

# A Note on Extensions of Infinitary Logic\*

Saharon Shelah <sup>†</sup>	Jouko Väänänen <sup>‡</sup>
Institute of Mathematics	Department of Mathematics
Hebrew University	University of Helsinki
Jerusalem, Israel	Helsinki, Finland

September 26, 2020

## Abstract

We show that a strong form of the so called Lindström's Theorem [4] fails to generalize to extensions of  $L_{\kappa\omega}$  and  $L_{\kappa\kappa}$ : For weakly compact  $\kappa$  there is no strongest extension of  $L_{\kappa\omega}$  with the  $(\kappa, \kappa)$ -compactness property and the Löwenheim-Skolem theorem down to  $\kappa$ . With an additional set-theoretic assumption, there is no strongest extension of  $L_{\kappa\kappa}$  with the  $(\kappa, \kappa)$ -compactness property and the Löwenheim-Skolem theorem down to  $< \kappa$ .

By a well-known theorem of Lindström [4], first order logic  $L_{\omega\omega}$  is the strongest logic which satisfies the compactness theorem and the downward Löwenheim-Skolem theorem. For weakly compact  $\kappa$ , the infinitary logic  $L_{\kappa\omega}$  satisfies both the  $(\kappa, \kappa)$ -compactness property and the Löwenheim-Skolem theorem down to  $\kappa$ . In [1] Jon Barwise pointed out that  $L_{\kappa\omega}$  is not maximal with respect to these properties, and asked what is the strongest logic based on a weakly compact cardinal  $\kappa$  which still satisfies the  $(\kappa, \kappa)$ -compactness property and some other natural conditions suggested by  $\kappa$ . We prove (Corollary 5) that for weakly compact  $\kappa$  there is no strongest extension of  $L_{\kappa\omega}$  with

---

\*We are indebted to Lauri Hella, Tapani Hyttinen and Kerkko Luosto for useful suggestions.

<sup>†</sup>Research partially supported by the United States-Israel Binational Science Foundation. Publication number [ShVa:726]

<sup>‡</sup>Research partially supported by grant 40734 of the Academy of Finland.

the  $(\kappa, \kappa)$ -compactness property and the Löwenheim-Skolem theorem down to  $\kappa$ . This shows that there is no extension of  $L_{\kappa\omega}$  which would satisfy the most obvious generalization of Lindström's Theorem. A stronger result (Theorem 11) is proved under an additional assumption.

We use the notation and terminology of [2, Chapter II] as much as possible. We will work with concrete logics such as first order logic  $L_{\omega\omega}$ , infinitary logic  $L_{\kappa\lambda}$  and their extensions  $L_{\omega\omega}(\{Q_i : i \in I\})$  and  $L_{\kappa\lambda}(\{Q_i : i \in I\})$  by generalized quantifiers. Therefore it is not at all critical which definition of a logic one uses as long as these logics are included and some basic closure properties are respected. We use  $\mathcal{L} \leq \mathcal{L}'$  to denote the sublogic relation. Let  $\mathcal{P}$  be a property of logics. A logic  $\mathcal{L}^*$  is *strongest extension of  $\mathcal{L}$  with  $\mathcal{P}$* , if

1.  $\mathcal{L} \leq \mathcal{L}^*$ ,
2.  $\mathcal{L}^*$  has property  $\mathcal{P}$ ,

and whenever  $\mathcal{L}$  is a sublogic  $\mathcal{L}'$  and  $\mathcal{L}'$  has property  $\mathcal{P}$ , then  $\mathcal{L}' \leq \mathcal{L}^*$ .

Let  $\mathcal{L}$  be a logic and  $\kappa$  and  $\lambda$  infinite cardinals.  $\mathcal{L}$  is  $(\kappa, \lambda)$ -compact if for all  $\Phi \subseteq \mathcal{L}$  of power  $\kappa$ , if each subset of  $\Phi$  of cardinality  $< \lambda$  has a model, then  $\Phi$  has a model.  $\mathcal{L}$  is  $\kappa$ -compact if it is  $(\kappa, \omega)$ -compact.  $\mathcal{L}$  is *weakly  $\kappa$ -compact* if  $\mathcal{L}$  is  $(\kappa, \kappa)$ -compact.  $\mathcal{L}$  is *fully compact* if it is  $\kappa$ -compact for all  $\kappa$ .  $\mathcal{L}$  has the *Löwenheim-Skolem property down to  $\kappa$* , denoted by  $LS(\kappa)$  if every  $\phi \in \mathcal{L}$  which has a model, has a model of cardinality  $\leq \kappa$ . If every sentence  $\phi \in \mathcal{L}$  which has a model, has a model of cardinality  $< \kappa$ , we say that  $\mathcal{L}$  satisfies  $LS(< \kappa)$ . The order-type of the well-ordering  $R$  is denoted by  $otp(R)$ .

**Theorem 1** [4] *The logic  $L_{\omega\omega}$  is strongest extension of  $L_{\omega\omega}$  with  $\aleph_0$ -compactness and  $LS(\aleph_0)$ .*

Let  $C$  be a class of cardinals. Let

$$\mathfrak{A} \models Q_C^{\text{cf}} xy\phi(x, y, \vec{z}) \iff otp(\{\langle a, b \rangle : \mathfrak{A} \models \phi(a, b, \vec{c})\}) \in C.$$

By [9],  $L_{\omega\omega}(Q_C^{\text{cf}})$  is always fully compact. For  $C$  an interval we use the notation  $Q_{[\kappa, \lambda]}^{\text{cf}}$  and  $Q_{[\kappa, \lambda]}$ .

**Proposition 2** *There is no strongest  $\kappa$ -compact extension of  $L_{\omega\omega}$ . In fact:*

1. *there are fully compact logics  $\mathcal{L}_n$ ,  $n < \omega$ , such that  $\mathcal{L}_n \leq \mathcal{L}_{n+1}$  for all  $n < \omega$ , but no  $\aleph_0$ -compact logic extends each  $\mathcal{L}_n$ .*

2. *There is an  $\aleph_0$ -compact logic  $\mathcal{L}_1$  and a fully compact logic  $\mathcal{L}_2$  such that no  $\aleph_0$ -compact logic extends both  $\mathcal{L}_1$  and  $\mathcal{L}_2$ .*

**Proof.** Let  $\mathcal{L}_n = L_{\omega\omega}(\{Q_{[\aleph_\omega, \infty]}^{\text{cf}}\} \cup \{Q_{\aleph_l}^{\text{cf}} : l < n\})$ . By [9], each  $\mathcal{L}_n$  is fully compact. Clearly, no  $\aleph_0$ -compact logic can extend each  $\mathcal{L}_n$ .

For the second claim, let  $\mathcal{L}_1$  be the logic  $L_{\omega\omega}(Q_1)$ , where  $Q_1$  is the quantifier “there exists uncountable many” introduced by Mostowski [8]. This logic is  $\aleph_0$ -compact [3], see [2, Chapter IV] for more recent results. Let  $\mathcal{L}_2$  be the logic  $L_{\omega\omega}(Q_B)$ , where  $Q_B$  is the quantifier “there is a branch” introduced by Shelah [10]. More exactly,

$$Q_B x y t u M(x) T(y) (t \leq u)$$

if and only if  $\leq_T$  is a partial order of  $T \subseteq M$  and there are  $D, \leq_D, f$  and  $B$  such that:

1.  $\leq_D$  is a total order of  $D \subseteq M$
2.  $f : \langle T, \leq_T \rangle \rightarrow \langle D, \leq_D \rangle$  is strictly increasing
3.  $\forall s \in D \exists p \in T (f(p) = s)$
4.  $B \subseteq T$  is totally ordered by  $\leq_T$
5.  $\forall b \in B ((p \in T \ \& \ p \leq_T b) \rightarrow p \in B)$
6.  $\forall s \in D \exists b \in B (s \leq_D f(b))$ .

The reader is referred to [10] for a proof of the full compactness of  $\mathcal{L}_2$ .

Suppose there were an  $\aleph_0$ -compact logic  $\mathcal{L}$  containing both  $\mathcal{L}_1$  and  $\mathcal{L}_2$  as a sublogic. It is easy to see that the class of countable well-orders can be expressed as a relativized pseudoelementary class in  $\mathcal{L}$ . This contradicts  $\aleph_0$ -compactness of  $\mathcal{L}$ .  $\square$

Lauri Hella pointed out that by elaborating the proof of claim (2) of the above proposition, we can make  $\mathcal{L}_1$  fully compact. It was proved in [11] that, assuming GCH, there is no strongest extension of  $\mathcal{L}_{\omega\omega}$  which is  $\aleph_0$ -compact. Our proof of (1) of the above proposition essentially occurs in a note, based on a suggestion of Paolo Lipparini, added after Theorem 8 of [11].

**Proposition 3** *Suppose  $\kappa > \aleph_0$ . There is no strongest extension of  $L_{\kappa+\omega}$  with  $LS(\kappa)$*

**Proof.** Let  $\mathcal{L}_1 = L_{\kappa+\omega}(Q_{\aleph_0}^{\text{cf}})$  and  $\mathcal{L}_2 = L_{\kappa+\omega}(Q_{[\aleph_1, \kappa]}^{\text{cf}})$ . By using standard arguments with elementary chains of submodels, it is easy to see that both  $\mathcal{L}_1$  and  $\mathcal{L}_2$  have  $\text{LS}(\kappa)$ , but the consistent sentence

$R$  is a linear order with no last element  $\wedge$

$$\neg Q_{\aleph_0}^{\text{cf}} xyR(x, y) \wedge \neg Q_{[\aleph_1, \kappa]}^{\text{cf}} xyR(x, y)$$

has no models of size  $\leq \kappa$ .  $\square$

It was proved in [11] that there is no strongest extension of  $\mathcal{L}_{\omega\omega}$  with  $\text{LS}(\omega)$ .

**Lemma 4** *Suppose  $\kappa$  is weakly compact. Then  $L_{\kappa\omega}(Q_{\aleph_0}^{\text{cf}})$  and  $\mathcal{L}_{\kappa\omega}(Q_{[\aleph_1, \kappa]}^{\text{cf}})$  are weakly  $\kappa$ -compact. Moreover, if  $\kappa > \omega$ , these logics satisfy  $\text{LS}(\kappa)$ .*

**Proof.** The claim concerning  $\text{LS}(\kappa)$  is proved with a standard elementary chain argument. We prove the weak compactness of  $\mathcal{L}_{\kappa\omega}(Q_{[\aleph_1, \kappa]}^{\text{cf}})$ . The case of  $L_{\kappa\omega}(Q_{\aleph_0}^{\text{cf}})$  is similar, but easier. For this end, suppose  $T$  is a set of sentences of  $\mathcal{L}_{\kappa\omega}(Q_{[\aleph_1, \kappa]}^{\text{cf}})$  and  $|T| = \kappa$ . We may assume  $T \subseteq \kappa$ . If  $\alpha < \kappa$ , then we assume that there is a model  $\mathfrak{M}_\alpha \models T \cap \alpha$ . In view of  $\text{LS}(\kappa)$ , it is not a loss of generality to assume that  $\mathfrak{M}_\alpha = \langle \kappa, R_\alpha \rangle$ , where  $R_\alpha \subseteq \kappa \times \kappa$ . Let  $R(\alpha, \beta, \gamma) \iff R_\alpha(\beta, \gamma)$ . By weak compactness there is a transitive  $M$  of cardinality  $\kappa$  such that

$$\langle H(\kappa), \epsilon, T, R \rangle \prec_{L_{\kappa\kappa}} \langle M, \epsilon, T^*, R^* \rangle$$

and  $\kappa \in M$ . Let  $\mathfrak{M} = \langle M, S \rangle$ , where  $S(x, y) \iff R^*(\kappa, x, y)$ . We claim that  $\mathfrak{M} \models T$ . We need only worry about the cofinality-quantifier. Cofinalities  $< \kappa$  can be expressed in  $L_{\kappa\kappa}$ , so they are preserved both ways. Therefore also cofinality  $\kappa$  is preserved, and no other cofinalities can occur as the models have cardinality  $\kappa$ .  $\square$

Since the logics  $L_{\kappa\omega}(Q_{\aleph_0}^{\text{cf}})$  and  $\mathcal{L}_{\kappa\omega}(Q_{[\aleph_1, \kappa]}^{\text{cf}})$  cannot both be a sublogics of a logic with  $\text{LS}(\kappa)$ , we get from the above lemma:

**Corollary 5** *Suppose  $\kappa > \omega$  is weakly compact. Then there is no strongest weakly  $\kappa$ -compact extension of  $L_{\kappa\omega}$  with  $\text{LS}(\kappa)$ .*

The logic  $L_{\kappa\omega}$  actually satisfies the property  $\text{LS}(< \kappa)$  which is stronger than  $\text{LS}(\kappa)$ . To prove a result like the above corollary for the property

$LS(< \kappa)$  we have to work a little harder. At the same time we extend the proof to extensions of  $L_{\kappa\kappa}$ . Here the cofinality quantifiers  $Q_C^{\text{cf}}$  will not help as  $Q_{\{\lambda\}}^{\text{cf}}$  is definable in  $L_{\kappa\kappa}$  for  $\lambda < \kappa$ . Therefore we use more refined order-type quantifiers.

**Definition 6** Let  $L_{\kappa\lambda}(Q)$  denote the formal extension of  $L_{\kappa\lambda}$  by the generalized quantifier symbol  $Qxy\phi(x, y, \vec{z})$ . If  $\mathcal{Y}$  is a class of ordinals, we get a logic  $L_{\kappa\lambda}(Q, \mathcal{Y})$  from  $L_{\kappa\lambda}(Q)$  by defining the semantics by

$$\mathfrak{A} \models Qxy\phi(x, y, \vec{c}) \iff \text{otp}(\{\langle a, b \rangle : \mathfrak{A} \models \phi(a, b, \vec{c})\}) \in \mathcal{Y}.$$

If  $\phi \in L_{\kappa\lambda}(Q, \mathcal{Y})$  and  $\mathfrak{A} \models \phi$ , we say that  $\mathfrak{A} \models \phi$  holds in the  $\mathcal{Y}$ -interpretation.

If  $\mathfrak{A}$  is a model, then

$$o(\mathfrak{A}, \mathcal{Y}, \kappa, \lambda)$$

is the supremum of all  $\text{otp}(\{\langle a, b \rangle : \mathfrak{A} \models \phi(a, b, \vec{c})\})$  where  $\phi \in L_{\kappa\lambda}(\mathcal{Y})$ ,  $\vec{c} \in A^{<\lambda}$  and  $\{\langle a, b \rangle : \mathfrak{A} \models \phi(a, b, \vec{c})\}$  is well-ordered.

**Lemma 7** Suppose  $\kappa \geq \lambda$ ,  $\phi \in L_{\kappa\lambda}(Q)$ ,  $\mathfrak{A}$  is a model,  $\vec{a} \in A^{<\lambda}$ , and  $\mathcal{Y} \cap o(\mathfrak{A}, \mathcal{Y}, \kappa, \lambda) = \mathcal{Y}$ . Then  $\mathfrak{A} \models \phi(\vec{a})$  in the  $\mathcal{Y}$ -interpretation if and only if  $\mathfrak{A} \models \phi(\vec{a})$  in the  $\mathcal{Y}'$ -interpretation.

**Proof.** This is a straightforward induction of the length of the formula  $\phi$ .  $\square$

**Lemma 8** 1. Suppose  $\kappa > \omega$ ,  $\phi \in L_{\kappa\kappa}(Q)$ , and  $\phi$  has a model  $\mathfrak{A}$  in the  $\mathcal{Y}$ -interpretation. Then there is a submodel  $\mathfrak{B}$  of  $\mathfrak{A}$  of cardinality  $\leq 2^\kappa$  and  $\mathcal{Y}' \subseteq (2^\kappa)^+$  such that  $\mathcal{Y}' \cap \kappa = \mathcal{Y}$  and  $\mathfrak{B} \models \phi$  in the  $\mathcal{Y}'$ -interpretation.

2. Suppose  $\kappa = \kappa^{<\kappa}$ ,  $T \subseteq L_{\kappa\kappa}(Q)$ ,  $|T| \leq \kappa$  and  $T$  has a model  $\mathfrak{A}$  in the  $\mathcal{Y}$ -interpretation. Then for all  $\xi < \kappa^+$  there is a submodel  $\mathfrak{B}$  of  $\mathfrak{A}$  of cardinality  $\leq \kappa$  and  $\mathcal{Y}' \subseteq \kappa^+$  such that  $\mathcal{Y} \cap \xi = \mathcal{Y}' \cap \xi$  and  $\mathfrak{B} \models T$  in the  $\mathcal{Y}'$ -interpretation.

**Proof.** We may assume  $|A| \geq 2^\kappa$ . Let us expand  $\mathfrak{A}$  by

1. A well-ordering  $\prec$  the order-type of which exceed all the order-types of well-orderings definable by subformulas of  $\phi$  with parameters in  $A$ .
2. A new predicate  $P$  which contains those elements  $d$  of  $A$  for which  $\text{otp}(\{\langle a, b \rangle : a \prec b \prec d\}) \in \mathcal{Y}$

3. A prediacte  $F$  which codes an isomorphism from each well-ordering, definable by a subformula of  $\phi$  with parameters in  $A$ , onto an initial segment of  $\prec$ .

Let  $\langle \mathfrak{A}, \prec, P, F \rangle$  be the expanded structure and  $\langle \mathfrak{B}, \prec^*, P^*, F^* \rangle$  an  $L_{\kappa\kappa}$ -elementary substructure of it of cardinality  $\leq 2^\kappa$ . Let

$$\mathcal{Y}' = \{\text{otp}(\{ \langle a, b \rangle \in B^2 : a \prec^* b \prec^* d \} : d \in P^*)\}$$

It is easy to see that  $\mathfrak{B} \models \phi$  in the  $\mathcal{Y}'$ -interpretation. The second claim is proved similarly.  $\square$

Let  $\pi$  be a canonical bijection from triples of ordinals to ordinals such that  $\pi[\kappa^3] = \kappa$  for each infinite cardinal  $\kappa$ . We say that a pair  $(\delta_1, Z_1)$  codes a pair  $(\delta_2, Z_2)$  if  $\delta_1, \delta_2$  are ordinals,  $Z_1 \subseteq \delta_1$ ,  $Z_2 \subseteq \delta_2$  and there is a bijection  $f : \delta_2 \rightarrow \delta_1$  such that

1.  $\delta_1$  is closed under  $\pi$
2.  $\pi(0, \alpha, \beta) \in Z_1 \iff f(\alpha) < f(\beta)$
3.  $\pi(1, 0, \alpha) \in Z_1 \iff f(\alpha) \in Z_2$ .

Suppose  $\kappa$  is an uncountable cardinal. The *weakly compact ideal* on  $\kappa$  is the ideal of subsets of  $\kappa$  generated by the sets  $\{\alpha : (H(\alpha), \epsilon, A \cap H(\alpha)) \models \neg\phi\}$ , where  $A \subseteq H(\kappa)$  and  $\phi$  is a  $\Pi_1^1$ -sentence such that  $(H(\alpha), \epsilon, A \cap H(\alpha)) \models \phi$ .

**Definition 9** *A cardinal  $\kappa$  satisfies  $\diamond(WC)$  if it is weakly compact and there is a sequence  $\langle A_\alpha : \alpha < \kappa \rangle$  such that*

1.  $A_\alpha \subseteq \alpha$  for  $\alpha < \kappa$ .
2.  $(\forall A \subseteq \kappa)(\{\lambda < \kappa : A_\lambda = A \cap \lambda\} \in \mathcal{I}^+)$ , where  $\mathcal{I}$  is the weakly compact ideal on  $\kappa$ .

We make the following remarks without proof: If  $\kappa$  is measurable  $> \omega$ , then  $\kappa$  satisfies  $\diamond(WC)$ . If  $\kappa$  is weakly compact  $> \omega$ , then there is a generic extension which preserves all cardinals and in which  $\kappa$  satisfies  $\diamond(WC)$ . If  $V=L$ , then every weakly compact cardinal  $> \omega$  satisfies  $\diamond(WC)$ .

**Theorem 10** *Suppose  $\kappa > \omega$  satisfies  $\diamond(WC)$  and  $2^\kappa = \kappa^+$ . Then there is no strongest extension of  $L_{\kappa\kappa}$  for which  $\kappa$  is weakly compact and which has  $LS(< \kappa)$ .*

**Proof.** We shall construct two sets  $\mathcal{Y}^1, \mathcal{Y}^2 \subseteq \kappa^+$  such that the logics  $L_{\kappa\kappa}(Q, \mathcal{Y}^i)$  are weakly  $\kappa$ -compact for and these logics satisfy  $\text{LS}(< \kappa)$ , but no logic containing both  $L_{\kappa\kappa}(Q, \mathcal{Y}^1)$  and  $L_{\kappa\kappa}(Q, \mathcal{Y}^2)$  satisfies  $\text{LS}(< \kappa)$ . The sets  $\mathcal{Y}^i$  are constructed by induction together with ordinals  $\xi_\alpha^i < \kappa^+$  such that:

$$\begin{aligned} \mathcal{Y}^i &= \bigcup_{\alpha < \kappa^+} \mathcal{Y}_\alpha^i \\ \mathcal{Y}_0^i &= \emptyset & \xi_0^i &= 0 \\ \mathcal{Y}_\alpha^i &= \mathcal{Y}_\beta^i \cap \xi_\alpha^i & \text{for } \alpha < \beta \\ \xi_\alpha^i &\leq \xi_\beta^i & \text{for } \alpha < \beta \\ \mathcal{Y}_\nu^i &= \bigcup_{\alpha < \nu} \mathcal{Y}_\alpha^i, & \xi_\nu^i &= \bigcup_{\alpha < \nu} \xi_\alpha^i, \text{ for } \nu = \bigcup \nu \\ \mathcal{Y}_\alpha^1 \cap \mathcal{Y}_\alpha^2 &= \emptyset & \text{for } \alpha < \kappa \end{aligned}$$

First we define  $\mathcal{Y}_\alpha^i$  for  $\alpha < \kappa$  in such a way that  $L_{\kappa\kappa}(\mathcal{Y}^i)$  will in the end have the property  $\text{LS}(< \kappa)$ .

Let  $S_1, S_2$  be a partition of the set of cardinals  $< \kappa$  into two stationary sets. Let  $\{\phi_\nu^i : \nu \in S_i\}$  list all  $L_{\kappa\kappa}(Q)$ -sentences so that each sentence is listed as  $\phi_\nu^i$  for stationary many  $\nu \in S_i$ .

Suppose  $\alpha = \lambda + 1$  and  $\xi_\lambda^i = \lambda$ . Suppose  $\lambda \in S_i$ .

**Case 1.** Suppose that  $(\lambda, A_\lambda)$  codes some pair  $(\xi, Z)$ . In this case we let

$$\begin{aligned} \mathcal{Y}_\alpha^i &= \mathcal{Y}_\lambda^i \cup (Z \setminus \lambda), \xi_\alpha^i = \xi \\ \mathcal{Y}_\alpha^{3-i} &= \mathcal{Y}_\lambda^{3-i}. \end{aligned}$$

**Case 2.** Otherwise we let  $\xi_\alpha^i = \lambda$ ,  $\mathcal{Y}_\alpha^i = \mathcal{Y}_\lambda^i$ ,  $\mathcal{Y}_\alpha^{3-i} = \mathcal{Y}_\lambda^{3-i}$ .

Suppose then  $\alpha = \lambda + 2$ ,  $\xi_\lambda^i = \lambda \in S_i$  and we have defined  $\xi_{\lambda+1}^i$  and  $\mathcal{Y}_{\lambda+1}^i$ .

**Case 3.** The sentence  $\phi_\lambda^i$  has a model in the  $\mathcal{Y}$ -interpretation for some  $\mathcal{Y} \subseteq \kappa^+$  with  $\mathcal{Y} \cap \xi_{\lambda+1}^i = \mathcal{Y}_{\lambda+1}^i$ . By Lemma 8 part 2,  $\phi_\lambda^i$  has a model  $\mathfrak{A}$  of cardinality  $< \kappa$  in the  $\mathcal{Y}$ -interpretation for some  $\mathcal{Y} \subseteq \kappa$  of cardinality  $< \kappa$  with  $\mathcal{Y} \cap \xi_{\lambda+1}^i = \mathcal{Y}_{\lambda+1}^i$ . Let  $\mu$  be minimal such that  $\phi_\lambda^i \in \mathcal{L}_{\mu\mu}(\mathcal{Y})$ . Let  $\xi_{\lambda+2}^i = o(\mathfrak{A}, \mathcal{Y}, \mu, \mu)$  and  $\mathcal{Y}_{\lambda+2}^i = \mathcal{Y}$ . Let  $\mathcal{Y}_{\lambda+2}^{3-i} = \mathcal{Y}_{\lambda+1}^{3-i}$ .

**Case 4.** Otherwise  $\xi_{\lambda+1}^i = \xi_\lambda^i$ ,  $\mathcal{Y}_\alpha^i = \mathcal{Y}_{\lambda+1}^i$ ,  $\mathcal{Y}_\alpha^{3-i} = \mathcal{Y}_{\lambda+1}^{3-i}$ .

Finally for all other  $\alpha \leq \kappa$  we let  $\xi_\alpha^i$  and  $\mathcal{Y}_\alpha^i$  be defined canonically.

This ends the construction of  $\mathcal{Y}_\alpha^i$  for  $\alpha \leq \kappa$ . Note that  $\mathcal{Y}_\kappa^1 \cap \mathcal{Y}_\kappa^2 = \emptyset$ . Moreover, if  $\phi_\nu^i$  has a model in the  $\mathcal{Y}$ -interpretation for some  $\mathcal{Y} \supseteq \mathcal{Y}_\kappa^i$ , then, by construction,  $\phi_\nu^i$  has a model of cardinality  $< \kappa$  in the  $\mathcal{Y}_\kappa^i$ -interpretation.

Let  $\mathcal{Y}_{\kappa+1}^i = \mathcal{Y}_{\kappa}^i \cup \{\kappa\}$  and  $\xi_{\kappa+1}^i = \kappa + 2$ . Next we shall define  $\mathcal{Y}_{\alpha}^i$  and  $\xi_{\alpha}^i$  for  $\kappa + 1 < \alpha < \kappa^+$ . For this, let  $\langle T_{\alpha} : \kappa < \alpha < \kappa^+ \rangle$  enumerate all  $L_{\kappa\kappa}(Q)$ -theories of cardinality  $\leq \kappa$  in a language of cardinality  $\leq \kappa$  which satisfy the condition that every subset of cardinality  $< \kappa$  has a model in the  $\mathcal{Y}_{\kappa}^i$ -interpretation. Here we use the assumption  $2^{\kappa} = \kappa^+$ . We may assume  $T_{\alpha} \subseteq \kappa$  for all  $\alpha$ .

Suppose  $\mathcal{Y}_{\beta}^i$  and  $\xi_{\beta}^i$  have been defined for  $\beta < \alpha$ . If  $\alpha = \cup \alpha$ ,  $\mathcal{Y}_{\alpha}^i$  and  $\xi_{\alpha}^i$  are defined canonically. So assume  $\alpha = \beta + 1$ . Let  $T : H(\kappa) \rightarrow H(\kappa)$  be the function  $T(a) = T_{\beta} \cap a$ . If  $a \in H(\kappa)$ , then  $T(a)$  has a model  $\mathfrak{B}_a$  in the  $\mathcal{Y}_{\kappa}^i$ -interpretation. By construction, we may assume  $\mathfrak{B}_a \in H(\kappa)$ . Let  $B : H(\kappa) \rightarrow H(\kappa)$  be the function  $B(a) = \mathfrak{B}_a$ . Let  $Z \subseteq \kappa$  code  $(\xi_{\beta}^i, \mathcal{Y}_{\beta}^i)$ . By  $\diamond(\text{WC})$ ,  $W = \{\lambda < \kappa : A_{\lambda} = Z \cap \lambda\} \in \mathcal{I}^+$ , where  $\mathcal{I}$  is the weakly compact ideal on  $\kappa$ . Let  $A : \kappa \rightarrow H(\kappa)$  be the function  $A(\alpha) = A_{\alpha}$ . By the definition of  $\mathcal{I}$ , there are a transitive set  $M$  and  $Y^*, R^*$  such that

$$\langle H(\kappa), \epsilon, A, W, \mathcal{Y}_{\kappa}^i, B, T \rangle \prec_{\kappa\kappa} \langle M, \epsilon, A^*, W^*, Y^*, B^*, T^* \rangle$$

and  $\kappa \in W^*$ . Now  $A^*(\kappa) = Z$  and, by construction,  $Y^* \cap \xi_{\beta}^i = \mathcal{Y}_{\beta}^i$

It is clear now that  $B(\kappa)$  is a model of  $T_{\alpha}$  in the  $Y^*$ -interpretation. By Lemma 8 there is a model  $\mathfrak{B}$  of cardinality  $\leq \kappa$  of  $T_{\beta}$  in the  $Y^{**}$ -interpretation for some  $Y^{**}$  with  $Y^{**} \cap \xi_{\beta}^i = \mathcal{Y}_{\beta}^i$ . Let  $\xi_{\alpha}^i = o(\mathfrak{B}, Y^{**}, \kappa, \kappa)$  and  $\mathcal{Y}_{\alpha}^i = Y^{**} \cap \xi_{\alpha}^i$ .

Finally, let  $\mathcal{Y}^i = \bigcup_{\alpha < \kappa^+} \mathcal{Y}_{\alpha}^i$ .

**Claim 1.**  $L_{\kappa\kappa}(\mathcal{Y}^i)$  satisfies the  $\text{LS}(< \kappa)$ -property.

Suppose  $\phi$  is a sentence of  $L_{\kappa\kappa}(\mathcal{Y}^i)$  with a model. Let  $\lambda \in S_i$  such that  $\xi_{\lambda}^i = \lambda$  and  $\phi_{\lambda}^i = \phi$ . By the construction of  $\mathcal{Y}_{\lambda+2}^i$  there is a model of  $\phi$  of cardinality  $< \kappa$ .

**Claim 2.**  $L_{\kappa\kappa}(\mathcal{Y}^i)$  is weakly  $\kappa$ -compact.

Suppose  $T \subseteq L_{\kappa\kappa}(\mathcal{Y}^i)$  is given and every subset of  $T_{\alpha}$  of cardinality  $< \kappa$  has a model in the  $\mathcal{Y}^i$ -interpretation. Then  $T = T_{\alpha}$  for some  $\alpha$ . By construction, every subset of  $T_{\alpha}$  of cardinality  $< \kappa$  has a model in the  $\mathcal{Y}^i \cap \kappa$ -interpretation. Thus the definition of  $\mathcal{Y}_{\alpha}^i$  is made so that  $T_{\alpha}$  has a model  $\mathfrak{B}$  in the  $\mathcal{Y}$ -interpretation for some  $\mathcal{Y}$  such that  $\mathcal{Y} \cap o(\mathfrak{B}, \mathcal{Y}, \kappa, \kappa) = \mathcal{Y}^i \cap o(\mathfrak{B}, \mathcal{Y}, \kappa, \kappa)$ . Thus by Lemma 7,  $\mathfrak{B} \models T_{\alpha}$  in the  $\mathcal{Y}^i$ -interpretation. The Claim is proved.

We can now finish the proof of the theorem. In a logic in which both the quantifier  $Q_{\mathcal{Y}^1}$  and  $Q_{\mathcal{Y}^2}$  are definable, we can say that the order-type of a well-ordering is in  $\mathcal{Y}^1 \cap \mathcal{Y}^2$ . Thus such a logic cannot satisfy  $\text{LS}(< \kappa)$ .  $\square$



It is interesting to note that a proof like above would not be possible for the following stronger Löwenheim-Skolem property: A *filter-family* is a family  $\mathcal{F} = (\mathcal{F}(A))_{A \neq \emptyset}$ , where  $\mathcal{F}(A)$  is always a filter on the set  $A$ . Luosto [6] defines the concept of a  $(\kappa^+, \omega)$ -*neat* filter family. We will not repeat the definition here, its elements are closure under bijections, fineness,  $\kappa^+$ -completeness, normality and upward relativizability (all defined in [6]). Suppose  $\mathcal{L}$  is a logic of the form  $L_{\kappa\lambda}(\vec{Q})$  for some sequence  $\vec{Q}$  of generalized quantifiers. We say that  $\mathcal{L}$  has the  $\mathcal{F}, \kappa$ -*persistence property*, if for all models  $\mathfrak{A}$  and  $B \in \mathcal{F}(A)$ , we have  $\mathfrak{A} \upharpoonright B \prec \mathfrak{A}$ . Luosto proves that if  $\mathcal{L}_1$  and  $\mathcal{L}_2$  both satisfy the  $\mathcal{F}, \kappa$ -persistence property, then there is  $\mathcal{L}_3$  such that  $\mathcal{L}_1 \leq \mathcal{L}_3$ ,  $\mathcal{L}_2 \leq \mathcal{L}_3$  and  $\mathcal{L}_3$  satisfies the  $\mathcal{F}, \kappa$ -persistence property. Lipparini [5] proves a similar result for families of limit ultrafilters related closely to compactness.

Tapani Hyttinen pointed out that the assumption  $2^\kappa = \kappa^+$  is not needed in Theorem 10, if  $\kappa$  is assumed to be measurable.

## References

- [1] Jon Barwise, Axioms for abstract model theory, *Ann. Math. Logic*, 7, 1974, 221–265.
- [2] Model-theoretic logics, Barwise, J. and Feferman, S., *Perspectives in Mathematical Logic*, Springer-Verlag, New York, 1985, xviii+893.
- [3] Gebhard Fuhrken, Skolem-type normal forms for first-order languages with a generalized quantifier, *Fund. Math.*, 54, 1964, 291–302.
- [4] Per Lindström, On extensions of elementary logic, *Theoria*, 35, 1969, 1–11.
- [5] Paolo Lipparini, Limit ultrapowers and abstract logics, *Journal of Symbolic Logic* vol. 52 (1987), 437–454.
- [6] Kerkko Luosto, Filters in abstract model theory, Ph.D. Thesis, University of Helsinki, 1992, 81 pages.
- [7] Janos Makowsky and Saharon Shelah, The theorems of Beth and Craig in abstract model theory. II. Compact logics, *Archiv für Mathematische Logik und Grundlagenforschung*, 21, 1981, 13–35.

- [8] Andrzej Mostowski, On a generalization of quantifiers, *Fund. Math.*, 44, 1957, 12–36
- [9] Saharon Shelah, Generalized quantifiers and compact logic, *Trans. Amer. Math. Soc.*, 204, 1975, 342–364.
- [10] Saharon Shelah, Models with second-order properties. I. Boolean algebras with no definable automorphisms, *Ann. Math. Logic, Annals of Mathematical Logic*, 14, 1978, 1, 57–72.
- [11] Marek Waclawek, On ordering of the family of logics with Skolem-Löwenheim property and countable compactness property, in: *Quantifiers: Logics, Models and Computation*, Vol. 2, (Michał Krynicki, Marcin Mostowski and Lesław Szczurba editors), Kluwer Academic Publishers, Dordrecht, Boston, London, 1995, pp.229–236.