Delta_2}).$$

So by the choice of  $\Gamma_0$ ,

$$(M^{\Delta_2}, \sigma^{\Delta_2}) \models \theta_{\Delta_1}(\langle \mathbf{b}_i^{\Delta_2}, \mathbf{b}'_i^{\Delta_2}, \mathbf{c}_i^{\Delta_2} : i < r \rangle \mathbf{a}^{\Delta_2}).$$

Now, let  $\mathbf{b}, \mathbf{b}', \mathbf{c} \in M^{\Delta_2}$  witness this sentence; then  $\sigma^{\Delta_2}(\mathbf{b}') = \mathbf{c}$ . Then

$$(M^{\Delta_2}, \sigma^{\Delta_2}) \models \tau_{\Delta_1}(\mathbf{b}, \mathbf{b}', \mathbf{c}, \langle \mathbf{b}_i^{\Delta_2} \mathbf{b}_i^{\Delta_2} \mathbf{c}_i^{\Delta_2} : i < k \rangle \mathbf{a}^{\Delta_2}).$$

By the definition of obstruction,

$$\sigma^{\Delta_2}(\text{tp}_{\Delta_3}(\mathbf{b}'/A^{\Delta_2} \cup \mathbf{a}^{\Delta_2})) \neq \text{tp}_{\Delta_3}(\mathbf{c}/A^{\Delta_2} \cup \mathbf{a}^{\Delta_2}).$$

Since  $\sigma^{\Delta_2}$  fixes  $A^{\Delta_2} \cup \mathbf{a}^{\Delta_2}$  pointwise, this contradicts that  $\sigma^{\Delta_2}(\mathbf{b}') = \mathbf{c}$  and we finish.  $\square$

Finally we have the main result.

**Theorem 5.8** *If  $T$  is a stable theory,  $T_{\text{Aut}}$  has a model companion if and only if  $T$  admits no obstructions.*

**Proof** We showed in Lemma 5.4 that if  $T_{\text{Aut}}$  has a model companion then there is no obstruction. If there is no obstruction, Lemma 4.7 implies  $T$  does not have the finite cover property. By Lemma 5.3 for every formula  $\psi(\mathbf{x}, \mathbf{y}, \mathbf{z})$  there is an  $L_\sigma$ -formula  $\theta_\psi(\mathbf{z})$ —write out condition 3 of Lemma 5.3—which for any  $(M, \sigma) \models T_{\text{Aut}}$  holds of any  $\mathbf{a}$  in  $M$  if and only if there exists  $(N, \sigma)$ ,  $(M, \sigma) \subseteq (N, \sigma) \models T_{\text{Aut}}$  and

$$(N, \sigma) \models (\exists \mathbf{x}\mathbf{y})[\psi(\mathbf{x}, \mathbf{y}, \mathbf{a}) \wedge \sigma(\mathbf{x}) = \mathbf{y}].$$

Thus, the class of existentially closed models of  $T_{\text{Aut}}$  is axiomatized by the sentences  $(\forall \mathbf{z})\theta_\psi(\mathbf{z}) \rightarrow (\exists \mathbf{x}\mathbf{y})[\psi(\mathbf{x}, \mathbf{y}, \mathbf{a}) \wedge \sigma(\mathbf{x}) = \mathbf{y}]$ . (We can restrict to formulas of the form  $\psi(\mathbf{x}, \sigma(\mathbf{x}), \mathbf{a})$  by the standard trick (Kikyo and Pillay [5]; Chatzidakis and Hrushovski [3]).  $\square$

Kikyo and Pillay [5] note that if a strongly minimal theory has the definable multiplicity property then  $T_{\text{Aut}}$  has a model companion. In view of Theorem 5.8, this implies that if  $T$  has the definable multiplicity property, then  $T$  admits no obstructions. Kikyo and Pillay conjecture that for a strongly minimal set, the converse holds: if  $T_{\text{Aut}}$  has a model companion then  $T$  has the definable multiplicity property. They prove this result if  $T$  is a finite cover of a theory with the finite multiplicity property. It would follow in general from a positive answer to the following question.

**Question 5.9** *If the  $\omega$ -stable theory  $T$  with finite rank does not have the definable multiplicity property, must it admit obstructions?*

Pillay has given a direct proof that if a strongly minimal  $T$  has the definable multiplicity property, then  $T$  admits no obstructions. Pillay has provided an insightful reworking of the ideas here in a note which is available on his website [8]. Here is a final question.

**Question 5.10** *Can  $T_{\text{Aut}}$  for an  $\aleph_0$ -categorical stable  $T$  admit obstructions?*

## References

- [1] Baldwin, J. T., *Fundamentals of Stability Theory*, Perspectives in Mathematical Logic, Springer-Verlag, Berlin, 1988. [Zbl 0685.03024](#). [MR 89k:03002](#). 133
- [2] Chatzidakis, Z., and A. Pillay, “Generic structures and simple theories,” *Annals of Pure and Applied Logic*, vol. 95 (1998), pp. 71–92. [Zbl 0929.03043](#). [MR 2000c:03028](#). 129

- [3] Chatzidakis, Z., and U. Hrushovski, “The model theory of difference fields,” *Transactions of AMS*, vol. 351 (1999), pp. 2997–3071. [129](#), [141](#)
- [4] Kikyo, H., “Model companions of theories with an automorphism,” *The Journal of Symbolic Logic*, vol. 65 (2000), pp. 1215–22. [Zbl 0967.03031](#). [MR 2001k:03072](#). [129](#), [139](#)
- [5] Kikyo, H., and A. Pillay, “The definable multiplicity property and generic automorphisms,” *Annals of Pure and Applied Logic*, vol. 106 (2000), pp. 263–73. [Zbl 0967.03030](#). [MR 2001j:03068](#). [141](#)
- [6] Kikyo, H., and S. Shelah, “The strict order property and generic automorphisms,” *The Journal of Symbolic Logic*, vol. 67 (2002), pp. 214–16. [Zbl 1002.03025](#). [MR 2002m:03051](#). [130](#)
- [7] Lascar, D., “Les beaux automorphismes,” *Archive for Mathematical Logic*, vol. 31 (1991), pp. 55–68. [Zbl 0766.03022](#). [MR 92i:03035](#). [129](#), [137](#)
- [8] Pillay, A., “Notes on model companions of stable theories with an automorphism,” <http://www.math.uiuc.edu/People/pillay.html>, 2001. [141](#)
- [9] Shelah, S., *Classification Theory and the Number of Nonisomorphic Models*, 2d edition, Studies in Logic and the Foundations of Mathematics, North-Holland Publishing Co., Amsterdam, 1990. [Zbl 0713.03013](#). [MR 91k:03085](#). [130](#), [131](#), [132](#), [133](#), [136](#)
- [10] Shelah, S., “On model completion of  $T_{\text{aut}}$ ,” preprint, 2003. [130](#)
- [11] Winkler, P., “Model completeness and Skolem expansions,” pp. 408–64 in *Model Theory and Algebra: A Memorial Tribute to Abraham Robinson*, edited by D. H. Saracino and V. B. Weispfenning, Lecture Notes in Mathematics, Springer-Verlag, Berlin, 1975. [Zbl 0307.00005](#). [MR 52:7889](#). [129](#), [139](#)

### Acknowledgments

We thank Assaf Hasson for his comments on the manuscript. We also thank the referee for his insightful remarks and, in particular, for forcing us to make explicit and prove Lemma 3.9 which may have independent interest. This work was partially supported by NSF grant DMS-9803496 and DMS-0100594. This is paper 759 in Shelah’s bibliography. The authors thank Mittag-Leffler Institute, Rutgers University, and the Israel Science Foundation for partial support of this research.

Department of Mathematics, Statistics and Computer Science  
 University of Illinois at Chicago  
 851 S Morgan St (M/C 249)  
 Chicago IL 60607-7045  
[jbaldwin@uic.edu](mailto:jbaldwin@uic.edu)

Einstein Institute of Mathematics  
 The Hebrew University of Jerusalem  
 ISRAEL  
 and  
 Department of Mathematics  
 Hill Center-Busch Campus  
 Rutgers, The State University of New Jersey  
 Piscataway NJ 08854-8019  
[shelah@math.huji.ac.il](mailto:shelah@math.huji.ac.il)