

DETAILS ON SH:74

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ABSTRACT. We give details on a claim from [She78] (continuing [She71]+
[?])

Key words and phrases.

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Theorem -1.1. 3

Let $\lambda_1 \geq \mu_1, \lambda \geq \mu$, then the following are equivalent

- (A) $(\lambda_1, \mu_1) \rightarrow_{\aleph_0} (\lambda, \mu)$
- (B) $(\lambda_1, \mu_1) \rightarrow_{\leq \mu} (\lambda, \mu)$
- (C) We can find functions $f_\ell : \lambda^\ell \rightarrow \mu$ for $\ell < \omega$ such that: if (n, E) is not an identity of (λ_1, μ_1) then $\langle f_\ell : \ell \leq n \rangle$ witness that it is not an identity of (λ, μ) .

REMARK :

- (1) If \square the sets of (n, E) which are not identifies of (λ_1, μ_1) is recursive, we can add (D) $(\lambda_1, \mu_1) \rightarrow_1 (\lambda, \mu)$.
- (2) We can weaken \square to: \square^+ there is a recursive set of identities, including the one failing for (λ_1, μ_1) and included in the one holding for (λ, μ) (check).

REMARK: Recall $(\lambda_1, \mu_1) \rightarrow_{\leq \kappa} (\lambda, \mu)$ mean that if T is a (first order theory of cardinality $\leq \kappa$, with the distinguish predicates P_1, P_2 and every finite $T' \subseteq T$ has a model M with $|P_1^M| = \lambda_1, |P_2^M| = \mu_1$ then T has a model N with $|P_1^N| = \lambda, |P_2^N| = \mu$

Proof. of theorem 3

The proof is by showing (for our given $\lambda_1, \mu_1, \lambda, \mu$)

$$(A) \Rightarrow (C) \Rightarrow (B) \Rightarrow (A)$$

$(B) \Rightarrow (A)$ trivially.

$(A) \Rightarrow (C)$: Let $\langle (n_i, E_i) : i < \omega \rangle$ list the identities. Let $m_i = \text{Max}\{i, n_0, \dots, n_i\}$

For each i , choose if possible f_ℓ^i , an ℓ - place function from λ_1 to μ_1 for $\ell \leq m_i$ exemplifying $\lambda_1 \not\rightarrow (n_i, E_i)_{\mu_1}$; if impossible f_ℓ^i is chosen identically zero. We do it by induction on i and so without loss of generality

$$\otimes_1 \quad i < j < \omega \ \& \ \ell \leq m_i \Rightarrow f_\ell^j \text{ refine } f_\ell^i \quad \text{i.e. } f_\ell^j(\bar{x}) = f_\ell^j(\bar{y}) \rightarrow f_\ell^i(\bar{x}) = f_\ell^i(\bar{y}) \text{ (recall } \mu^n = \mu)$$

Let $M_1 = (\lambda_1, \mu_1 \dots f_\ell^i, \dots)$, so it has $i < \omega, \ell \leq m_i$ universe λ_1 and vocabulary $\{F_\ell^i : i < \omega, \ell \leq m_i\}$ where F_ℓ^i is an ℓ -place function.

This is not exactly right so let M_2 be defined by

- (a) M_2 has universe λ_1
relations:
- (b) $P_1^{M_2} = \lambda_1$
- (c) $P_2^{M_2} = \mu_1$
- (d) F_ℓ an $(\ell+1)$ - place function such that for $i < \omega, \forall \bar{x}[F_\ell(\bar{x}, i) = f_\ell^i(\bar{x})]$ otherwise zero
- (e) $P_3^{M_2} = \omega$
- (f) $c_n^{M_2} = n$
- (g) $<^{M_2}$ - the usual order on λ_1

Let M_2^+ be the expansion of M_2 by Skolem functions.

Lastly let $T = \text{Th}(M_2) \cup \{c_n < c \wedge P_3(c) : n < \omega\}$. Clearly.

⊗₂ T is first order, countable and every finite subset has a (λ_1, μ_1) -model.

As we are assuming (1) there is a model N of T , with

$$|P_1^N| = \|N\| = \lambda, |P_2^N| = \mu;$$

so without loss of generality $P_1^N = \lambda, P_2^N = \mu$. Let $f_\ell^* : {}^\ell\lambda \rightarrow \lambda$ be

$$f_\ell^*(\bar{x}) = F_\ell^N(\bar{x}, c^N)$$

⊗₃ if (n_i, E_i) is not an identity of (λ_1, μ_1) then $\langle f_\ell^* : \ell \leq n_i \rangle$ witness it is not an identity of (λ, μ)

[Why? because $M_2 \models \langle f_\ell^j : \ell \leq n_i \rangle$ witness (n_i, E_i) is not an identity of (λ_1, μ_1) ” is known when $j = i$ by its choice, and if $j \leq i$ by ⊗₁, which means

⊠ $M_2 \models (\forall y)(c_i \leq y \in P_3^{M_2} \rightarrow \langle F_\ell(-, y) : i \leq m_i \rangle$ is a witness to (n_i, E_i) being not an identity),

this (⊠) is expressed by a first order sentence ψ_i which M_2 satisfied hence $\psi_i \in T$ hence $N \models \psi_i$.

In particular use $y = c$ in N recalling $N \models [c_i < c \ \& \ P_3(c)]$ so we have

$$\langle F_\ell^N(-, c) : \ell \leq n_i \rangle \text{ witness } (P_1^N, P_2^N) = (\lambda, \mu)$$

fail the identity (n_i, E_i) which mean that $\langle f_\ell^* : \ell \leq n_i \rangle$ witness (λ, μ) fail the identity (n_i, E_i) . So we have gotten (3).

(C) \Rightarrow (B) Exactly as in [She71]. Define an equivalence relation E on $\bigcup_n {}^n\lambda$ as follows

$$\bar{b}E\bar{c} \text{ iff } \bigvee_n [\bar{b}, \bar{c} \in {}^n\lambda \ \& \ f_n(\bar{b}) = f_n(\bar{c})]$$

By [She71], Lemma 1 it suffice to show (*) of [[She71] p. 194] which is stright.

REMARK: in (3) we have $\langle f_\ell : \ell < \omega \rangle$ for all possible (n, E) not just for each $(n, E) \dots$

□

REFERENCES

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