

Infinitary Logics and A.E.C.

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Abstract

We prove here that every a.e.c. with LST number $\leq \kappa$ and vocabulary τ of cardinality $\leq \kappa$ can be defined in the logic $\mathbb{L}_{\beth_2(\kappa)^+, \kappa^+}(\tau)$. In this logic an a.e.c. is therefore an EC class unlike in the Presentation Theorem. Furthermore, we study a connection between the sentences defining an a.e.c. and the infinitary logic L_κ^1 .

Introduction

The Presentation Theorem [5] is central to the development of stability for abstract elementary classes: notably, it enables Ehrenfeucht-Mostowski techniques for classes that have large enough models. This has as almost immediate consequence the link between categoricity and stability, and constitutes the beginning of relatively advanced classification/stability theory in that wider setting.

The Presentation Theorem provides a way to capture an a.e.c. as a PC-class: by expanding its vocabulary of the AEC with infinitely many function symbols, an a.e.c. may be axiomatized by an infinitary formula. Although for the stability-theoretical applications mentioned this expansion is harmless, the question as to whether it is possible to axiomatize an a.e.c. with a (necessarily infinitary) sentence *in the same vocabulary* of the a.e.c. has been asked in various ways in the past. Here we provide a positive solution: given an a.e.c \mathcal{K} we provide an infinitary sentence *in the same original vocabulary* $\varphi_{\lambda+2,0}$ whose models are exactly those in \mathcal{K} . Therefore, unlike the situation

in the Presentation Theorem, here the class turns out to be an EC Class, not a PC class.

The main idea is that a “canonical tree of models”, each of size the LST-number of the class, the tree of height ω ends up providing enough tools; the sentence essentially describes all possible maps from elements of this tree into arbitrary potential models in the class. A combinatorial device (a partition theorem theorem on well-founded trees due to Komjath and the first author [3]) is necessary for our proof.

We prove the following two theorems:

Theorem (Theorem 2.1). *Let $\lambda = \beth_2(|\tau| + \kappa)^+$, where $\kappa = \text{LST}(\mathcal{K})$. Then there is a sentence $\psi_{\mathcal{K}}$ in the logic $\mathbb{L}_{\lambda^+, \kappa^+}(\tau)$ such that $\mathcal{K} = \text{Mod}(\psi_{\mathcal{K}})$.*

and

Theorem (Here, a reduced version of Theorem 3.1). *If $M_1 \subseteq M_2$ are $\tau = \tau_{\mathcal{K}}$ -structures, then TFAE:*

- $M_1 \prec_{\mathcal{K}} M_2$
- if $\bar{a} \in {}^{\kappa}M_1$ then there are \bar{b} , N and f such that
 1. $\bar{b} \in {}^{\kappa}M_1$ and $N \in \mathcal{M}_1$
 2. $\text{Rang}(\bar{a}) \subseteq \text{Rang}(\bar{b})$
 3. f is an isomorphism from N onto $M_1 \upharpoonright \text{Rang}(\bar{b})$
 4. $M_2 \models \varphi_{N, \lambda+1, 1}[\langle f(a_{\alpha}^*) \mid \alpha < \kappa \rangle]$.

1 Canonical trees for a.e.c.’s

Fix \mathcal{K} for the remainder of this paper an a.e.c. with vocabulary $\tau = \tau(\mathcal{K})$ and $\text{LST}(\mathcal{K}) = \kappa$. Let λ be the cardinal $\beth_2(\kappa + |\tau|)^+$.

Without loss of generality we may assume that all models in \mathcal{K} are of cardinality $\geq \kappa$. Furthermore, we will use an “empty model” called M_{empt} with the property that $M_{\text{empt}} \prec_{\mathcal{K}} M$ for all $M \in \mathcal{K}$.

Notation 1.1. *We fix the following notation, models and elements in the rest of this paper.*

- We first fix a sequence of (different) elements $\langle a_{\alpha}^* \mid \alpha < \kappa \times \omega \rangle$ in some model in \mathcal{K} .
- Given a model M , we denote by $\bar{\alpha}_M$ a sequence of ordinals $\langle \alpha_k[M] \mid k < n \rangle$ for some $n < \omega$, where for each $k < n$, $\alpha_k[M] < \kappa$.

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- $S_{\bar{\alpha}[M]} := \bigcup_{\kappa < n} [\kappa \times \kappa, \kappa \times \kappa + \alpha_\kappa[M]]$.

We now define the **canonical tree** of \mathcal{K} :

- $\mathcal{M}_n := \{M \in \mathcal{K} \mid \text{for some } \bar{\alpha} = \bar{\alpha}_M \text{ of length } n, M \text{ has universe } \{\alpha_\alpha^* \mid \alpha \in S_{\bar{\alpha}[M]}\} \text{ and } m < n \Rightarrow M \upharpoonright S_{\bar{\alpha} \upharpoonright m[M]} \prec_{\mathcal{K}} M\}$ (and $\mathcal{M}_0 = \{M_{\text{empt}}\}$),
- $\mathcal{M} = \mathcal{M}_{\mathcal{K}} := \bigcup_n \mathcal{M}_n$; this is a tree with ω levels under $\prec_{\mathcal{K}}$ (equivalently under \subseteq).

And some further notation for the rest of the proof:

- Notation 1.2.**
1. $\bar{x}_n := \langle x_\alpha \mid \alpha < \kappa \times n \rangle$,
 2. $\bar{x}_{=n} := \langle x_\alpha \mid \alpha \in [\kappa \times n, \kappa \times (n+1)] \rangle$.

We now define by induction on $\gamma < \lambda^+$ formulas

$$\varphi_{M,\gamma,n}(\bar{x}_n),$$

for every n and $M \in \mathcal{M}_n$ (when $n = 0$ we may omit M).

Case 1 : $\gamma = 0$

If $n = 0$ then the formula $\varphi_{0,0}$ is \top (the sentence denoting “truth”).

Assume $n > 0$. Then

$$\varphi_{M,0,n} := \bigwedge \text{Diag}_{\kappa}^n(M),$$

where $\text{Diag}_{\kappa}^n(M)$ is the set $\{\varphi(x_{\alpha_0}, \dots, x_{\alpha_{n-1}}) \mid \alpha_0, \dots, \alpha_{n-1} \in S_{\bar{\alpha}[M]}, \varphi(y_0, \dots, y_{n-1}) \text{ is an atomic or a negation of an atomic formula and } M \models \varphi(\alpha_{\alpha_0}^*, \dots, \alpha_{\alpha_{n-1}}^*)\}$.

Case 2 : γ a limit ordinal

Then

$$\varphi_{M,\gamma,n}(\bar{x}_n) := \bigwedge_{\beta < \gamma} \varphi_{M,\beta,n}(\bar{x}_n).$$

Case 3 : $\gamma = \beta + 1$

Let $\varphi_{M,\gamma,n}(\bar{x}_n)$ be the formula

$$\forall \bar{z}_{[\kappa]} \bigvee_{\substack{N \succ_{\mathcal{K}} M \\ N \in \mathcal{M}_{n+1}}} \exists \bar{x}_{=n} \left[\varphi_{N,\beta,n+1}(\bar{x}_{n+1}) \wedge \bigwedge_{\alpha < \alpha_n[N]} \bigvee_{\delta \in S[N]} z_\alpha = x_\delta \right]$$

Note: all the formulas constructed belong to $\mathbb{L}_{\lambda^+, \kappa^+}(\tau)$. When $n = 0$ our formulas are really *sentences* $\varphi_{\gamma,0}$, for $\gamma < \lambda^+$. These sentences may be understood as “external approximations” to the a.e.c. \mathcal{K} . Our first aim is to prove how these approximations end up characterizing the a.e.c. \mathcal{K} .

2 Characterizing \mathcal{K} by its canonical sentence

In this section we prove the first main theorem:

Theorem 2.1. *There is a sentence $\psi_{\mathcal{K}}$ in the logic $\mathbb{L}_{\lambda^+, \kappa^+}(\tau)$ such that $\mathcal{K} = \text{Mod}(\psi_{\mathcal{K}})$.*

Our first aim in this section is to prove that every model $M \in \mathcal{K}$ satisfies $\varphi_{\gamma, 0}$, for all $\gamma < \lambda^+$.

In order to achieve this, we prove the following (more elaborate) statement, by induction on γ .

Claim 2.2. *If $M \in \mathcal{K}$, $n < \omega$, $N \in \mathcal{M}_n$, $f : N \rightarrow M$ is a $\prec_{\mathcal{K}}$ -embedding (if $n = 0$, f is empty) then $M \models \varphi_{N, \gamma, n}[\langle f(a_{\alpha}^* \mid \alpha < \kappa \times n) \rangle]$.*

PROOF Let first $\gamma = 0$. Then we have either $n = 0$ in which case trivially $M \models \varphi_{0, 0}(= \top)$ or $n > 0$. In the latter case $\varphi_{N, 0, n} := \bigwedge \text{Diag}_{\kappa}^n(N)$; if $f : N \rightarrow M$ is a $\prec_{\mathcal{K}}$ -embedding, M satisfies this sentence as it satisfies each of the formulas $\varphi(y_0, \dots, y_{n-1})$ satisfied in N by the images of the $\prec_{\mathcal{K}}$ -map f .

The case γ limit ordinal is an immediate consequence of the induction hypothesis.

Let now $\gamma = \beta + 1$ and assume that for every $M \in \mathcal{K}$, $n < \omega$, $N \in \mathcal{M}_n$, if $f : N \rightarrow M$ is a $\prec_{\mathcal{K}}$ -embedding then $M \models \varphi_{N, \beta, n}[\langle f(a_{\alpha}^* \mid \alpha < \kappa \times n) \rangle]$. Now, fix $M \in \mathcal{K}$, $n < \omega$, $N \in \mathcal{M}_n$ and $f : N \rightarrow M$ a \mathcal{K} -embedding. We want to check that $M \models \varphi_{N, \gamma, n}[\langle f(a_{\alpha}^* \mid \alpha < \kappa \times n) \rangle]$, i.e. we need to verify that

$$M \models \forall \bar{z}_{[\kappa]} \bigvee_{\substack{N' \succ_{\mathcal{K}} N \\ N' \in \mathcal{M}_{n+1}}} \exists \bar{x}_{=n} \left[\varphi_{N', \beta, n+1}(\bar{x}_n \widehat{\sim} \bar{x}_{=n}) \wedge \bigwedge_{\alpha < \alpha_{n+1}[N']} \bigvee_{\delta \in S[N']} z_{\alpha} = x_{\delta} \right]$$

when \bar{x}_n is replaced in M by $\langle f(a_{\alpha}^* \mid \alpha < \kappa \times n) \rangle$.

So let $\bar{c}_{[\kappa]} \in M$. By the LST axiom, there is some $M' \prec_{\mathcal{K}} M$ containing both $\bar{c}_{[\kappa]}$ and $\langle f(a_{\alpha}^* \mid \alpha < \kappa \times n) \rangle$, with $|M'| = \kappa$. By the isomorphism axioms there is $N' \succ_{\mathcal{K}} N$, $N' \in \mathcal{M}_{n+1}$, isomorphic to M' through an isomorphism f' extending f . We may now apply the induction hypothesis to N' , f' : since $f' : N' \rightarrow M$ is a $\prec_{\mathcal{K}}$ -embedding, we have that $M \models \varphi_{N', \beta, n+1}[\langle a_{\alpha}^* \mid \alpha < \kappa \times (n+1) \rangle]$. But this enables us to conclude: N' is a witness in the disjunction on models $\prec_{\mathcal{K}}$ -extending N , and the existential $\exists \bar{x}_{=n}$ is witnessed by $\langle a_{\alpha}^* \mid \alpha \in [\kappa \times n, \kappa \times (n+1)] \rangle$. As the original M' had been chosen to include the sequence $\bar{c}_{[\kappa]}$, the last part of the formula holds. □Claim 2.2

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Now we come to the main point:

Claim 2.3. *If M is a τ -model and $M \models \varphi_{\lambda+2,0}$ then $M \in \mathcal{K}$.*

PROOF Let $\mathcal{N} := \{N_* \subseteq M \mid N_* \text{ has cardinality } \kappa \text{ and for some } N \in \mathcal{M}_1 \text{ there is a bijective } f : \kappa \rightarrow N_* \text{ such that } M \models \varphi_{N,\lambda,1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle]\}$. In particular, the previous f is an isomorphism from N to N_* .

We prove first

$$N_1^* \subseteq N_2^* \quad (N_\ell^* \in \mathcal{N}) \quad \text{then } N_1^* \prec_{\mathcal{K}} N_2^*. \quad (1)$$

To see this, choose $(N_\eta^\ell, f_\eta^\ell)$ for $\ell = 1, 2$ and $\eta \in \text{ds}(\lambda) := \{\nu \mid \nu \text{ a decreasing sequence of ordinals } < \lambda\}$ by induction on $\ell g(\eta)$ such that

1. $N_\eta^\ell \in \mathcal{M}_{\ell g(\eta)+1}$
2. f_η^ℓ embeds N_η^ℓ into M : $f_\eta^\ell(N_\eta^\ell) \subseteq M$
3. $M \models \varphi_{N_\eta^\ell, \text{last}(\eta), \ell g(\eta)+1}[\langle f_\eta^\ell(a_\alpha^*) \mid \alpha < \kappa \times (\ell g(\eta) + 1) \rangle]$ where $\text{last}(\langle \rangle) = \lambda$, $\text{last}(\nu \frown \langle \alpha \rangle) = \alpha$
4. if $\nu \triangleleft \eta$ then $N_\nu^\ell \prec_{\mathcal{K}} N_\eta^\ell$ and $f_\nu^\ell \subseteq f_\eta^\ell$
5. if $\ell g(\eta) = 0$ then $f_\eta^\ell(N_\eta^\ell) = N_\ell^*$
6. $N_\eta^1 \subseteq N_\eta^2$ and $\nu \triangleleft \eta \Rightarrow N_\nu^2 \subseteq N_\eta^1$.

Carrying the induction is standard: if $\ell g(\eta) = 0$ let f_η^ℓ be a one-to-one function from $a_\alpha^* \mid \alpha < \kappa$ onto N_ℓ^* ; as $\|N_\ell^*\| = \kappa$ there is a model N_η^ℓ with universe $a_\alpha^* \mid \alpha < \kappa$ such that f_η^ℓ is an isomorphism from N_η^ℓ onto N_ℓ^* .

If $\ell g(\eta) = n = m + 1$ we first choose (f_η^1, N_η^1) . We use

for $\ell = 1$ and from the inductive definition of $\varphi_{N_{\eta \upharpoonright m, \text{last}(\eta \upharpoonright m)}, m}^1$ with $\bar{z}_{[\kappa]} \mapsto \langle f_{\eta \upharpoonright m}^2(a_\alpha^*) \mid \alpha < \kappa \times m \rangle$ (using renaming), the $\bar{x}_{=n}$ gives us the map f_η^1 , with domain N_η^1

Now to choose (f_η^2, N_η^2) we use the inductive definition of $\varphi_{N_{\eta \upharpoonright m, \text{last}(\eta \upharpoonright m)}, m}^2$ with $\bar{z}_{[\kappa]} \mapsto \langle f_{\eta \upharpoonright n}^1(a_\alpha^*) \mid \alpha < \kappa \times n \rangle$; as before, the sequence $\bar{x}_{=n}$ gives us the map f_η^2 , with domain N_η^2 .

Let us now check why having carried the induction suffices.

We apply the partition theorem on well founded trees due to Komjath and the first author [3]. In [2], Gruenhut and the first author provide the following useful form.

Theorem 2.4 (Komjath-Shelah, [3]). *Let α be an ordinal and μ a cardinal. Set $\lambda = \left(|\alpha|^{\mu^{\aleph_0}}\right)^+$ and let $F(\text{ds}(\lambda^+)) \rightarrow \mu$ be a colouring of the tree of finite*

descending sequences of ordinals $< \lambda$. Then there are an embedding $\varphi : \text{ds}(\alpha) \rightarrow \text{ds}(\lambda)$ and a function $c : \omega \rightarrow \mu$ such that for every $\eta \in \text{ds}(\alpha)$ of length $n + 1$

$$F(\varphi(\eta)) = c(n).$$

We apply it with number of colours μ equal to $\kappa^{|\tau|+\kappa} = 2^\kappa$; therefore $(2^\kappa)^{\aleph_0} = 2^\kappa$. We thus obtain a sequence $(\eta_n)_{n < \omega}$, $\eta_n \in \text{ds}(\lambda)$ such that:

$$k \leq m \leq n, \ell \in \{1, 2\} \Rightarrow N_{\eta_m \upharpoonright k}^\ell = N_{\eta_n \upharpoonright k}^\ell.$$

We therefore obtain $(N_k^\ell, g_{k,n}^\ell)_{k \leq n}$ such that

- $N_k^1 \subseteq N_k^2 \subseteq N_{k+1}^1$
- $g_{k,n}^\ell$ is an isomorphism from N_k^ℓ onto $N_{\eta_n \upharpoonright k}^\ell$
- $g_{k,n}^1 \subseteq g_{k,n}^2 \subseteq g_{k+1,n}^1$.

Hence $N_n^\ell \prec_{\mathcal{K}} N_{n+1}^\ell$ and so $\langle N_n^\ell \mid n < \omega \rangle$ is $\prec_{\mathcal{K}}$ -increasing. Let $N_\ell := \bigcup_n N_n^\ell$. Then clearly $N_1 = N_2$; call this model N . Since we then have $N_n^1 \prec_{\mathcal{K}} N$, $N_n^2 \prec_{\mathcal{K}} N$ and $N_n^1 \subseteq N_n^2$ by the coherence axiom for A.E.C.s we have that $N_n^1 \prec_{\mathcal{K}} N_n^2$. In particular, when $n = 0$ we get that $N_1^* \prec_{\mathcal{K}} N_2^*$.

Finally, we also have that

$$N \text{ is cofinal in } [M]^{\leq \kappa}. \quad (2)$$

This is true, since $M \models \varphi_{\lambda+2,0}$

Finally, putting together (1) and (2), we conclude that every τ -model M such that $M \models \varphi_{\lambda+2,0}$ must be in the class (use the union axiom for a.e.c.'s).

□_{Lemma 2.3}

Lastly, we complete the proof of Theorem 2.1: Claims 2.2 and 2.3 provide the definability in the class, as clearly $\varphi_{\gamma,0} \in \mathbb{L}_{\lambda^+, \kappa^+}(\tau_{\mathcal{K}})$.

□_{Theorem 2.1}

3 Strong embeddings and definability

We now focus on the relation $\prec_{\mathcal{K}}$ of our a.e.c. \mathcal{K} : we characterize it in $\mathbb{L}_{\lambda^+, \kappa^+}$. We prove an analog of a ‘‘Tarski-Vaught’’ criterion for a.e.c.’s.

Theorem 3.1. *Let \mathcal{K} be an a.e.c., $\tau = \tau(\mathcal{K})$, $\kappa = \text{LST}(\mathcal{K})$, $\lambda = \beth_2(\kappa + |\tau|)$. Then, given τ -models $M_1 \subseteq M_2$, the following are equivalent:*

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- (A) $M_1 \prec_{\mathcal{K}} M_2$
- (B) if $\bar{a}_\ell \in \kappa^{\geq}(M_\ell)$ for $\ell = 1, 2$ and $\gamma < \lambda$ then there are \bar{b}_ℓ, N_ℓ and f_ℓ for $\ell = 1, 2$ such that:
- for $\ell = 1, 2$,
 - (a) $\bar{b}_\ell \in \kappa^{\geq}(M_\ell)$ and $N_\ell \in \mathcal{M}_\ell$
 - (b) $\text{Rang}(\bar{a}_\ell) \subseteq \text{Rang}(\bar{b}_\ell)$
 - (c) f_ℓ is an isomorphism from N_ℓ onto $M_\ell \upharpoonright \text{Rang}(\bar{b}_\ell)$
 - (d) $\text{Rang}(\bar{b}_1) \subseteq \text{Rang}(\bar{b}_2)$
 - (e) $N_1 \subseteq N_2$
 - (f) $M_\ell \models \varphi_{N_\ell, \gamma, \ell}[\langle f_\ell(a_\alpha^*) \mid \alpha < \kappa \ell \rangle]$.
- (C) if $\bar{a} \in \kappa^{\geq}(M_1)$ then there are \bar{b}, N and f such that
- (a) $\bar{b} \in \kappa^{\geq}(M_1)$ and $N \in \mathcal{M}_1$
 - (b) $\text{Rang}(\bar{a}) \subseteq \text{Rang}(\bar{b})$
 - (c) f is an isomorphism from N onto $M_1 \upharpoonright \text{Rang}(\bar{b})$
 - (d) $M_2 \models \varphi_{N, \lambda+1, 1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle]$.

PROOF (A) \Rightarrow (B): Let $\bar{a}_\ell \in \kappa^{\geq}(M_\ell)$ for $\ell = 1, 2$ and let $\gamma < \lambda$. Choose first $N_1^* \prec_{\mathcal{K}} M_1$ of cardinality $\leq \kappa$ including $\text{Rang}(\bar{a}_1)$ and next, choose $N_2^* \prec_{\mathcal{K}} M_2$ including $N_1 \cup \bar{a}_2$. Let \bar{b}_ℓ enumerate N_ℓ^* and let (N_1, f_1, N_2, f_2) be such that

1. $N_1 \in \mathcal{M}_1, N_2 \in \mathcal{M}_2, N_1 \subseteq N_2$ and
2. f_ℓ is an isomorphism from N_ℓ onto N_ℓ^* for $\ell = 1, 2$.

This is possible: since $M_1 \prec_{\mathcal{K}} M_2$ and $N_\ell^* \prec_{\mathcal{K}} M_\ell$ for $\ell = 1, 2$, we also have that $N_1^* \prec_{\mathcal{K}} N_2^*$. Therefore there are corresponding models $N_1 \subseteq N_2$ in the canonical tree, at levels 1 and 2 (as these must satisfy $N_1 \prec_{\mathcal{K}} N_2$).

We then have that $f_\ell : N_\ell \rightarrow M_\ell$ is a \mathcal{K} -embedding from elements N_1 and N_2 in the canonical tree \mathcal{M} . By Claim 2.2, we may conclude that

$$M_1 \models \varphi_{N_1, \gamma, 1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle]$$

and

$$M_2 \models \varphi_{N_2, \gamma, 2}[\langle f(a_\alpha^*) \mid \alpha < \kappa \times 2 \rangle],$$

for each $\gamma < \kappa$.

(B) \Rightarrow (C): let $\bar{a} \in \kappa^{\geq}(M_1)$. We need $\bar{b}, N \in \mathcal{M}_1$ and $f : N \rightarrow M_1 \upharpoonright \text{Rang}(\bar{b})$ such that

$$M_2 \models \varphi_{N, \lambda+1, 1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle]. \quad (3)$$

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(B) provides a model $N = N_1 \in \mathcal{M}_1$ and elements $\bar{b} = \bar{b}_1$, as well as an isomorphism $f : N \rightarrow \text{Rang}(\bar{b})$. We now check that (B) also implies 3.

Recall the definition of $\varphi_{N,\lambda+1,1}$ (as applied to $[\langle f(a_\alpha^* \mid \alpha < \kappa) \rangle]$). This formula holds in M_2 if for every $\bar{c}_{[\kappa]}$ (of size κ) in M_2 , **for some** $\prec_{\mathcal{K}}$ -extension N' of N in M_2 we have that

$$M_2 \models \exists \bar{x}_{=2} \varphi_{N',\lambda,2}[\langle f(a_\alpha^* \mid \alpha < \kappa) \rangle \bar{x}_{=2}] \quad (4)$$

and the elements $\bar{c}_{[\kappa]}$ are “covered” by the list of elements (of length $\kappa \times 2$) $\langle f(a_\alpha^* \mid \alpha < \kappa) \rangle \bar{x}_{=2}$. But the remaining part of clause (B) provides just this: there is *some* $N' = N_2 \in \mathcal{M}_2$, extending $N = N_1$ such that for each $\gamma < \lambda$, and an isomorphism f' from N' into some $\prec_{\mathcal{K}}$ -submodel N^* of M_2 containing $\text{Rang}(\bar{c}_{[\kappa]})$ such that $M_2 \models \varphi_{N',\gamma,2}[\langle f'(a_\alpha^* \mid \alpha < \kappa \times 2) \rangle]$. The submodel N' witnesses the disjunction on models and $\langle f'(a_\alpha^*) \mid \alpha \in [\kappa, \kappa \times 2] \rangle$ witnesses the existential $\bar{x}_{=2}$.

(C) \Rightarrow (A): assuming (C) means that for every κ -tuple \bar{a} from M_1 there are a model $N \in \mathcal{M}_1$, a κ -tuple \bar{b} from M_1 containing \bar{a} and an isomorphism from N onto $M_1 \upharpoonright \text{Rang}(\bar{b})$ such that

$$M_2 \models \varphi_{N,\lambda+1,1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle].$$

This means that for each \bar{c} included in M_2 (of length κ) there are some extension N' of N with $N' \in \mathcal{M}_2$ and some \bar{d} included in M_2 , of length κ , such that

$$M_2 \models \varphi_{N',\lambda,2}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle \bar{d}]$$

and such that $\text{Rang}(\bar{c}) \subseteq \text{Rang}([\langle f(a_\alpha^*) \rangle] \bar{d})$.

So let, for each $\bar{a} \in {}^{\kappa}M_1$,

- $N_{\bar{a}}^* \prec_{\mathcal{K}} M_1$ be such that $|N_{\bar{a}}^*| = \kappa$,
- $\text{Rang}(\bar{a}) \subseteq N_{\bar{a}}^*$.

Furthermore let for each such \bar{a} ,

$$\mathcal{N}_{\bar{a}} := \{N^* \subseteq M_2 \mid N_{\bar{a}} \prec_{\mathcal{K}} N^* \text{ and } \|N^*\| = \kappa\}.$$

Then for each \bar{a} , $\mathcal{N}_{\bar{a}}$ is cofinal among subsets of M_2 of cardinality κ . By the union axiom, $N_{\bar{a}} \prec_{\mathcal{K}} \bigcup \mathcal{N}_{\bar{a}} = M_2$.

On the other hand, allowing \bar{a} to range over $[M_1]^{\leq \kappa}$, we have a system \mathcal{N}^0 of \subseteq -submodels of M_1 whose union is M_1 (again by (C)). Using a well-founded tree argument as in the proof of 1, this is really a cofinal $\prec_{\mathcal{K}}$ system, and therefore for each \bar{a} , $N_{\bar{a}} \prec_{\mathcal{K}} \bigcup \mathcal{N}^0 = M_1$ and we may conclude that $M_1 \prec_{\mathcal{K}} M_2$.

□ Theorem 3.1

4 Around the logic of an a.e.c.

The logic L_κ^1 from Shelah's paper [6] satisfies Interpolation and a weak form of compactness: strong undefinability of well-order. Furthermore, it satisfies a Lindström-like maximality theorem for these properties (as well as union of ω -chains of models). The logic L_κ^1 , however, lacks a well-defined syntax. Väänänen and Villaveces [7] have produced a logic with a clearly defined (and relatively simple) syntax, whose Δ -closure (a notion appearing first in [4]) is L_κ^1 , and which satisfies several of the good properties of that logic (of course, strong undefinability of well-order but also closure under unions of chains). Also, Dzamonja and Väänänen have linked chain logic [1] to L_κ^1 .

All of these logics are close to our constructions in this paper: the sentence $\varphi_{\lambda+2,0}$ belongs to $\mathbb{L}_{\lambda^+, \kappa^+}$ and L_μ^1 lies in between two logics of the form L_{μ, \aleph_0} and $L_{\mu, \mu}$. Our sentence $\varphi_{\lambda+2,0}$ belongs to L_μ^1 . However, it is not clear if this is the minimal logic for which this is the case.

The question of which is the internal logic of an a.e.c. remains still partially open.

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