STABLE THEORIES AND REPRESENTATION OVER SETS [COSH:919]

MORAN COHEN AND SAHARON SHELAH

ABSTRACT. In this paper we give characterizations of the stable and \aleph_0 -stable theories, in terms of an external property called representation. In the sense of the representation property, the mentioned classes of first-order theories can be regarded as "not very complicated".

§0. Introduction

0.A. **Motivation and main results.** Our motivation to investigate the properties under consideration in this paper comes from the following

Thesis: It is very interesting to find dividing lines and it is a fruitful approach in investigating quite general classes of models. A "natural" dividing property "should" have equivalent internal, syntactical, and external properties. (see [4] for more)

Of course, we expect the natural dividing lines will have many equivalent definitions by internal and external properties.

The class of stable (complete first order theories) T is well known (see [8]), it has many equivalent definitions by "internal, syntactical" properties, such as the order property. As for external properties, one may say "for every $\lambda \geq |T|$ for some model M of T we have $\mathbf{S}(M)$ has cardinality $> \lambda$ " is such a property (characterizing instability). Anyhow, the property "not having many κ -resplendent models (or equivalently, having at most one in each cardinality)" is certainly such an external property (see [5]).

Here we deal with another external property, representability. The results are phrased below, and the full definition appears in Definition 1.2, but first consider a simplified version. We say that a a model M is \mathfrak{k} -representable for a class \mathfrak{k} when there exists a structure $\mathbf{I} \in \mathfrak{k}$ with the same universe as M such that for any n and two sequences of length n from M, if they realize the same quantifier free type in \mathbf{I}

 $^{1991\ \}textit{Mathematics Subject Classification.}\ \ 03\text{C}45 (Primary),\ 03\text{C}55.$

Key words and phrases. model theory, classification theory, stability, representation.

The author would like to thank the ISF for partially supporting this research. This article is [CoSh:919] in the First author's list of publications. The authors thank Assaf Hasson for his remarks on early versions of this paper.

then they realize the same (first order) type in M. Of course, T is \mathfrak{k} -representable if every model of T is \mathfrak{k} -representable. We prove, e.g. that T is stable iff it is $\mathfrak{k}_{\kappa}^{\mathrm{unary}}$ -representable for some κ where $\mathfrak{k}_{\kappa}^{\mathrm{unary}}$ is the class of structures with exactly κ unary functions (and nothing else).

There are some relatives characterizing "T is \aleph_0 -stable, T is superstable and $\kappa(T) \leq \kappa, |T| = \mu$ ".

Originally we hoped to characterize "T is dependent", using classes which have in addition a linear order. However this would have as application to the existence of enough indiscernible sequences. This was proved for strongly dependent T, in [6], but for a general dependent theory by Kaplan-Shelah [2], it is false.

Still this raises various further questions

Problem:

- (1) Can we characterize, by representability "T is strongly dependent", similarly for the various relatives (see [Sh:863])
- (2) For a natural number n, what is the class of T representable by $\mathfrak{t}_{\kappa}^{n}$ of structures with just κ n-place functions (or relations).

The main results presented in this paper are:

Characterization of stable theories (Theorem 2.1):

For a complete first-order theory T, the following conditions are equivalent

- (1) T is stable
- (2) T is representable in $\operatorname{Ex}^1_{|T|^+,|T|}(\mathfrak{k}^{\operatorname{eq}})$ (cf. definitions 1.2, 1.19, and 1.5).
- (3) For some cardinals $\mu_1, \kappa_1, \mu_2, \kappa_2$, it holds that T is representable in $\operatorname{Ex}^1_{\mu_1,\kappa_1}(\operatorname{Ex}^2_{\mu_2,\kappa_2}(\mathfrak{k}^{\operatorname{eq}}))$ (cf. definition 1.20).
- (4) T is representable in $\operatorname{Ex}_{\mu,\kappa}^{0}(\mathfrak{t}^{\operatorname{eq}})$ for some cardinals μ,κ (cf. definition 1.17).
- (5) T is representable in $\operatorname{Ex}_{0,|T|}^{0}(\mathfrak{k}^{\operatorname{eq}})$.

Characterization of \aleph_0 -stable theories (Theorem 3.1):

For a complete first-order theory T, the following conditions are equivalent

- (1) T is \aleph_0 -stable.
- (2) T is representable in $\operatorname{Ex}_{\aleph_0,\aleph_0}^2(\mathfrak{k}^{\operatorname{eq}})$.
- (3) T is representable in $\operatorname{Ex}_{\aleph_0,2}^1(\mathfrak{k}^{\operatorname{eq}})$.
- (4) T is representable in $\text{Ex}_{\aleph_0,2}^{0,\text{lf}}(\mathfrak{k}^{\text{eq}})$ (see definition 1.18)

In the attempt to extend the framework of representation it seemed natural, initially, to conjecture that if we consider representation over linear orders rather than over sets, we could find an analogous characterizations for dependent theories. However, such characterizations would imply strong theorems on existence of

indiscernible sequences. Lately (see [2]), some dependent theories were discovered for which it is provably "quite hard to find indiscernible subsequences", implying that this conjecture would fail in its original formulation.

0.B. Set-theoretic Preliminaries. We use the greek letters μ, κ, λ for cardinal numbers. T will be used to denote a first-order theory, \mathfrak{C}_T a monster model for it.

The set-theoretic facts that are used in this article are stated here for the sake of completeness. We also prove the special version of the Δ -system lemma that is used later.

Theorem 0.1. (Fodor) Let λ a regular uncountable cardinal, $f: \lambda \to \lambda$ such that $f(\alpha) < \alpha$ for all $0 < \alpha < \lambda$, Then $\{\alpha < \lambda : f(\alpha) = \beta\}$ is a stationary set of λ for some $\beta < \lambda$. (see [1])

Corollary 0.2. If $f : \lambda \to \mu$, $\lambda > \mu$ regular, then $f^{-1}(\{\alpha\})$ is a stationary set of λ for some $\alpha < \mu$.

Theorem 0.3. Let λ be a regular cardinal, $|W| = \lambda, |S_t| < \mu$ (for $t \in W$) such that $\chi < \lambda \to \chi^{<\mu} < \lambda$. Then:

- (1) (The Δ -system lemma) There exist $W' \subseteq W$, $|W'| = \lambda$ and S such that $s \neq t \to S_t \cap S_s = S$ holds for all $s, t \in W'$.
- (2) If $\bar{z}_t = \langle z_t^{\alpha} : \alpha < \alpha(t) \rangle$ enumerates S_t for $t \in W$, then we can add:
 - (a) $t \in W' \Rightarrow \alpha(t) = \alpha_0$ holds for some α_0 .
 - (b) For some $U \subseteq \alpha_0$ it holds that $s \neq t \in W' \Rightarrow \bar{z}_t \upharpoonright U = \bar{z}_s \upharpoonright U$, $U = \{\alpha < \alpha_0 : z_t^{\alpha} = z_s^{\alpha}\}.$
 - (c) For some equivalence relation E on α_0 it holds that $t \in W' \Rightarrow z_t^{\alpha} = z_t^{\beta} \leftrightarrow (\alpha, \beta) \in E$.

Proof.

- (1) See [1, 8].
- (2) The map $t \to \alpha(t)$ fulfills the assumptions of Theorem 0.1 $(\alpha(t) < \mu < \lambda)$, therefore (a) holds for some $W_0 \subseteq W$. By part 1 there exist $S \subseteq \{z_t^\alpha: \alpha < \alpha_0, t \in W_0\}$, $W_1 \subseteq W_0$ of cardinality λ , such that $S = S_t \cap S_s$ for all $t \neq s$. Define a map $W_1 \ni t \mapsto U_t$ where: $U_t = \{\alpha < \alpha_0 : z_t^\alpha \in S\}$, Since the range has cardinality $2^{|\alpha_0|} \le 2^{<\mu} < \lambda$ this map also fulfills the assumptions of Theorem 0.1, and we get that for some $W_2 \subseteq W_1$ of cardinality λ and U it holds that $t \in W_2 \to U_t = U$. The range of the map $t \to \bar{z}_t \upharpoonright U$ is US whose cardinality is $\leq |\alpha_0|^{|\alpha_0|} < \lambda$, and by another use of Theorem 0.1 we get $W_3 \subseteq W_2$ of cardinality λ such that (b) holds. The map $t \to E_t$ where $E_t = \{(\alpha, \beta) : z_t^\alpha = z_t^\beta, \alpha, \beta < \alpha_0\}$ has cardinality at most $|\alpha_0|^{|\alpha_0|}$ and again by Theorem 0.1 the result holds for some E and $W' \subseteq W_3$ of cardinality λ , now W' is as required.

0.C. Model-theoretic and Stability-theoretic preliminaries. This subsection is organized in three parts - General, Stable and \aleph_0 -Stable theories.

General. For the rest of this paper, T is assumed to be a complete first-order theory. $\mathbf{EC}(T)$ is the elementary class of T, i.e. the class of all models satisfying T. $\mathbf{EC}_{\lambda}(T)$ denotes the class of models of cardinality λ satisfying T. We will use the name "structure" for any triple consisting of: a set (the structure's universe or domain), a vocabulary (i.e. function symbols and relation symbols with prescribed arities), and an interpretation relation. Structures will be denoted by \mathbf{I}, \mathbf{J} and their domains by I, J, respectively. The word "model" will only be used for the elements of $\mathbf{EC}(T)$. Models will be denoted by M, N and their domains by |M|, |N| respectively.

The German letter \mathfrak{k} will denote a class of structures in a fixed vocabulary $\tau_{\mathfrak{k}}$.

Remark 0.4. Although any model is a structure, we shall use the words "structure" and "model" carefully since in this paper we use structures to analyse models of the theory T. Usually we will deal with quantifier-free properties of the structures, but with general first-order properties of the models.

Definition 0.5. (1) For a model M and set $A \subseteq |M|$, $\mathbf{S}^m(A, M)$ denotes all the complete m-types in M over A.

(2) If $M = \mathfrak{C}_T$, we may omit it.

4

Definition 0.6. For a set $A \subseteq \mathfrak{C}_T$, let $\mathbf{FE}(A)$ denote the set of formulas $\varphi(x, y, \overline{a})$ such that $T \models "\varphi(x, y, \overline{a})$ is an equivalence relation with finitely many classes"

Definition 0.7. The formula $\varphi(\overline{x}, \overline{a})$ divides over a set A iff there exists a sequence $\langle \overline{a}_n : n < \omega \rangle$ such that $\mathbf{tp}(\overline{a}_n, A) = \mathbf{tp}(\overline{a}, A)$ for all $n < \omega$, but there exists an $m < \omega$ such that $\models \neg \exists \overline{x} \bigwedge_{n \in w} \varphi(\overline{x}, \overline{a}_n)$ holds for all $w \in [\omega]^m$.

The type $p(\overline{x})$ forks over A if there exist formulas $\varphi_i(\overline{x}, \overline{a}_i)$ (i < n), such that for all i < n, φ_i divides over A and $p(\overline{x}) \vdash \bigvee_{i < n} \varphi_i(\overline{x}, \overline{a}_i)$.

Theorem 0.8 (Forking preserved under elementary maps). If $p(\overline{x})$ forks over A, and f is an elementary map in M, $Dom(f) \supseteq Dom(p) \cup A$, then f(p) forks over f(A). (see [8, III.1.5])

Definition 0.9. (1) We say that a formula $\varphi(x, \overline{c})$ (with parameters from \mathfrak{C}) is almost over $A \subseteq \mathfrak{C}$ iff for some $E(x,y) \in \mathbf{FE}(A)$ and some $d_i \in \mathfrak{C}$, (i < n) it holds that $T \models E(x, d_i) \leftrightarrow \varphi(x, \overline{c})$.

(2) A formula is over $A \subseteq \mathfrak{C}$ iff it is equivalent in T to a formula with parameters taken only from A.

Theorem 0.10. (see [8] III;2.2(2)) There are (up to logical equivalence mod T) at most |T| + |A| formulas almost over A.

Theorem 0.11. (the order property) T is unstable if and only if there exist a formula $\varphi(\overline{x}, \overline{y})$ and a sequence $\langle \overline{a}_n : n < \omega \rangle$ such that $\models \varphi(\overline{a}_i, \overline{a}_j)^{\text{if}(i < j)}$ holds for all $i, j < \omega$. (see [8, II.2.13])

Stable theories.

Definition 0.12. (1) T is called κ -stable iff $|\mathbf{S}^m(A, M)| \leq \kappa$ for every $M \models T$, set $A \subseteq M$, $|A| \leq \kappa$ and $m < \omega$.

(2) T is called stable if T is κ -stable for some κ .

For the rest of this subsection, T is assumed to be stable.

Theorem 0.13. (Type definability in stable theories) For every formula $\varphi(\overline{x}, \overline{y})$ there exists another formula $\psi_{\varphi}(\overline{y}, \overline{z})$ such that: For all $\overline{b} \in \mathfrak{C}$ there exists a \overline{c} such that if $A \subseteq M$, $|A| \ge 2$ then

$$\models \varphi(\overline{b}, \overline{a}) \Leftrightarrow \models \psi_{\varphi}(\overline{a}, \overline{c})$$

for all $\overline{a} \in A$ (see [8, II.2.2]).

Theorem 0.14. For stable T, distinct types $p, q \in \mathbf{S}(B)$ non-forking over $A \subseteq B$, there exists $E \in \mathbf{FE}(A)$ (cf. Definition 0.6) such that

$$p(x) \cup q(y) \vdash \neg E(x, y)$$

(see [8, III.2.9(2)])

Lemma 0.15. For a stable T, $A \subset M \models T$, and $a \in M$, there exists $B \subseteq A$ such that:

- |B| < |T|
- For every $\varphi(\overline{x}, \overline{c})$ over A which is almost over B there exists $\theta(\overline{x}, \overline{d})$ over B such that $\models \forall \overline{x} (\theta(\overline{x}, \overline{d}) \leftrightarrow \varphi(\overline{x}, \overline{c}))$.
- $\mathbf{tp}(a, A)$ does not fork over B.

Proof. First, define an increasing sequence B_n by induction on n.

Let $|B_0| < \kappa(T) \le |T|^+$, $B_0 \subseteq A$ such that $\mathbf{tp}(a, A)$ does not fork over B_0 . Now assume B_n was defined and let

$$S_n := \{ \varphi(\overline{x}, \overline{c}) \in \mathcal{L}_T : \overline{c} \subseteq A, \varphi(\overline{x}, \overline{c}) \text{ is almost over } B_n \}$$

By Theorem 0.10 there exist at most $|T| + |B_n| = |T|$ non equivalent formulas almost over B_n . Therefore w.l.o.g $|S_n| \le |T|$ and define B_{n+1} as follows:

$$B_{n+1} := B_n \cup \{ \overline{c} : \varphi(\overline{x}, \overline{c}) \in S_n \}$$

That the required properties of $B := \bigcup_{n < \omega} B_n$ hold is easily verified.

Definition 0.16. $\kappa(T)$ is the first cardinal such that for any model $M \models T$, increasing sequence of sets $\langle A_i \subseteq |M| : i \le \kappa \rangle$ and type $p(\bar{x}) \in \mathbf{S}^m(A_{\kappa})$ it holds for some $i < \kappa$ that $p \upharpoonright A_{i+1}$ does not fork over A_i .

Theorem 0.17. T and $\kappa(T)$ satisfy the following (still assuming T is stable):

- (1) $\kappa(T) \leq |T|^+$.
- (2) For every $B \subseteq |M|$ and type $p(\bar{x}) \in \mathbf{S}^m(B)$, there exists $A \subseteq B, |A| < \kappa(T)$ such that p does not fork over A.

 \aleph_0 -Stable theories. For the rest of this subsection, T is assumed to be \aleph_0 -stable.

Theorem 0.18. Let $p \in \mathbf{S}(A)$. For a given $B \supseteq A$ there are only finitely many non-forking extensions of p in $\mathbf{S}(B)$. (see [8], III)

Corollary 0.19. For $p \in \mathbf{S}(A)$ there exists a finite $B \subseteq A$ such that p is the unique non-forking extension of $p \upharpoonright B$ to $\mathbf{S}(A)$.

Proof. By 0.18 there are finitely many non-forking extensions of p
vert B in $\mathbf{S}(A)$, therefore there exists a finite $B_0 \subseteq A$ such that $q_0
vert B_0 \neq q_1
vert B_0$ holds for every distinct $q_0, q_1 \in \mathbf{S}(A)$ non-forking extensions of p. Also p does not fork over some finite $B_1 \subseteq A$. Now, the conclusion easily follows for $B = B_0 \cup B_1$.

Theorem 0.20. A formula $\varphi(\overline{x}, \overline{c})$ is equivalent to a formula over B if and only if $\varphi(\overline{x}, f(\overline{c})) \equiv \varphi(\overline{x}, \overline{c})$ holds for every $f \in Aut(\mathfrak{C}/B)$. (see [8, III.2.3(2)])

§1. Structure classes and representation

Convention 1.1. (1) The vocabulary is a set of individual constants, (partial) function symbols and finitary and relation symbols (=predicates), each with the number of places (= arity), so for a function symbol F, arity_{\tau}(F) is the number of places. Individual constants may be considered as 0-place function symbols; here function symbols are interpreted as partial functions. (2) A structure $\mathbf{I} = \langle \tau, I, \models \rangle$ is a triple of vocabulary, universe(domain) and the interpretation relation for the vocabulary: let $|\mathbf{I}| = I$, $||\mathbf{I}||$ the cardinality of I and $\tau_{\mathbf{I}} = \tau$; \mathbf{I} is called a \tau-structure, so $\mathbf{I} \in \mathfrak{k} \Rightarrow \mathbf{I}$ is a \tau_\mathbf{t}-structure. (3) \mathbf{k} denotes a class of structures in a given vocabulary \tau_\mathbf{t}. (4) \mathcal{L}_{qf}^\tau\$ denotes all the quantifier-free formulas with terms from \tau_\mathbf{t}. That is, finite boolean combinations of atomic formulas, where atomic formulas (for \tau) have the form $P(\sigma_0, ..., \sigma_{n-1})$ or \sigma_0 = \sigma_1 for some n-ary predicate $P \in \tau$, \sigma_0 ... are terms in the closure of the variable by function (and partial function) symbols. (5) If \Delta is a set of formulas in the vocabulary \tau, \mathbf{I} a \tau-structure,

 $\bar{a} = \langle a_i : i < \alpha \rangle \in {}^{\alpha} |\mathbf{I}|, then$

$$\mathbf{tp}_{\Delta}\left(\bar{a},B,\mathbf{I}\right) = \left\{\varphi(\bar{x},\bar{b}): \varphi(\bar{x},\bar{y}) \in \Delta: \mathbf{I} \models \varphi(\bar{a},\bar{b}), \ \bar{b} \in {}^{\ell\!\!g(\bar{y})}B\right\}$$

1.A. **Defining representations.** We now reach the central definitions

Definition 1.2. Consider a structure \mathbf{I} and a set of formulas $\Delta \subseteq \mathcal{L}_{\mathbf{I}}$. (in our use here they are either as in 1.1 (4) (for structures \mathbf{I}) or just first order).

(1) For a structure \mathbf{J} , a function $f: |\mathbf{I}| \to |\mathbf{J}|$ is called a Δ -representation of \mathbf{I} in \mathbf{J} iff

$$\mathbf{tp}_{\mathrm{qf}}(f(\overline{a}),\emptyset,\mathbf{J}) = \mathbf{tp}_{\mathrm{qf}}(f(\overline{b}),\emptyset,\mathbf{J}) \ \Rightarrow \ \mathbf{tp}_{\Delta}(\overline{a},\emptyset,\mathbf{I}) = \mathbf{tp}_{\Delta}(\overline{b},\emptyset,\mathbf{I})$$

for any two sequences $\overline{a}, \overline{b} \in {}^{<\omega}I$ of equal length.

- (2) We say that \mathbf{I} is Δ -represented in a class of models \mathfrak{k} if there exists a $\mathbf{J} \in \mathfrak{k}$ such that \mathbf{I} is Δ -represented in \mathbf{J} .
- (3) For two classes of structures $\mathfrak{k}_0, \mathfrak{k}$ we say that \mathfrak{k}_0 is Δ -represented in \mathfrak{k} if every $\mathbf{I} \in \mathfrak{k}_0$ is Δ -represented in \mathfrak{k} .
- (4) We say that a first-order theory T is Δ -represented in \mathfrak{k} if $\mathbf{EC}(T)$ is Δ -represented in \mathfrak{k} . We may omit Δ from the notation for the set of first order formulas or use qf for the set of quantifier-free formulas.
- **Definition 1.3.** (1) For a structure \mathbf{I} , $\Delta \subseteq \mathcal{L}_T$ and sequences $\overline{a}_t(t \in I)$ from \mathfrak{C}_T we say that $\mathbf{a} = \langle \overline{a}_t : t \in I \rangle$ is a Δ -indiscernible sequence over A in M if $\mathbf{tp}_{\Delta}(\overline{a}_{\overline{t}}, A, M) = \mathbf{tp}_{\Delta}(\overline{a}_{\overline{s}}, A, M)$ holds for all $\overline{s}, \overline{t} \subseteq I$ with the same quantifier-free type in \mathbf{I} .
 - (2) For a class of structures \mathfrak{k} , $M \models T$ and subset $A \subseteq |M|$ we denote by $\operatorname{Ind}_{\Delta}(\mathfrak{k}, A, M)$ the class of structures $\mathbf{a} = \langle \overline{a}_t : t \in I \rangle$ ($I \in \mathfrak{k}$, $\overline{a}_t \subset |M|$) which are Δ -indiscernible in M over $A \subseteq |M|$.

Convention 1.4. We omit the respective symbol from the above notation in the specific cases $\Delta = \mathcal{L}(\tau_M)$, $M = \mathfrak{C}$ and $A = \emptyset$.

Definition 1.5. \mathfrak{t}^{eq} denotes the class of structures of the vocabulary $\{=\}$.

1.B. The free algebras $\mathcal{M}_{\mu,\kappa}$.

Definition 1.6. For a given structure \mathbf{I} , we define the structure $\mathcal{M}_{\mu,\kappa}(\mathbf{I})$ as the structure whose vocabulary is $\tau_{\mathbf{I}} \cup \langle F_{\alpha,\beta} : \alpha < \mu, \beta < \kappa \rangle$, with a β -ary function symbol $F_{\alpha,\beta}$ for all $\alpha < \mu, \beta < \kappa$. (the vocabulary of \mathbf{I} includes a unary relation symbol I for the structure's universe, and we will assume $F_{\alpha,\beta} \notin \tau_{\mathbf{I}}$). The universe for this

¹ This definition also appears in [3],[7, II]

 $^{^2}$ for a sequence \overline{s} we use the shorthand $\overline{a}_{\overline{s}}:=a_{s_0} ^\frown \dots ^\frown a_{s_{\ell a(\overline{s})-1}}$

 $structure is^3$:

$$\mathcal{M}_{\mu,\kappa}(\mathbf{I}) = \bigcup_{\gamma \in \mathrm{Ord}} \mathcal{M}_{\mu,\kappa,\gamma}(\mathbf{I})$$

Where $\mathcal{M}_{\zeta} = \mathcal{M}_{\mu,\kappa,\zeta}(\mathbf{I})$ is defined as follows:

- $\mathcal{M}_0(\mathbf{I}) := |\mathbf{I}|$
- For limit ζ : $\mathcal{M}_{\zeta}(\mathbf{I}) = \bigcup_{\xi < \zeta} \mathcal{M}_{\xi}(\mathbf{I})$
- For $\zeta = \gamma + 1$

$$\mathcal{M}_{\zeta} = \mathcal{M}_{\gamma} \cup \left\{ F_{\alpha,\beta}(\overline{b}) : \overline{b} \in {}^{\beta}\mathcal{M}_{\gamma}, \ \alpha < \mu, \ \beta < \kappa \right\}$$

Where $F_{\alpha,\beta}(\overline{b})$ is treated as a formal object. The symbols in $\tau_{\mathbf{I}}$ have the same interpretation as in \mathbf{I} . In particular, α -ary functions may be interpreted as $(\alpha + 1) - ary$ relations. The β -ary function $F_{\alpha,\beta}(\overline{x})$ is interpreted as the mapping $\overline{a} \mapsto F_{\alpha,\beta}(\overline{a})$ for all $\overline{a} \in {}^{\beta} |\mathcal{M}_{\mu,\kappa}(I)|$, where $F_{\alpha,\beta}(\overline{a})$ on the right side of the mapping is the formal object. If $\mu = \kappa = \aleph_0$ we may omit them.

Definition 1.7. We denote $\mathbf{reg}(\kappa) = \begin{cases} \kappa & \kappa = \mathrm{cf}\kappa \\ \kappa^+ & otherwise \end{cases}$ for every cardinal κ .

Remark 1.8. Since $\operatorname{reg}(\kappa) \geq \kappa$ is regular, for all $\beta < \kappa$ and sequence of terms $\sigma_i(\overline{c}_i) \in \mathcal{M}_{\mu,\operatorname{reg}(\kappa)}$, $(i < \beta)$ there exists $\gamma < \operatorname{reg}(\kappa)$ such that $\sigma_i(\overline{c}_i) \in \mathcal{M}_{\gamma}$ for all $i < \beta$. Therefore $F_{\alpha,\beta}(\langle \sigma_i(\overline{c}_i) : i < \beta \rangle) \in \mathcal{M}_{\gamma+1} \subseteq \mathcal{M}_{\mu,\operatorname{reg}(\kappa)}$, hence $\mathcal{M}_{\mu,\kappa}(S) = \mathcal{M}_{\mu,\kappa,\operatorname{reg}(\kappa)}(S)$ and particularly $\mathcal{M}_{\mu,\kappa}(S)$ is a set (though defined as a class).

Observation 1.9. $\|\mathcal{M}_{\mu,\kappa}(S)\| \leq (\mu + |S|)^{\langle \mathbf{reg}(\kappa)}$. $|\mathcal{M}_{\gamma}| \leq (\mu + |S|)^{\langle \mathbf{reg}(\kappa)}$ can be proved by induction on $\gamma \leq \mathbf{reg}(\kappa)$.

Definition 1.10. For a sequence $\overline{a} \subseteq \mathcal{M}_{\mu,\kappa}(S)$ we define its closure under subterms as the set $\operatorname{cl}(\overline{a})$ defined by induction on the construction of the term and the sequence length as

 $\operatorname{cl}(\overline{a}) := \overline{a} \text{ for } \overline{a} \subseteq S. \text{ If } \ell g(\overline{a}) = 1 \text{ and } a_0 = F_{\alpha,\beta}(\overline{b}) \text{ then } \operatorname{cl}(\overline{a}) := \{a_0\} \cup \{\operatorname{cl}(b_i) : i < \beta\}. \text{ Otherwise, } \operatorname{cl}(\overline{a}) := \bigcup \{\operatorname{cl}(a_i) : i < \ell g(\overline{a})\}.$

Observation 1.11. For any set S, and sequence $\overline{a} \subset (\mathcal{M}_{\mu,\kappa}(S))$ of length $< \operatorname{reg}(\kappa)$, the closure of \overline{a} under subterms has cardinality less $< \operatorname{reg}(\kappa)$.

Remark 1.12. If λ is regular, $\lambda \geq \kappa$ then for every set S and sequence $\overline{a} \in {}^{<\lambda}(\mathcal{M}_{\mu,\kappa}(S))$ there exist a subset $S' \subseteq S$ of cardinality $\chi < \lambda$, and a term $\overline{\sigma}(\overline{b}) \in \mathcal{M}_{\mu,\kappa}(S')$ such that $\overline{a} = \overline{\sigma}(\overline{b})$.

Definition 1.13. Denote $\theta_{\mu,\kappa} := (\mathbf{reg}(\kappa) + \mu)^{<\mathbf{reg}(\kappa)}$. (In particular, $\theta_{\mu,\kappa} = \|\mathcal{M}_{\mu,\kappa}(\mathbf{I})\|$ for $\|\mathbf{I}\| = \mathbf{reg}(\kappa)$.)

³This defines a set and not a proper class by remark 1.8.

Definition 1.14. Consider a free algebra $\mathcal{M}(S)$. We shall say that a set \mathcal{A} of $\tau_{\mathcal{M}}$ -terms, is a minimal system of terms for \mathcal{M} , if and only if for every term $\sigma(\overline{v}) \in \mathcal{M}(S)$ there exists a single $\sigma'(\overline{x}) \in \mathcal{A}$ such that for some $\overline{u} \in S$ without repetitions it holds that $\sigma(\overline{v}) = \sigma'(\overline{u})$.

Observation 1.15. Every free algebra has a minimal system of terms. This follows from the axiom of choice.

1.C. Extensions of classes of structures.

Discussion 1.16. For a class of structures \mathfrak{k} , we define several classes of structures that are based on \mathfrak{k} .

Definition 1.17. $\operatorname{Ex}_{\mu,\kappa}^{0}(\mathfrak{k})$ is the class of structures \mathbf{I}^{+} which, for some $\mathbf{I} \in \mathfrak{k}$ satisfy $|\mathbf{I}^{+}| = |\mathbf{I}|; \tau_{\mathbf{I}^{+}} = \tau_{\mathbf{I}} \cup \{P_{\alpha} : \alpha < \mu\} \cup \{F_{\beta} : \beta < \kappa\}$ for new unary relation symbols P_{α} and new unary function symbols F_{β} ; such that if $\mu > 0$ then $\langle P_{\alpha}^{\mathbf{I}^{+}} : \alpha < \mu \rangle$ is a partition of $|\mathbf{I}|$; and, $\langle F_{\beta}^{\mathbf{I}^{+}} : \beta < \kappa \rangle$ are **partial** unary functions.

Definition 1.18. $\operatorname{Ex}_{\mu,\kappa}^{0,\operatorname{lf}}(\mathfrak{k})$ is the class of structures in $\operatorname{Ex}_{\mu,\kappa}^{0}(\mathfrak{k})$ for which the closure of every element under the new functions is finite. (If stands for "locally finite").

Definition 1.19. $\operatorname{Ex}_{\mu,\kappa}^1(\mathfrak{k})$ is the class of structures in $\operatorname{Ex}_{\mu,\kappa}^0(\mathfrak{k})$ for which $F_{\beta}(P_{\alpha}) \subseteq P_{<\alpha} := \bigcup_{\gamma < \alpha} P_{\gamma}$ holds for every $\alpha < \mu, \beta < \kappa$.

Definition 1.20. $\operatorname{Ex}_{\mu,\kappa}^2(\mathfrak{k})$ is the class of structures of the form $\mathbf{I}^+ = \mathcal{M}_{\mu,\kappa}(\mathbf{I})$, for some $\mathbf{I} \in \mathfrak{k}$ (cf. Definition 1.6).

Convention 1.21. $\operatorname{Ex}_{\mu,\kappa}$ will be one of the above classes.

1.D. Some properties of representation and extension classes. Let us note several properties of representation

Observation 1.22. Let I_1, I_2, I_3 be structures. If $f: I_1 \to I_2$ is a Δ -representation of I_1 in I_2 and $g: I_2 \to I_3$ is an $\mathcal{L}_{I_2}^{qf}$ -representation of I_2 in I_3 , then $g \circ f$ is a Δ -representation of I_1 in I_3 .

Observation 1.23. $\operatorname{Ex}_{\mu,\kappa}^{0,\operatorname{lf}}(\mathfrak{k}), \operatorname{Ex}_{\mu,\kappa}^{1}(\mathfrak{k}) \subseteq \operatorname{Ex}_{\mu,\kappa}^{0}(\mathfrak{k}).$

Observation 1.24. For κ finite, $\operatorname{Ex}_{\mu,\kappa}^{0,\operatorname{lf}}(\mathfrak{k}) \supseteq \operatorname{Ex}_{\mu,\kappa}^{1}(\mathfrak{k})$. Let $\mathbf{I} \in \operatorname{Ex}_{\mu,\kappa}^{1}(\mathfrak{k})$ and fix $a \in |\mathbf{I}|$. To see this, consider the tree formed by finite sequences $\bar{\eta} \in {}^{<\omega}[\kappa]$ for which $(F_{\eta_{\ell-1}} \circ \ldots \circ F_{\eta_0})(a)$ is well-defined (i.e. $(F_{\eta_s} \circ \ldots \circ F_{\eta_0})(a) \in \operatorname{Dom} F_{\eta_{s+1}}$ for every $s < \ell - 1$, where $\ell = \ell g \bar{\eta}$). Now, since it holds that $F_{\beta}(P_{\alpha}) \subseteq P_{<\alpha}$ for any $\beta < \kappa, \alpha < \mu$, and by the well-ordering of the ordinals each branch of this tree is finite. König's Lemma implies that the tree is finite, and since there is a map from

the tree onto the closure of a under the functions, we see that \mathbf{I} is actually locally finite.

Observation 1.25. \mathfrak{k} is af-representable in all the extension classes of \mathfrak{k} defined above, for any two cardinals μ, κ .

Observation 1.26. The classes $\operatorname{Ex}_{\mu_1+\mu_2,\kappa_1+\kappa_2}(\mathfrak{k})$ and $\operatorname{Ex}_{\mu_2,\kappa_2}(\operatorname{Ex}_{\mu_1,\kappa_1}(\mathfrak{k}))$ are qf-representable in each-other (for $\operatorname{Ex} = \operatorname{Ex}^0$, $\operatorname{Ex}^{0,\operatorname{lf}}$, Ex^1 , Ex^2).

Observation 1.27. If $\mu_2 \leq \mu_1$, $\kappa_2 \leq \kappa_1$ then $\operatorname{Ex}_{\mu_2,\kappa_2}(\mathfrak{k})$ is qf-representable in $\operatorname{Ex}_{\mu_1,\kappa_1}(\mathfrak{k})$ (Ex = Ex⁰, Ex^{0,lf}, Ex¹, Ex²).

Observation 1.28. $\operatorname{Ex}^2_{\mu_2,\kappa_2}(\operatorname{Ex}^1_{\mu_1,\kappa_1}(\mathfrak{k}))$ is qf-representable in $\operatorname{Ex}^1_{\mu_1,\kappa_1}(\operatorname{Ex}^2_{\mu_2,\kappa_2}(\mathfrak{k}))$, for any $\mu_1,\mu_2,\kappa_2,\kappa_1 \geq \aleph_0$.

Proof. Indeed, Note that $\operatorname{Ex}_{\mu_1,\kappa_1}^1(\cdot)$ expands each structure of the class $\mathfrak k$ by preserving the structure and enriching it with a partition and partial functions between classes of the partition (particularly such that the order of the classes is preserved). On the other hand, $\operatorname{Ex}_{\mu_2,\kappa_2}^2(\cdot)$ extends the structure by adding formal terms. Now, consider structures $\mathbf I \in \mathfrak k$, $\mathbf I^+ \in \operatorname{Ex}_{\mu_1,\kappa_1}^1(\mathfrak k)$, $\mathbf I^{++} \in \operatorname{Ex}_{\mu_2,\kappa_2}^2(\operatorname{Ex}_{\mu_1,\kappa_1}^1(\mathfrak k))$, such that $\mathbf I^{++}$ is formed from $\mathbf I^+$ which is formed from $\mathbf I$ in the obvious way, using the extension classes we defined. In order to show that $\mathbf I^{++}$ is qf-representable in the latter class, we extend $\mathbf I^+$ to a new structure $\mathbf I^*$ by defining the partition P'_{α} as follows: we will take the partition $\{P_{\alpha}\}$ (of $|\mathbf I^+|$ only) and extend it to a partition of $|\mathbf I^{++}|$. Define (as sets) $P'_n = P_{n-1}$ for all $2 \le n < \omega$, $P'_{\alpha} = P_{\alpha}$ for all $\omega \le \alpha$, and $P'_0 = |\mathbf I^{++}| \setminus |\mathbf I^+|$. Note that we did not extend the domain of the structure, nor have we changed any of the partial or full functions. It should be clear that the identity map $\mathbf I^{++} \to \mathbf I^*$ is a qf-representation as required.

Observation 1.29. $\operatorname{Ex}_{2^{\kappa},\kappa}^{1}(\mathfrak{k}^{\operatorname{eq}})$ is *qf-representable in* $\operatorname{Ex}_{0,\kappa}^{1}(\mathfrak{k}^{\operatorname{eq}})$.

Proof. Let $\langle I, \langle P_\alpha : \alpha < 2^\kappa \rangle, \langle F_\beta : \beta < \kappa \rangle \rangle$ be the vocabulary of \mathbf{I}^+ . Without loss of generality $\left| P_0^{\mathbf{I}^+} \right| \geq 2$ (every model such that $\left| P_0^{\mathbf{I}^+} \right| = 1$ can be represented in such a model). We select two distinct $t_0, t_1 \in P_0^{\mathbf{I}^+}$ and let $h: 2^\kappa \to \mathcal{P}(\kappa)$ a bijection. Consider the structure $\mathbf{I}' = \langle I, \langle F_\beta : \beta < \kappa \rangle, \langle G_\beta : \beta < \kappa \rangle \rangle$ whose universe is $|\mathbf{I}^+|$, $F_\beta^{\mathbf{I}'} = F_\beta^{\mathbf{I}^+}$ and also define for all $\gamma < \kappa, \ x \in P_\gamma^{\mathbf{I}^+}$,

$$G_{\beta}^{\mathbf{I}'}(x) = \begin{cases} t_0 & \gamma \in h(\beta) \\ t_1 & \gamma \notin h(\beta) \end{cases}$$

it is easy to verify that the identity is a $\mathcal{L}_{\mathbf{I}^+}^{\mathrm{qf}}$ -representation of \mathbf{I}^+ in \mathbf{I}' .

⁴Here, I means that $P_0 \in \tau_{\mathbf{I}'}, P_0^{\mathbf{I}'} = I$

Observation 1.30. $\operatorname{Ex}_{\mu,\aleph_0}^2(\mathfrak{k}^{\operatorname{eq}})$ is qf-representable in $\operatorname{Ex}_{\mu,2}^1(\mathfrak{k}^{\operatorname{eq}})$ for $\mu \geq \aleph_0$; (the 2 here means only two unary functions).

Proof. Consider a structure from $\operatorname{Ex}^2_{\mu,\aleph_0}(\mathfrak{k}^{\operatorname{eq}})$. By the definition, we may assume that the structure is $\mathcal{M}_{\mu,\aleph_0}(S)$ for some $S \in \mathfrak{k}^{\operatorname{eq}}$. Let $\langle \sigma_\alpha(\overline{x}_\alpha) : \alpha < \mu \rangle$ be a minimal system of terms of $\mathcal{M}_{\mu,\aleph_0}(S)$ (see definition 1.14, ω is the upper bound on function symbol arities, μ is the number of functions). W.l.o.g $\ell g(\overline{x}_0) = 1$, $\sigma_0(x_0) = x_0$. We now construct a structure \mathbf{I}^+ whose vocabulary is $\langle I, f_{\operatorname{last}}, f_{\operatorname{head}}, \langle P_\beta : \beta < \mu \rangle \rangle$ and whose universe is

$$I^{\mathbf{I}^{+}} = \{ \langle \alpha, i, s_0 \dots s_i \rangle : \alpha < \mu, \ i < \ell g(\overline{x}_{\alpha}), \ s_0 \dots s_i \in S \}$$

Let $\langle\langle \alpha_{\beta}, i_{\beta} \rangle : \beta < \mu \rangle$ enumerate the pairs $\{\langle \alpha, i \rangle : \alpha < \mu, i < \ell g(\overline{x}_{\alpha})\}$ in increasing lexicographic order. Let $P_{\beta}^{\mathbf{I}^{+}}$ be the set of sequences in $|\mathbf{I}^{+}|$ whose head is $\langle \alpha_{\beta}, i_{\beta} \rangle$, and let

$$f_{\text{last}}^{\mathbf{I}^{+}}(\langle \alpha, i, s_{0} \dots s_{i} \rangle) := \langle 0, 0, s_{i} \rangle$$

$$f_{\text{head}}^{\mathbf{I}^{+}}(\langle \alpha, i, s_{0} \dots s_{i} \rangle) := \langle \alpha, i - 1, s_{0} \dots s_{i-1} \rangle \quad (i > 0)$$

we define a map $h: \mathcal{M}_{\mu,\aleph_0}(S) \to |\mathbf{I}^+|$ as follows:

$$h(\sigma_{\alpha}(\overline{v})) = \langle \alpha, \ell g(\overline{x}_{\alpha}) - 1, v_0, \dots, v_{\ell q(\overline{x}_{\alpha}) - 1} \rangle$$

That h is a qf-representation of $\mathcal{M}_{\mu,\aleph_0}(S)$ in \mathbf{I}^+ is easy to verify.

Definition 1.31. We say that a function f with domain and range contained in a structure \mathbf{I} is a partial automorphism when for every sequence $\overline{a} \in |\mathbf{I}|$ of members of $\mathrm{Dom}(f)$, it holds that $\mathbf{tp}_{\mathrm{cf}}(\overline{a},\emptyset,\mathbf{I}) = \mathbf{tp}_{\mathrm{cf}}(f(\overline{a}),\emptyset,\mathbf{I})$.

§2. Stable theories

The central result for this section is

Theorem 2.1. Let T be a complete first-order theory, then the following conditions are equivalent:

- (1) T is stable
- (2) T is representable in $\operatorname{Ex}_{0,|T|}^1(\mathfrak{t}^{\operatorname{eq}})$ (cf. definitions 1.2, 1.19, and 1.5).
- (3) T is representable in $\operatorname{Ex}^1_{|T|^+,|T|}(\mathfrak{t}^{\operatorname{eq}})$.
- (4) T is representable in $\operatorname{Ex}_{2^{|T|},|T|}^1(\mathfrak{k}^{\operatorname{eq}})$.
- (5) For some cardinals $\mu_1, \kappa_1, \mu_2, \kappa_2$, it holds that T is representable in $\operatorname{Ex}^1_{\mu_1, \kappa_1}(\operatorname{Ex}^2_{\mu_2, \kappa_2}(\mathfrak{k}^{\operatorname{eq}}))$ (cf. definition 1.20).
- (6) T is representable in $\operatorname{Ex}_{\mu,\kappa}^0(\mathfrak{t}^{\operatorname{eq}})$ for some cardinals μ,κ (cf. definition 1.17).
- (7) T is representable in $\operatorname{Ex}_{0,|T|}^{0}(\mathfrak{k}^{\operatorname{eq}})$.

Remark 2.2. We can add in the above (better as this is a condition only on one $EC_{\lambda}(T)$, and similarly for 2,3,4,6,7)

5'. For some $\mu_1, \kappa_1, \mu_2, \kappa_2$, and $\lambda > (\mu_1 + |T|)^{<\kappa_1} + (\mu_2 + |T|)^{<\kappa_2}$, it holds that $\mathbf{EC}_{\lambda}(T)$ is representable in $\mathrm{Ex}^1_{\mu_1,\kappa_1}(\mathrm{Ex}^2_{\mu_2,\kappa_2}(\mathfrak{k}^\mathrm{eq}))$

Proof. Theorem 2.15 below proves $1 \Rightarrow 3$. Observation 1.27 implies $2 \Rightarrow 3 \Rightarrow 4$ immediately. Similarly, $4 \Rightarrow 5$, $7 \Rightarrow 6$ are immediate by properties of representations. $5 \Rightarrow 1$ follows from Theorem 2.5. $4 \Rightarrow 2$ follows from Observations 1.22,1.29. So far we have equivalence of conditions 1-5. Now, $2 \Rightarrow 7$ is immediate since $\operatorname{Ex}_{0,|T|}^0(\mathfrak{k}^{\operatorname{eq}}) \supseteq \operatorname{Ex}_{0,|T|}^1(\mathfrak{k}^{\operatorname{eq}})$. We leave $6 \Rightarrow 1$ without a complete proof, since it is very similar to $5 \Rightarrow 1$.

2.A. Stability of representable theories.

Discussion 2.3. We shall first prove the first direction of the main theorem. Namely, that a theory which is representable in $\operatorname{Ex}^2_{\mu_2,\kappa_2}\left(\operatorname{Ex}^1_{\mu_1,\kappa_1}(\mathfrak{k}^{\operatorname{eq}})\right)$ is stable. The method relies on the combinatorial properties of models of stable theories, particularly that all order indiscernibles are indiscernible sets.

Theorem 2.4. Assume that

- (1) T is representable in $\operatorname{Ex}_{\mu_2,\kappa_2}^2\left(\operatorname{Ex}_{\mu_1,\kappa_1}^1(\mathfrak{k}^{\operatorname{eq}})\right)$, for fixed cardinals $\mu_1,\mu_2,\kappa_1,\kappa_2$.
- (2) $\mu \ge \mathbf{reg}(\kappa_2) + \mu_1^+$.
- (3) $\lambda > \mu + \theta_{\mu_2,\kappa_2} + \kappa_1$ a regular cardinal (see definition 1.13).
- (4) $\lambda > \chi^{<\mu}$ for all cardinals $\chi < \lambda$.

Then, for every sequence $\bar{b} = \langle \bar{b}_{\alpha} : \alpha < \lambda \rangle \subseteq {}^{<\mu}[\mathfrak{C}_T]$ of length $<\mu$ there exists $S \in [\lambda]^{\lambda}$ such that $\langle \bar{b}_{\alpha} : \alpha \in S \rangle$ is an indiscernible set.

Proof. Let $M \models T$ such that $\bar{b}_{\alpha} \in {}^{<\mu}|M|$ for all $\alpha < \lambda$ and assume that $f: M \to \mathbf{I}^+ := (\mathcal{M}_{\mu_2,\kappa_2}(I), P_{\alpha}, F_{\beta})_{\alpha < \mu_1,\beta < \kappa_1}$ is a representation, $I = \bigcup_{\alpha < \mu_1} P_{\alpha}$ where $\langle P_{\alpha} : \alpha < \mu_1 \rangle$ is a partition of $I, F_{\beta} : P_{\alpha} \to \bigcup_{\gamma < \alpha} P_{\gamma}$ (So, $\text{Dom}(F_{\beta}) = I$) and let $\bar{a}_{\alpha} = f(\bar{b}_{\alpha})$ for all $\alpha < \lambda$.

w.l.o.g we can add the following assumptions

- Each \overline{a}_{α} is closed under subterms in $\mathcal{M}_{\mu_2,\kappa_2}(I)$ since $\mu \geq \kappa_2$ is regular, so apply 1.11 for (μ,κ_2) .
- The set $\{F_{\beta}: \beta < \mu_1\}$ is closed under composition, (including the empty composition which is the identity, recall that those are unary functions)
- Each \overline{a}_{α} is closed under the partial functions F_{β} : To find the closure of \overline{a}_{α} under the functions we need to add at most μ_1 elements, so the closure of \overline{a}_{α} is $< \mu$.
- $\ell g \overline{a}_{\alpha} = \xi = |\xi|$ for all $\alpha < \lambda$: since $\lambda = \bigcup_{\xi < \mu} \{\alpha < \lambda : \xi = \ell g \overline{a}_{\alpha}\}$ and $\lambda > \mu$ is regular, and by reordering.

The rest of the proof is by taking subsequences of the original sequence, while preserving the length λ , as follows (in brackets we note the common property of the sought subsequence):

First subsequence (sequences constructed by the same terms): By Remark 1.12, for each $i < \xi, \alpha < \lambda$ there exist terms $\sigma_{\alpha,i}(\overline{x}_{\alpha,i})$ in the language of $\mathcal{M}_{\mu_2,\kappa_2}$ and sequences $\overline{t}_{\alpha,i} \in {}^{\langle \mathbf{reg}(\kappa_2)I}$ such that $a_{\alpha,i} = \sigma_{\alpha,i}(\overline{t}_{\alpha,i})$, and also $|\bigcup {\{\overline{x}_{\alpha,i:\alpha<\lambda,i<\xi}\}}| \leq \mathbf{reg}(\kappa_2)$. since $\lambda > [\theta_{\mu_2,\kappa_2}]^{\xi}$ is regular, there exist $\langle \sigma_i(\overline{x}_i) : i < \xi \rangle \in$, $S_0 \in [\lambda]^{\lambda}$ such that $\langle \sigma_{\alpha,i}(\overline{x}_{\alpha,i}) : i < \xi \rangle = \langle \sigma_i(\overline{x}_i) : i < \xi \rangle$ for all $\alpha \in S_0$.

Second subsequence (the quantifier free type of \overline{a}_{α} relative to the P_{α}): since $(\kappa_1)^{\xi} < \lambda$, there exists a $S_1 \in [S_0]^{\lambda}$ such that the function

$$\alpha \mapsto \{(i,\beta) \in \xi \times \kappa_1 : a^i_\alpha \in P_\beta\}$$

is constant on S_1 (denote this constant as the relation R_1).

Third subsequence - (the quantifier free type of \overline{a}_{α} relative to the F_{α}): since $\xi^{\mu_1+\xi} \leq \xi^{<\mu} < \lambda$, there exists a $S_2 \in [S_1]^{\lambda}$ such that the function

$$\alpha \mapsto \left\{ (\beta, \zeta_0, \zeta_1) : \zeta_0, \zeta_1 < \xi, \ \beta < \mu_1, \ F_\beta(a_\alpha^{\zeta_0}) = a_\alpha^{\zeta_1} \right\}$$

is constant on S_2 (denote this constant as the relation R_2).

Final subsequence: By the Δ -system theorem (0.3(2)) there exist $S_3 \in [S_2]^{\lambda}$, $U \subseteq \xi$, $E \subseteq \xi \times \xi$ such that:

- $\overline{a}_{\alpha} \upharpoonright U = \overline{a}_{\beta} \upharpoonright U$ for all $\alpha, \beta \in S_3$.
- E is an equivalence relation such that for all $\alpha \in S_3$: $a^i_{\alpha} = a^j_{\alpha} \leftrightarrow (i,j) \in E$.
- $a_{\alpha}^{i} = a_{\beta}^{j} \to i, j \in U$ for all $\alpha \neq \beta \in S_3$.

We now show that for any finite $\overline{u}, \overline{v} \subseteq S_3$ of length ℓ without repetition, it holds that $\overline{a}_{\overline{v}}$ and $\overline{a}_{\overline{u}}$ have the same quantifier-free type in \mathbf{I}^+ .

Let $\varphi(\overline{x}_{\ell \times \xi})$ an atomic formula. By symmetry, it suffices to show that $\varphi(\overline{a}_{\overline{u}}) \to \varphi(\overline{a}_{\overline{v}})$.

• Case $\varphi(\overline{x}_{\ell \times \xi}) = "\sigma_1(\overline{x}_{\ell \times \xi}) = \sigma_2(\overline{x}_{\ell \times \xi})"$:

proof is carried by induction on the complexity of the term σ_1 .

- For $\sigma_1(\overline{x}_{\ell \times \xi}) = F_{\alpha,\beta}(\overline{\sigma}_1^*(\overline{x}_{\ell \times \xi}))$ it follows from properties of the free algebra that for some sequence of terms $\overline{\sigma}_2^*(\overline{x}_{\ell \times \xi})$ it holds that $\sigma_2(\overline{x}_{\ell \times \xi}) = F_{\alpha,\beta}(\overline{\sigma}_2^*(\overline{x}_{\ell \times \xi}))$ and also $\sigma_{1,i}^*(\overline{a}_{\overline{u}}) = \sigma_{2,i}^*(\overline{a}_{\overline{u}})$ for all $i < \alpha$. The induction hypothesis implies that $\sigma_{1,i}^*(\overline{a}_{\overline{v}}) = \sigma_{2,i}^*(\overline{a}_{\overline{v}})$ and thus $\sigma_1(\overline{a}_{\overline{v}}) = \sigma_2(\overline{a}_{\overline{v}})$ as required.
- For $\sigma_1(\overline{x}_{\ell \times \xi}) = F_{\alpha_1^*}(\sigma_1^*(\overline{x}_{\ell \times \xi}))$, the validity of $\varphi(\overline{a}_{\overline{u}})$ implies that $\sigma_2(\overline{a}_{\overline{u}}) = \sigma_1(\overline{a}_{\overline{u}}) \in I$. It is easy to verify (by induction on the complexity of the term) that the terms $\sigma_s(s=1,2)$ contains only symbols from $\overline{x}_{\ell \times \xi}, F_{\alpha}$ (since $\text{Dom}(F_{\alpha}) \subseteq I$). Now, for a finite sequence of

ordinals $\overline{\alpha}$, denote $F_{\overline{\alpha}} := F_{\alpha_0} \circ \ldots \circ F_{\alpha_{\ell g(\overline{\alpha})}}$, $(F_{\langle \rangle} - \text{the identity})$. It is easy to verify that the term $\sigma_s(\overline{x}_{\ell \times \xi})$ takes the form $F_{\overline{\alpha}_s}(x_{i_s,\zeta_s})$. for some sequence $\overline{\alpha}$.

And the formula φ can be rewritten as:

$$F_{\overline{\alpha}_1}(x_{i_1,\zeta_1}) = F_{\overline{\alpha}_2}(x_{i_2,\zeta_2})$$

Since the family $\langle F_{\alpha}: \alpha < \mu_1 \rangle$ is closed under composition (see above), there exists an $\beta_s < \mu_1$ such that $F_{\overline{\alpha}_s} = F_{\beta_s}$. The sequences $\overline{a}_{u_{i_s}}$ are closed under $\langle F_{\alpha}: \alpha < \mu_1 \rangle$, hence for some $\zeta_s^* < \xi$ it holds that $F_{\beta_s}(a_{u_{i_s},\zeta_s}) = a_{u_{i_s},\zeta_s^*}$ and $a_{u_{i_1},\zeta_1^*} = a_{u_{i_2},\zeta_2^*}$. The former implies $\langle \beta_s,\zeta_s,\zeta_s^* \rangle \in R_2$ and the latter implies that $\zeta_1^*,\zeta_2^* \in U$ and $\langle \zeta_1^*,\zeta_2^* \rangle \in E$. Now, since $\overline{a}_{v_{i_1}} \upharpoonright U = \overline{a}_{v_{i_2}} \upharpoonright U$ it follows that $F_{\beta_s}(a_{v_{i_s},\zeta_s}) = a_{v_{i_s},\zeta_s^*}$ and $a_{v_{i_1},\zeta_1^*} = a_{v_{i_2},\zeta_2^*}$ so easily $\models \varphi(\overline{a}_{\overline{v}})$.

• $\varphi(\overline{x}_{\ell \times \xi}) = P_{\alpha}(\sigma(\overline{x}_{\ell \times \xi}))$: $\models \varphi(\overline{a}_{\overline{v}})$ implies that $\sigma(\overline{x}_{\ell \times \xi}) = F_{\overline{\alpha}}(x_{i,\zeta})$ for some $i < \ell, \zeta < \xi$. Now by the closure of the functions under composition, formula is equivalent to $P_{\alpha}(F_{\beta}(x_{i,\zeta}))$. And for some ζ^* we get that $F_{\beta}(a_{u_i,\zeta}) = a_{u_i,\zeta^*}$ and $P_{\alpha}(a_{u_i,\zeta^*})$ implying $\langle \beta, \zeta, \zeta^* \rangle \in R_2$ and $\langle \alpha, \zeta^* \rangle \in R_1$, respectively. Similar arguments give $\models \varphi(\overline{a}_{\overline{u}})$.

Theorem 2.5. If T is representable in $\operatorname{Ex}_{\mu_2,\kappa_2}^2(\operatorname{Ex}_{\mu_1,\kappa_1}^1(\mathfrak{t}^{\operatorname{eq}}))$, then T is stable.

Proof. Assume towards contradiction that T is unstable. By the order property (Theorem 0.11), and compactness, we can construct a sequence $\langle \overline{a}_i : i < \lambda \rangle$, where

$$\lambda = \beth_2 \left(\mu + \theta_{\mu_2,\kappa_2} + \kappa_1\right)^+, \ \mu = \mathbf{reg}(\kappa_2) + \mu_1^+$$

such that $\models \varphi(\overline{a}_i, \overline{a}_j)^{if(i < j)}$ holds for all $i, j < \lambda$.

Moreover, for a stable T,

Now by the assumptions let $f: M \to \mathbf{I}^+$ be a representation of M in $\operatorname{Ex}^2_{\mu_2,\kappa_2}(\operatorname{Ex}^1_{\mu_1,\kappa_1}(\mathfrak{k}^{\operatorname{eq}}))$. It is easily verified that the conditions in 2.4 hold. Hence, there exists $S \in [\lambda]^{\lambda}$ such that $\{\overline{a}_i : i \in S\}$ is an indiscernible set and particularly $\models \varphi(\overline{a}_i, \overline{a}_j) \leftrightarrow \varphi(\overline{a}_j, \overline{a}_i)$ holds for all $i, j \in S$, contradicting the assumption.

2.B. Stability implies representability. Here we turn to proving the other direction of the main result. We recall several facts about stable theories (from [8, II, III]).

Theorem 2.6. Monotonicity of forking: If $p(\overline{x})$ forks over A and $B \subseteq A$, $q \vdash p$ then q forks over B.

Symmetry: $\mathbf{tp}(\overline{a}, A \cup \overline{b})$ does not fork over A iff $\mathbf{tp}(\overline{b}, A \cup \overline{a})$ does not fork over A.

Transitivity: For sets $A \subseteq B \subseteq C$ such that $\mathbf{tp}(\overline{a}, C)$ does not fork over B, and $\mathbf{tp}(\overline{a}, B)$ does not fork over A it holds that $\mathbf{tp}(\overline{a}, C)$ does not fork over A.

Equivalence of Forking and Dividing: a type p forks over A iff it divides over A.

Definition 2.7. A set $C \subseteq \mathfrak{C}$ will be called strongly independent over A if

* For any $a \in C$, the type $\mathbf{tp}(a, A \cup (C \setminus \{a\}))$ is the unique extension in $\mathbf{S}(A \cup (C \setminus \{a\}))$ of $\mathbf{tp}(a, A)$ which does not fork over A.

Definition 2.8. We say that a sequence $\langle \mathcal{I}_{\alpha} : \alpha < \gamma \rangle$ ($\mathcal{I}_{\alpha} \neq \emptyset$, $\alpha < \gamma$), is a strongly independent decomposition of M of length γ iff for all $\alpha < \gamma$, it holds that \mathcal{I}_{α} is strongly independent over $\mathcal{I}_{<\alpha}$ (in M), and that $|M| = \mathcal{I}_{<\gamma}$, where $\mathcal{I}_{<\alpha} = \bigcup \mathcal{I}_{\beta} : \beta < \alpha$, of course.

Lemma 2.9 (Symmetry of strong independence). Let $a_1, a_2 \in \mathfrak{C}$, $A \supseteq B_1, B_2$ such that $\mathbf{tp}(a_i, A \cup \{a_{3-i}\})$ does not fork over B_i and $\mathbf{tp}(a_i, A)$ is the unique non-forking extension of $\mathbf{tp}(a_i, B_i)$ in $\mathbf{S}(A)$ for i = 1, 2; Then, $(*)_1 \Leftrightarrow (*)_2$ where:

 $(*)_i$ The type $\mathbf{tp}(a_i, B_i)$ has a unique extension to $A \cup \{a_{3-i}\}$ which is nonforking.

Proof. By symmetry it suffices to show that $\neg(*)_2 \Rightarrow \neg(*)_1$.

Assume that $p := \mathbf{tp}(a_2, B_2)$ has two different non-forking extensions $p_1, p_2 \in \mathbf{S}(A \cup \{a_1\})$.

Since both types are complete, there exists a formula $\varphi = \varphi(x, a_1, \overline{c})$ with $\overline{c} \subseteq A$ such that $\varphi \in p_1, \ \neg \varphi \in p_2$. Let b_1, b_2 realize p_1, p_2 , respectively.

So, $\mathbf{tp}(b_i, A) = p_i \upharpoonright A$ is a non-forking extension of p, by uniqueness it follows that $p_1 \upharpoonright A = p_2 \upharpoonright A$. Hence, for i < 2 there exist elementary maps F_i in \mathfrak{C} such that $F_i \upharpoonright A = \mathrm{id}_A$, $F_i(b_i) = a_2$.

Let $q_i \in \mathbf{S}(A \cup \{b_i\})$ be a non-forking extension of $\mathbf{tp}(a_1, B_1)$.

Then $F_i(q_i) \in \mathbf{S}(A \cup \{a_2\})$ is a non-forking extension of $\mathbf{tp}(a_1, B_1)$, (since $F_i \upharpoonright A = \mathrm{id}_A$, and non-forking is preserved under elementary maps).

Now, note that $\models \varphi(b_1, a_1, \overline{c}) \land \neg \varphi(b_2, a_1, \overline{c})$ which implies $\varphi(a_2, x, \overline{c}) \in F_1(q_1)$ and also $\neg \varphi(a_2, x, \overline{c}) \in F_2(q_2)$. This implies that $F_1(q_1)$, $F_2(q_2)$ are distinct extensions of $\mathbf{tp}(a_1, B_1)$, as needed.

Proposition 2.10. If $\langle \mathcal{I}_{\alpha} : \alpha < \gamma \rangle$ is a strongly independent decomposition of M, then every order-preserving refinement of this partition is also a strongly independent decomposition of M.

Remark 2.11. An order-preserving refinement is a partition $\langle \mathcal{J}_{\alpha} : \alpha < \gamma' \rangle$ which refines $\langle \mathcal{I}_{\alpha} : \alpha < \gamma \rangle$ such that for all $\alpha < \beta < \gamma$, $\alpha', \beta' < \gamma'$, $\mathcal{I}_{\alpha} \supseteq \mathcal{J}_{\alpha'} \mathcal{I}_{\beta} \supseteq \mathcal{J}_{\beta'}$ imply $\alpha' < \beta'$.

Proof. Using the basic properties of non-forking

Theorem 2.12. Assume T is stable, and let $A \subset B$ such that for every formula φ over B which is almost over A, φ is equivalent (in T) to a formula over A. If $p,q \in \mathbf{S}(B)$ are distinct and non-forking over A, there exists a $\varphi_*(x,\overline{c})$ over A such that $p \vdash \varphi_*, q \vdash \neg \varphi_*$.

Proof. By 0.14, there exists an equivalence relation $E \in \mathbf{FE}(A)$ such that $p(x) \cup q(y) \vdash \neg E(x, y)$.

Let $\{b_i : i < n(E)\} \subseteq \mathfrak{C}$ enumerate representatives for all the distinct equivalence classes of E and let

$$w := \{i < n(E) : p(x) \cup \{E(x, b_i)\} \text{ is consistent}\}$$

W.l.o.g assume that b_i realizes p for all $i \in w$. Let $\varphi(x) := \bigvee_{i \in w} E(x, b_i)$. It can be easily verified that $p(x) \vdash \varphi(x)$ and similarly, $q(x) \vdash \neg \varphi(x)$. We will show that $\varphi(x)$ is preserved by every $f \in \operatorname{Aut}(\mathfrak{C}/B)$:

Since p is over B and E is a formula over B, they are preserved by f and so:

- $p(x) \cup \{E(x,b_i)\} \Leftrightarrow p(x) \cup \{E(x,f(b_i))\}$ holds for all i < n(E).
- $\neg E(b_i, b_j)$ holds for every $i, j < n(E), i \neq j$ and hence also $\neg E(f(b_i), f(b_j))$.

Hence, f can be regarded as a permutation on $\{b_i/E : i \in w\}$, the equivalence classes of E in \mathfrak{C} :

$$f(\varphi(\mathfrak{C})) = f(\bigcup_{i \in w} b_i/E) = \bigcup_{i \in w} f(b_i)/E = \varphi(\mathfrak{C})$$

Consequently, $\models \varphi(x) \equiv f(\varphi(x))$. Now use Theorem 0.20, which gives the required equivalent formula.

Lemma 2.13. Let $\mu = |T|^+$. For a stable T, for each model $M \models T$ there exists a strongly independent decomposition of length $\gamma \leq \mu$

Remark 2.14. Since our definition of a decomposition does not allow empty sets, we use ordinals $\gamma \leq \mu$. Such a decomposition is trivial for $||M|| \leq \mu$ (e.g. by taking an enumeration of |M| and converting to singletons).

Proof. We choose by induction a sequence $\langle \mathcal{I}_{\alpha} : \alpha < \mu \rangle$ such that \mathcal{I}_{α} is strongly independent over $\mathcal{I}_{<\alpha}$, and is also maximal in $|M| \setminus \mathcal{I}_{<\alpha}$ with respect to this property for every $\alpha < \mu$. Indeed this process may end within at most μ steps, so we define $\gamma = \min (\{\alpha < \mu : \mathcal{I}_{\alpha} = \emptyset\} \cup \{\mu\}).$

Assume towards contradiction that the elements of M were not exhausted after μ iterations, then there exists an $a \in M \setminus \mathcal{I}_{<\mu}$. Recall that for a stable theory $\kappa(T) \leq \mu$ (see [8, III;3.2,3.3]) and so, by the definition of $\kappa(T)$ there exists a set

$$B \subseteq \mathcal{I}_{<\mu}, |B| < \kappa(T) \le \mu$$

such that $p(x) := \mathbf{tp}(a, \mathcal{I}_{<\mu})$ is non-forking over B, and by regularity of μ there exists $\alpha_0(*) < \mu$ such that $\mathcal{I}_{<\alpha_0(*)} \supseteq B$. Now, recalling Theorem 0.10, let

$$\Gamma := \left\{ \varphi(x; \overline{c}) : \varphi(x, \overline{c}) \text{ almost over } B, \ \varphi(x; \overline{y}) \in \mathcal{L}, \ \overline{c} \in {}^{\ell g \overline{y}} \mathcal{I}_{<\mu} \right\},\,$$

and choose $\Gamma_* \subseteq \Gamma$ a minimal set of representatives up to logical equivalence from Γ . By Theorem 0.10 we have $|\Gamma_*| \leq |B| + |T| < \operatorname{cf}(\mu)$, and since μ is regular, there exists $\alpha_1(*) < \mu$ such that $\overline{b} \subseteq \mathcal{I}_{<\alpha_1(*)}$ for all $\varphi(x,\overline{b}) \in \Gamma_*$. Let $\alpha(*) = \max{\{\alpha_0(*), \ \alpha_1(*)\}}$. Let $B' = \bigcup (\overline{b} : \varphi(\overline{x},\overline{b}) \in \Gamma_*) \cup B$ be such that $B' \subseteq \mathcal{I}_{<\alpha(*)}, |B'| \leq |T|$. We will now prove that the type $p \upharpoonright \mathcal{I}_{\leq \alpha(*)}$ is the unique extension in $\mathbf{S}(\mathcal{I}_{\leq \alpha(*)})$ of $p \upharpoonright \mathcal{I}_{<\alpha(*)}$.

Clearly $p \upharpoonright \mathcal{I}_{\leq \alpha(*)}$ extends $p \upharpoonright \mathcal{I}_{<\alpha(*)}$, now assume that $q \in \mathbf{S}(\mathcal{I}_{\leq \alpha(*)})$ extends $p \upharpoonright \mathcal{I}_{<\alpha(*)}$, and does not fork over $\mathcal{I}_{<\alpha(*)}$.

By the transitivity of non-forking, q does not fork over B. Assume towards contradiction that $p \neq q$. By fact 0.14 there exists $E \in \mathbf{FE}(B)$ such that $q(x) \cup p(y) \vdash \neg E(x,y)$, and particularly $q(x) \vdash \neg E(x,a)$.

The formula E(x, a) is almost over B, so by the choice of $\alpha_1(*)$ $(q \vdash \neg E(x, a))$ implies that $\neg E(x, a)$ is logically equivalent to a formula in x over $\mathcal{I}_{<\mu}$, there exist a $\bar{b} \subseteq \mathcal{I}_{<\alpha(*)}$ and $\varphi(x, \bar{b})$ logically equivalent to E(x, a).

Now since E(a, a) holds, we also get $\models \varphi(a, \overline{b})$, and since $\overline{b} \subseteq \mathcal{I}_{<\alpha(*)}$, we get $\varphi(x, \overline{b}) \in \mathbf{tp}(a, \mathcal{I}_{<\alpha(*)}) = q \upharpoonright \mathcal{I}_{<\alpha(*)}$, a contradiction.

So we have proved that $\mathbf{tp}(a, \mathcal{I}_{\leq \alpha(*)})$ is the unique non-forking extension of $\mathbf{tp}(a, \mathcal{I}_{<\alpha(*)})$ in $\mathbf{S}\left(\mathcal{I}_{\leq \alpha(*)} \setminus \{b\}\right)$. Recall from the choice of $\mathcal{I}_{\alpha(*)}$ that for all $b \in \mathcal{I}_{\alpha(*)}$, $\mathbf{tp}(b, \mathcal{I}_{\leq \alpha(*)} \setminus \{b\})$ is the unique extension in $\mathbf{S}\left(\mathcal{I}_{\leq \alpha(*)} \setminus \{a\}\right)$ which does not fork over $\mathcal{I}_{<\alpha(*)}$. Also, Lemma 2.9 implies that $\mathbf{tp}(b, \mathcal{I}_{\leq \alpha(*)} \setminus \{b\} \cup \{a\})$ is the unique non-forking extension of $\mathbf{tp}(b, \mathcal{I}_{<\alpha(*)})$ in $\mathbf{S}(\mathcal{I}_{\leq \alpha(*)} \setminus \{b\} \cup \{a\})$.

From the last two sentences it follows that the condition \circledast above holds for $\mathcal{I}_{\alpha(*)} \cup \{a\}$ over $\mathcal{I}_{<\alpha(*)}$, which contradicts the maximality of $\mathcal{I}_{\alpha(*)}$.

Theorem 2.15. A stable first order theory T is representable in $\operatorname{Ex}^1_{|T|^+,|T|}(\mathfrak{k}^{\operatorname{eq}})$.

Proof. Let $M \models T$. By lemma 2.13 we get a strongly independent decomposition of M: $\langle \mathcal{I}_{\alpha} : \alpha < \gamma \leq |T|^+ \rangle$. By Proposition 2.10 we can assume w.l.o.g that $|\mathcal{I}_1| = |\mathcal{I}_0| = 1$.

Define $\mathbf{I}^+ = \left\langle F_i^*, P_{\alpha}, F_{\varphi(x,\overline{y}),j} : \alpha < \gamma, i < |T|, \varphi(x,\overline{y}) \in \mathcal{L}, j < k_{\varphi(x,\overline{y})} \right\rangle \in \mathfrak{k}^{eq}$ as follows:

 $|\mathbf{I}^+| = |M|, P_{\alpha}^{\mathbf{I}^+} = \mathcal{I}_{\alpha} \text{ for all } \alpha < \gamma.$

By Lemma 0.15 we can find some $B(a) \subseteq \mathcal{I}_{<\alpha}$ for all $a \in \mathcal{I}_{\alpha}$, with the following properties:

(\mathbf{H}) $|B(a)| \leq |T|$; For every formula $\varphi(\bar{x}, \bar{c})$ over $\mathcal{I}_{<\alpha}$ which is almost over B(a), there exists a formula $\theta(\bar{x}, \bar{d})$ over B(a) such that $\models \forall \bar{x}(\theta(\bar{x}, \bar{d}) \leftrightarrow \varphi(\bar{x}, \bar{c}))$; if $\alpha > 0$, then $\mathcal{I}_0 \subseteq B(a)$; and, $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$ does not fork over B(a).

Now, we define the functions $\langle (F_i^*)^{\mathbf{I}^+} : i < |T| \rangle$ on $\text{Dom}(F_i^*) = |M| \setminus (\mathcal{I}_0 \cup \mathcal{I}_1)$ for all i < |T|. Fix $\alpha > 1$, $a \in \mathcal{I}_{\alpha}$. We define $F_i^*(a) = b_i(a)$, (i < |T|) for some enumeration $\langle b_i(a) : i < |T| \rangle$ of B(a), possibly with repetitions.

Now, let ψ_{φ} be a formula which is guaranteed to exist from Theorem 0.13 and define the partial unary functions

$$\left\{ F_{\varphi(x,\overline{y}),j}^{\mathbf{I}^+}(x) : j < k_{\varphi(x,\overline{y})} \right\}$$

as follows: Let $a \in \text{Dom}\left(F^{\mathbf{I}^+}_{\varphi(x,\overline{y}),j}\right) = |M| \setminus (\mathcal{I}_0 \cup \mathcal{I}_1)$, and let $\alpha < \mu$ be such that $a \in \mathcal{I}_{\alpha}$. Since $|\mathcal{I}_{<\alpha}| \geq 2$ for all $2 \leq \alpha < \mu$, it follows from the definition of ψ_{φ} (recall Theorem 0.13) that there exists $\overline{c}_a \in {}^{\ell g(\overline{z}_{\varphi})} \mathcal{I}_{<\alpha}$ such that $\models \varphi[a, \overline{b}] \Leftrightarrow \models \psi_{\varphi}[\overline{b}, \overline{c}_a]$. Now we define $F^{\mathbf{I}^+}_{\varphi(x,\overline{y}),j}(a) := (\overline{c}_a)_j$, for all $j < \ell g(\overline{z}_{\varphi})$. Thus we have defined \mathbf{I}^+ and we define $f: M \to \mathbf{I}^+$ as f(a) = a for all $a \in |M|$. Now, to prove that f is indeed a representation, it follows by Proposition 2.16 that if would suffice to prove that $\mathbf{tp}(h(\overline{a}),\emptyset,M) = \mathbf{tp}(\overline{a},\emptyset,M)$, for every $\overline{a} \subseteq \text{Dom}(h)$ and partial automorphism h of \mathbf{I}^+ with domain and range closed under functions.

Let $D_{\alpha} = \mathcal{I}_{\alpha} \cap \text{Dom}(h)$, $R_{\alpha} = \mathcal{I}_{\alpha} \cap \text{Rang}(h)$. It is easily verified that for $\alpha < \gamma$, $h \upharpoonright D_{\alpha}$ is a partial automorphism of \mathbf{I}^+ from D_{α} onto R_{α} . We will prove by induction on $\alpha < \gamma$ that

 $\boxtimes_{\alpha,n}$: $h(\mathbf{tp}(\overline{a}, D_{<\alpha}, M)) = \mathbf{tp}(h(\overline{a}), R_{<\alpha}, M)$ for all $\overline{a} \in D_{\alpha}$ of length n and don't have repetitions.

holds for all $n < \omega$.

Now, $\boxtimes_{\alpha,n}$ holds for $\alpha < 2$ since by the definition, $|\mathcal{I}_{\alpha}| = 1$. Now let $\alpha \geq 2$ and assume that $\boxtimes_{\beta,n}$ holds for all $n < \omega$, $\beta < \alpha$. We prove by induction on $n < \omega$ that $\boxtimes_{\alpha,n}$ holds.

First, $\boxtimes_{\alpha,1}$: Let $a \in \mathcal{I}_{\alpha}$, $\varphi = \varphi(x,\overline{c})$ a formula over $D_{<\alpha}$. W.l.o.g assume $\models \varphi[a,\overline{c}]$. by the definition of the functions it follows $\models \psi_{\varphi}[\overline{c},F_{\varphi,0}(a)\dots F_{\varphi,\ell g\overline{z}-1}(a)]$. This formula contains only constants from $D_{<\alpha}$, so by the induction hypothesis, $\models \psi_{\varphi}[h(\overline{c}),h(F_{\varphi,0}(a))\dots h(F_{\varphi,\ell g\overline{z}-1}(a))]$ holds. Since h is a partial automorphism (with closed range and domain) of \mathbf{I}^+ , h commutes with the functions on \mathbf{I}^+ so $\models \psi_{\varphi}[h(\overline{c}),F_{\varphi,0}(h(a))\dots F_{\varphi,\ell g\overline{z}-1}(h(a))]$ holds. By the definitions of $F_{\varphi,j}$ ($j < \ell g\overline{z}$), ψ_{φ} we get $M \models \varphi[h(a),h(\overline{c})]$, as needed.

For n > 1 we continue by induction, but first we state the following property of \mathbf{I}^+ (to be proven later)

(•): If $A \subseteq \mathbf{I}^+$ is closed under the partial functions in the vocabulary $\tau_{\mathbf{I}^+}$ then $A \cap \mathcal{I}_{\alpha}$ Is strongly independent over $A \cap \mathcal{I}_{<\alpha}$.

Now, let $\overline{a} \in D_{\alpha}$ of length n and $b \in D_{\alpha} \setminus \overline{a}$. By the induction hypothesis (on n), it follows that $h \upharpoonright (D_{<\alpha} \cup \overline{a})$ is elementary. By (\bullet) , D_{α} is strongly independent over $D_{<\alpha}$. Hence, $\mathbf{tp}(b, D_{\leq \alpha} \setminus \{b\})$ does not fork over $D_{<\alpha}$ and particularly $\mathbf{tp}(b, D_{<\alpha} \cup \overline{a})$ does not fork over $D_{<\alpha}$.

By the induction hypothesis, $h
cup (D_{<\alpha} \cup \overline{a})$ is elementary, and so $q := h(\mathbf{tp}(b, D_{<\alpha} \cup \overline{a}))$ does not fork over $h(D_{<\alpha}) = R_{<\alpha}$. Note that $q \supseteq h(\mathbf{tp}(b, D_{<\alpha}))$ and by $\boxtimes_{\alpha,1}$ (see above) $h(\mathbf{tp}(b, D_{<\alpha})) = \mathbf{tp}(h(b), R_{<\alpha})$ holds. Hence, q extends $\mathbf{tp}(h(b), R_{<\alpha})$ to a type over $R_{<\alpha} \cup \overline{a}$ and does not fork over $R_{<\alpha}$. Therefore there exists an extension $q \subseteq q' \in \mathbf{S}(R_{\leq\alpha} \setminus h(b))$ which does not fork over $R_{<\alpha}$.

Since R_{α} is closed under the partial functions, it follows from (\bullet) above, that R_{α} is strongly independent over $R_{<\alpha}$, meaning that $q' = \mathbf{tp}(h(b), R_{\leq \alpha} \setminus \{h(b)\})$. Now we reduce both types to the domain $R_{<\alpha} \cup h(\overline{a})$ to get

$$\mathbf{tp}(h(a), R_{<\alpha} \cup h(\overline{a})) = h(\mathbf{tp}(b, D_{<\alpha} \cup \overline{a}))$$

and the induction step on n:

$$\mathbf{tp}(h(b \overline{a}), R_{\leq \alpha}) = h(\mathbf{tp}(b \overline{a}, D_{\leq \alpha}))$$

Hence, f is a representation.

Proof of (\bullet) from the proof of Theorem 2.15, above. Let $A_{\alpha} = \mathcal{I}_{\alpha} \cap A$, $a \in A_{\alpha}$, and recall that $B(a) = \{F_i^*(a) : i < |T|\}$. We prove that $\mathbf{tp}(a, A \cap \mathcal{I}_{\leq \alpha} \setminus \{a\})$ is the unique non-forking extension of $\mathbf{tp}(a, A \cap \mathcal{I}_{\leq \alpha})$ in $\mathbf{S}(A \cap \mathcal{I}_{\leq \alpha} \setminus \{a\})$.

Since A is closed under the F_i^* , it follows that $B(a) \subseteq A$, and consequently $B(a) \subseteq A_{<\alpha}$. Also, $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$ does not fork over B(a) (Recall (\maltese) above). By transitivity of non forking $\mathbf{tp}(a, \mathcal{I}_{\leq \alpha} \setminus \{a\})$ which is a non-forking extension of $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$, does not fork over B(a) either. By the definition of B(a), we also get that every formula over $\mathcal{I}_{<\alpha}$ which is almost over B(a) is equivalent to a formula over B(a) (again, see (\maltese)).

Now, by monotonicity of non-forking we get that $\mathbf{tp}(a, A_{\leq \alpha} \setminus \{a\}) \subseteq \mathbf{tp}(a, \mathcal{I}_{\leq \alpha} \setminus \{a\})$ does not fork over $A_{\leq \alpha}$.

To prove uniqueness, let $q_0 \in \mathbf{S}(A_{\leq \alpha} \setminus \{a\})$ a non-forking extension of $\mathbf{tp}(a, A_{<\alpha})$. q_0 has a non-forking extension $q \in \mathbf{S}(\mathcal{I}_{\leq \alpha} \setminus \{a\})$. By transitivity, q does not fork over B(a). Recall that the functions F_i^* are defined so that every formula over $\mathcal{I}_{<\alpha}$ and almost over B(a) is equivalent to a formula over B(a). The types $q \upharpoonright \mathcal{I}_{<\alpha}$, $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$ are both non-forking over B(a). Since q extends $\mathbf{tp}(a, A_{<\alpha}) \supseteq$

 $\mathbf{tp}(a, B(a))$ we get that $q \upharpoonright \mathcal{I}_{<\alpha}$, $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$ (both non-forking over B(a)) agree on all formulas over B(a), and by theorem 2.12 this implies $q \upharpoonright \mathcal{I}_{<\alpha} = \mathbf{tp}(a, \mathcal{I}_{<\alpha})$. Now, since q is a non-forking extension of $\mathbf{tp}(a, \mathcal{I}_{<\alpha})$ and \mathcal{I}_{α} is strongly independent over $\mathcal{I}_{<\alpha}$ we get that $q = \mathbf{tp}(a, \mathcal{I}_{\leq\alpha} \setminus \alpha)$ and so

$$q_0 = q \upharpoonright (A_{\leq \alpha} \setminus \{a\}) = \mathbf{tp}(a, A_{\leq \alpha} \setminus \{a\})$$

as required.

Proposition 2.16. Let $\mathbf{I}^+ \in \operatorname{Ex}_{u,\kappa}^1(\mathfrak{k}^{\operatorname{or}})$ and $f: M \to \mathbf{I}^+$. If

$$h(f(\overline{a})) = f(\overline{b}) \to \mathbf{tp}(\overline{a}, \emptyset, M) = \mathbf{tp}(\overline{b}, \emptyset, M)$$

holds for every partial automorphism h of \mathbf{I}^+ with domain and range which are closed under the partial functions, and sequences $\overline{a}, \overline{b} \in M$. Then f is a representation.

Proof. Let f be as described above. Now assume towards contradiction that f is not a representation. Therefore there exist $\overline{a}, \overline{b} \in M$ which have different types in M such that the map $f(\overline{a}) \mapsto f(\overline{b})$ is a partial automorphism of \mathbf{I}^+ . It is possible to extend this partial automorphism to one with domain and range closed under the partial functions, contrary to the definition of f.

§3. \aleph_0 -stable theories

In this section we will prove the following result:

Theorem 3.1 (Characterization of \aleph_0 -stable theories). for a complete, countable first-order theory T, the following conditions are equivalent

- (1) T is \aleph_0 -stable.
- (2) T is representable in $\operatorname{Ex}_{\aleph_0,\aleph_0}^2(\mathfrak{k}^{\operatorname{eq}})$.
- (3) T is representable in $\operatorname{Ex}^1_{\aleph_0,2}(\mathfrak{k}^{\operatorname{eq}})$.
- (4) T is representable in $\text{Ex}_{\aleph_0,2}^{0,\text{lf}}(\mathfrak{k}^{\text{eq}})$ (see definition 1.18)

Proof. Theorem 3.5 gives $1 \Rightarrow 2$ and Observation 1.30 gives $2 \Rightarrow 3$. $3 \Rightarrow 4$ is immediate since $\operatorname{Ex}^1_{\aleph_0,2}(\mathfrak{k}^{\operatorname{eq}}) \subseteq \operatorname{Ex}^{0,\operatorname{lf}}_{\aleph_0,2}(\mathfrak{k}^{\operatorname{eq}})$ by Observation 1.24. $4 \Rightarrow 1$ is the content of Proposition 3.2.

Proposition 3.2. If T is representable in $\operatorname{Ex}_{\aleph_0,2}^{0,\operatorname{lf}}(\mathfrak{k}^{\operatorname{eq}})$ then T is \aleph_0 -stable.

Proof. To prove \aleph_0 -stability it suffices to show that $|S(B,M)| \leq \aleph_0$ for every model $M \models T$ and countable $B \subseteq |M|$. Now, suppose that $f: M \to \mathbf{I}$ is a representation, then for any $a, b \in M$ such that $\mathbf{tp}(a, B, M) \neq \mathbf{tp}(b, B, M)$, it also holds that $\mathbf{tp}_{qf}(f(a), f(B), \mathbf{I}) \neq \mathbf{tp}_{qf}(f(b), f(B), \mathbf{I})$, so if suffices to show that $|\mathbf{S}_{qf}(\mathbf{I}, A)| \leq \aleph_0$

for every structure $\mathbf{I} \in \operatorname{Ex}_{\aleph_0,2}^{0,\operatorname{lf}}(\mathfrak{k}^{\operatorname{eq}})$, and countable $A \subseteq I = |\mathbf{I}|$. Let \mathbf{I}, A be as above, then

- (1) Without loss of generality A is closed under the functions of I.
- (2) $\mathbf{tp}_{qf}(\overline{a}, \mathbf{I}) \in \mathbf{S}_{qf}(\mathbf{I})$ is determined by formulas of the following types

$$P_{\alpha}(\sigma(b)) \qquad (b \in \overline{a})$$

$$\sigma_{1}(b_{0}) = \sigma_{2}(b_{1}) \quad (b_{0}, b_{1} \in \overline{a})$$

$$\sigma_{1}(b_{0}) = b_{1} \qquad (b_{0} \in \overline{a}, b_{1} \in A)$$

for terms $\sigma, \sigma_1, \sigma_2 \in \tau_{\mathbf{I}}$, and so necessarily unary.

(3) Moreover, since **I** is locally finite, $\mathbf{tp}_{qf}(\overline{a}, \mathbf{I})$ is determined by a finite subset of these formulas.

So, the number of unary types over A is at most $|A|^{<\omega} \leq \aleph_0$

Convention 3.3. We assume for the rest of this section that T is stable in \aleph_0 .

Lemma 3.4. Let $M \models T$, and $\mathcal{I}_0 \subseteq |M|$. There exists a sequence of sets $\langle \mathcal{I}_n : 0 < n < \omega \rangle$ such that

- (1) For all $a \in \mathcal{I}_n$, $n < \omega$ there exists a finite $B_a \subseteq \mathcal{I}_{< n}$ such that $\mathbf{tp}(a, \mathcal{I}_{\le n} \setminus \{a\})$ is the unique non-forking extension of $\mathbf{tp}(a, B_a)$ in $\mathbf{S}(\mathcal{I}_{\le n} \setminus \{a\})$,
- (2) $\mathcal{I}_n \cap \mathcal{I}_{\leq n} = \emptyset$, and
- (3) $\mathcal{I}_{<\omega} = |M|$.

In particular, $\langle \mathcal{I}_n : n < \omega \rangle$ is a strongly independent decomposition of M.

Proof. We choose the sequence \mathcal{I}_n by induction on $n \geq 0$. The case n = 0 is given in the assumptions. For n > 0 assume that we have \mathcal{I}_n and choose \mathcal{I}_{n+1} as a set $\mathcal{I}_{n+1} \subseteq M \setminus \mathcal{I}_{\leq n}$ which is maximal (possibly empty) under the requirements 1,2 above. Now, assume towards contradiction that $a \in M \setminus \mathcal{I}_{<\omega}$. By Corollary 0.19, there exists a finite $B_a \subseteq \mathcal{I}_{<\omega}$ such that $\mathbf{tp}(a, \mathcal{I}_{<\omega})$ is the unique non-forking extension of $\mathbf{tp}(a, B_a)$ in $\mathbf{S}(\mathcal{I}_{<\omega})$. Similarly, for any $b \in \mathcal{I}_n$, let $B_b \subseteq \mathcal{I}_{< n}$ be such that $\mathbf{tp}(b, \mathcal{I}_{< n})$ is the unique non-forking extension of $\mathbf{tp}(b, B_b)$ in $\mathbf{S}(\mathcal{I}_{< n})$. Let $0 \leq n_* < \omega$ be minimal such that $B_a \subseteq \mathcal{I}_{< n_*}$.

Now, for any $b \in \mathbf{S}(\mathcal{I}_{n_*})$ it holds that $\mathbf{tp}(b, B_b)$ has a unique non-forking extension in $\mathbf{S}(\mathcal{I}_{\leq n_*} \setminus \{b\})$. Also, $\mathbf{tp}(a, B_a)$ has a unique non-forking extension in $\mathbf{S}(\mathcal{I}_{\leq n_*})$ (since it has a unique non-forking extension in $\mathbf{S}(\mathcal{I}_{<\omega})$). Now by the symmetry of strong independence (Lemma 2.9) it follows that $\mathcal{I}_{n_*} \cup \{a\}$ is strongly independent over $\mathcal{I}_{< n_*}$, which contradicts the maximality of \mathcal{I}_{n_*} .

Theorem 3.5. Let $M \models T$, $\lambda = ||M||$, \mathcal{I}_0 a set of indiscernibles in M. Then M can be represented in $\mathcal{M}_{\aleph_0\aleph_0}(\mathcal{I}_0 \cup \lambda) \in \operatorname{Ex}^2_{\aleph_0,\aleph_0}(\mathfrak{k}^{\operatorname{eq}})$ by an extension of the identity function on \mathcal{I}_0 .

Proof. Let $\langle \mathcal{I}_n : n < \omega \rangle$ as in Lemma 3.4. Let $g : ||M|| \to \lambda$ a one-to-one function. T is \aleph_0 -stable and so $\mathbf{S}^m(\emptyset)$ is countable for all $m < \omega$. For convenience we use

 22

the symbols $\{F_{p,n}: n < \omega, p \in \mathbf{S}^{<\omega}(\emptyset)\}$ as the function symbols of $\mathcal{M}_{\aleph_0\aleph_0}(\mathcal{I}_0 \cup \lambda)$, such that for each m-type p, $F_{p,n}$ is an m-ary function symbol.

We define an increasing sequence of one-to-one functions $f_i: \mathcal{I}_{\leq i} \to \mathcal{M}_{\aleph_0\aleph_0}(\mathcal{I}_0 \cup \lambda)$ by induction on $n < \omega$:

Define f_0 as the identity on \mathcal{I}_0 .

Assume that f_n was defined and now define $f_{n+1} \supseteq f_n$ as follows. For each $a \in \mathcal{I}_{n+1}$ recall B_a from Lemma 3.4. Let $\overline{c}_a \in {}^{\ell}(\mathcal{I}_{\leq n})$ enumerate B_a . Now define $p \in \mathbf{S}^{\ell+1}(\emptyset)$ and $f_{n+1}(a)$ as follows:

$$p := \mathbf{tp}(a \overline{c}_a, \emptyset, M)$$

$$f_{n+1}(a) := F_{p,n}(f_n(\overline{c}_a), g(a))$$

Let $f = \bigcup_{n < \omega} f_n$. We will use (proof is omitted) an analogue of Proposition 2.16 to show that f is a representation:

Proposition 3.6. Let $f: M \to \mathcal{M}(S)$. If

$$h(f(\overline{a})) = f(\overline{b}) \to \mathbf{tp}(\overline{a}, \emptyset, M) = \mathbf{tp}(\overline{b}, \emptyset, M)$$

holds for all $\overline{a}, \overline{b} \in M$ and every partial automorphism h of $\mathcal{M}(S)$ whose domain and range are closed under subterms, then f is a representation.

First note that $a \in \mathcal{I}_n$ and also $f(a) = F_{p,n}(f(\overline{c}_a), g(a))$, so $p = \mathbf{tp}(a \frown \overline{c}_a, \emptyset, M)$ and $\mathbf{tp}(a, \mathcal{I}_{\leq n} \setminus \{a\})$ is the unique non-forking extension of $\mathbf{tp}(a, \overline{c}_a)$.

We now show that f fulfills the conditions of the Proposition. Let h a partial automorphism of $\mathcal{M}(\mathcal{I}_0 \cup \lambda)$ with domain and range closed under the functions. Fix $n < \omega$ and sequences $\overline{a}, \overline{b} \in \mathcal{I}_{\leq n}$ such that $h(f(\overline{a})) = f(\overline{b})$. Since f is one-to-one, w.l.o.g $\overline{a}, \overline{b}$ are without repetition. We prove that $\mathbf{tp}(\overline{a}, \emptyset, M) = \mathbf{tp}(\overline{b}, \emptyset, M)$ by induction on n:

For n = 0: the Proposition holds since \mathcal{I}_0 is an indiscernible set.

For n = m + 1:: Proof by induction on $\ell = |\overline{a} \cap \mathcal{I}_n| = |\overline{b} \cap \mathcal{I}_n|$ (the latter equality is easy to verify).

For $\ell = 0$: This is the Proposition of the induction hypothesis (on n).

For $\ell = \ell_0 + 1$:: Let $a_0, b_0 \in \mathcal{I}_n$, $\overline{a}_1, \overline{b}_1 \in^{\ell} \mathcal{I}_n$, $\overline{b}_1, \overline{b}_2 \in^{<\omega} \mathcal{I}_{< n}$ such that $h(f(a_0\overline{a}_1\overline{a}_2)) = f(b_0\overline{b}_1\overline{b}_2)$. By the definition there exist $\overline{c}_{a_0}, \overline{c}_{b_0}$ such that $f(a_0) = F_{p,n}(f(\overline{c}_{a_0}), g(a_0)), f(b_0) = F_{p',n}(f(\overline{c}_{b_0}), g(b_0))$ for some sequences and types. Since Dom(h) is closed under subterms we get:

$$F_{p',n}(f(\overline{c}_{b_0}),g(b_0)) = f(b_0) = h(f(a_0)) = h(F_{p,n}(f(\overline{c}_{a_0}),g(a_0))) = F_{p,n}(h(f(\overline{c}_{a_0})),h(g(a_0)))$$

and by the definition of the free algebra p' = p and $h(f(\bar{c}_{a_0})) = f(\bar{c}_{b_0})$. The induction hypothesis implies that the map G defined as

$$G(\overline{a}_1) = \overline{b}_1, \ G(\overline{a}_2) = \overline{b}_2, \ G(\overline{c}_{a_0}) = \overline{c}_{b_0}$$

is elementary. Now, let $q = \mathbf{tp}(a_0, \overline{a}_1 \cup \overline{a}_2 \cup \overline{c}_{a_0})$. Since $\mathbf{tp}(a_0 \overline{c}_{a_0}) = p = p' = \mathbf{tp}(b_0 \overline{c}_{b_0})$ holds, it follows that $G(q) \upharpoonright \overline{c}_{b_0} = \mathbf{tp}(b_0, \overline{c}_{b_0})$. The definition of \mathcal{I}_n implies that $\mathbf{tp}(a_0, \mathcal{I}_{\leq n} \setminus \{a_0\})$ is non-forking over \overline{c}_{a_0} , and so is $\mathbf{tp}(a_0, \overline{a}_1 \cup \overline{a}_2 \cup \overline{c}_{a_0})$. On the other hand, since G is elementary, G(q) does not fork over \overline{c}_{b_0} . Let $\mathbf{S}(\mathcal{I}_{\leq n} \setminus \{b_0\}) \ni q' \supseteq G(q)$ a nonforking extension. Since $\mathbf{tp}(b_0, \mathcal{I}_{\leq n} \setminus \{b_0\})$ is the unique non-forking extension of $\mathbf{tp}(b_0, \overline{c}_{b_0})$, and by transitivity q' is also a non-forking extension, it follows that $q' = \mathbf{tp}(b_0, \mathcal{I}_{\leq n} \setminus \{b_0\})$ and after reduction $(\overline{b}$ is without repetitions, so $b_0 \notin \overline{b}_1$ and $\overline{b}_1 \cup \overline{b}_2 \cup \overline{b}_{a_0} \subseteq \mathcal{I}_{\leq n} \setminus \{b_0\}$):

$$G(q) = q' \upharpoonright \overline{b}_1 \cup \overline{b}_2 \cup \overline{c}_{b_0} = \mathbf{tp}\left(b_0, \overline{b}_1 \cup \overline{b}_2 \cup \overline{b}_{a_0}\right)$$

Hence, $G \cup \{(a_0, b_0)\}$ is elementary and the proof is complete.

§References

- [1] Thomas Jech. Set theory. Academic Press, New York, 1978.
- [2] Itay Kaplan and Saharon Shelah. Dependent examples. preprint.
- [3] Saharon Shelah. General non-structure theory. TBA. math.LO/1011.3576.
- [4] Saharon Shelah. Introduction to: classification theory for abstract elementary class. TBA. math.LO/0903.3428.
- [5] Saharon Shelah. On spectrum of κ -respondent models. *preprint*. A revised version will appear in "Non structure theory".
- [6] Saharon Shelah. Strongly dependent theories. Israel Journal of Mathematics, submitted. math.LO/0504197.
- [7] Saharon Shelah. Universal classes. In Classification theory (Chicago, IL, 1985), volume 1292 of Lecture Notes in Mathematics, pages 264–418. Springer, Berlin, 1987. Proceedings of the USA–Israel Conference on Classification Theory, Chicago, December 1985; ed. Baldwin, J.T.
- [8] Saharon Shelah. Classification theory and the number of nonisomorphic models, volume 92 of Studies in Logic and the Foundations of Mathematics. North-Holland Publishing Co., Amsterdam, xxxiv+705 pp, 1990.

MORAN COHEN AND SAHARON SHELAH

Institute of Mathematics, The Hebrew University of Jerusalem, 91904 Jerusalem, Israel

Institute of Mathematics, The Hebrew University of Jerusalem, 91904 Jerusalem, Israel, and Department of Mathematics, Rutgers University, New Brunswick, NJ 08854, USA

 $Email\ address:\ {\it shelah@math.huji.ac.il} \\ URL:\ {\it http://www.math.rutgers.edu/}{\sim} {\it shelah}$

 24