

## ON COUNTABLY CLOSED COMPLETE BOOLEAN ALGEBRAS

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ABSTRACT. It is unprovable that every complete subalgebra of a countably closed complete Boolean algebra is countably closed.

**Introduction.** A partially ordered set  $(P, <)$  is  $\sigma$ -closed if every countable chain in  $P$  has a lower bound. A complete Boolean algebra  $B$  is *countably closed* if  $(B^+, <)$  has a dense subset that is  $\sigma$ -closed. In [2] the first author introduced a weaker condition for Boolean algebras, *game-closed*: the second player has a winning strategy in the infinite game where the two players play an infinite descending chain of nonzero elements, and the second player wins if the chain has a lower bound. In [1], Foreman proved that when  $B$  has a dense subset of size  $\aleph_1$  and is game-closed then  $B$  is countably closed. (By Vojtáš [5] and Veličković [4] this holds for every  $B$  that has a dense subset of size  $2^{\aleph_0}$ .) We show that, in general, it is unprovable that game-closed implies countably closed. We construct a model in which a  $B$  exists that is game-closed but not countably closed. It remains open whether a counterexample exists in ZFC.

Being game-closed is a hereditary property: If  $A$  is a complete subalgebra of a game-closed complete Boolean algebra  $B$  then  $A$  is game-closed. It is observed in [3] that every game-closed algebra is embedded in a countably closed algebra; in fact,

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for a forcing notion  $(P, <)$ , being game-closed is equivalent to the existence of a  $\sigma$ -closed forcing  $Q$  such that  $P \times Q$  has a dense  $\sigma$ -closed subset. Hence the statement “every game-closed complete Boolean algebra is countably closed” is equivalent to the statement “every complete subalgebra of a countably closed complete Boolean algebra is countably closed”.

Below we construct (by forcing) a model of ZFC+GCH and in it a partial ordering  $P$  of size  $\aleph_2$  such that  $B(P)$ , the completion of  $P$ , is not countably closed, but  $B(P \times Col)$  is, where  $Col$  is the Lévy collapse of  $\aleph_2$  to  $\aleph_1$  (with countable conditions).

**Theorem.** *It is consistent that there exists a partial ordering  $(P, <)$  such that  $B(P)$  is not countably closed but  $B(P \times Col)$  is countably closed.*

### Forcing Conditions.

We assume that the ground model satisfies *GCH*.

We want to construct, by forcing, a partially ordered set  $(P, <_P)$  of size  $\aleph_2$  that has the desired properties. We shall use as forcing conditions countable approximations of  $P$ . One part of a forcing condition will thus be a countable partial ordering  $(A, <_A)$  with the intention that  $A$  be a subset of  $P$  and that the relation  $<_A$  on  $A$  be the restriction of  $<_P$ . As  $P$  will have size  $\aleph_2$ , we let  $P = \omega_2$ , and so  $A$  is a countable subset of  $\omega_2$ .

The second part of a forcing condition will be a countable set  $B \subset A \times Col$ , a countable approximation of a dense set in the product ordering  $P \times Col$ . The third part of a forcing condition will be a countable set  $C$  of countable descending chains in  $A$  that have no lower bound. Finally, a forcing condition includes a function that guarantees that the limit of the  $B$ 's is  $\sigma$ -closed (and so  $P \times Col$  has a  $\sigma$ -closed dense subset).

Whenever we use  $<$  without a subscript, we mean the natural ordering of ordinal numbers.

**Definition.** For any set  $X$ ,  $Col(X)$  is the set of all countable functions  $q$  such that  $dom(q) \in \omega_1$  and  $range(q) \subset X$ ;  $Col = Col(\omega_2)$ .

**Definition.** The set  $R$  of forcing conditions  $r$  consists of quadruples  $r = ((A_r, <_r), B_r, C_r, F_r)$  such that

- (1)  $A_r$  is a countable subset of  $\omega_2$ ,
  - (2)  $(A_r, <_r)$  is a partially ordered set,
  - (3) if  $b <_r a$  then  $a < b$ ,
  - (4)  $B_r$  is a countable subset of  $A_r \times Col(A_r)$ , and for every  $(p, q) \in B_r$ ,  
 $p \in \text{range}(q)$ ,
  - (5)  $C_r$  is a countable set of countable sequences  $\{a_n\}_{n=0}^\infty$  in  $A_r$  with the property that  $a_0 >_r a_1 >_r \dots >_r a_n >_r \dots$  and that  $\{a_n\}_n$  has no lower bound in  $A_r$ ,
  - (6)  $F_r$  is a function of two variables,  $\{a_n\}_n \in C_r$  and  $(p, q) \in B_r$  such that  $p \geq a_0$ , and  $\text{range}(F_r) \subset \omega$ . If  $m = F_r(\{a_n\}_n, (p, q))$  then for every  $(p', q') \in B_r$  stronger than  $(p, q)$ ,
- (\*) if  $p' <_r a_m$  then  $p' \perp_r \{a_n\}_n$  (i.e.  $p' \perp_r a_k$  for some  $k$ ).

If  $r, s \in R$  then  $r <_R s$  ( $r$  is stronger than  $s$ ) if

- (7)  $A_r \supseteq A_s$ ,
- (8)  $<_r$  and  $<_s$  agree on  $A_s$ , and  $\perp_r$  and  $\perp_s$  agree on  $A_s$ ; i.e. if  $a, b \in A_s$  then  $a <_r b$  iff  $a <_s b$  and  $a \perp_r b$  iff  $a \perp_s b$  for all  $a, b \in A_s$ ,
- (9)  $B_r \supseteq B_s$ ,
- (10)  $C_r \supseteq C_s$ ,
- (11)  $F_r \supseteq F_s$ .

The relation  $<_R$  on  $R$  is a partial ordering. We shall prove that the forcing extension by  $R$  contains a desired example  $(P, <_P)$ . Assuming the GCH in the ground model, the forcing  $R$  preserves cardinals and  $V^R$  is a model of  $ZFC + GCH$ ; this follows from the next two lemmas:

**Lemma 1.**  $R$  is  $\sigma$ -closed.

*Proof.* Let  $\{r_n\}_n$  be a sequence of conditions such that  $r_0 >_R r_1 >_R \cdots >_R r_n >_R \cdots$ . We show that  $\{r_n\}_n$  has a lower bound.

Assuming that for each  $n$ ,  $r_n = ((A_n, <_n), B_n, C_n, F_n)$ , we let  $A_r = \bigcup_{n=0}^{\infty} A_n$ ,  $B_r = \bigcup_{n=0}^{\infty} B_n$ ,  $C_r = \bigcup_{n=0}^{\infty} C_n$ ,  $F_r = \bigcup_{n=0}^{\infty} F_n$  and  $<_r = \bigcup_{n=0}^{\infty} <_n$ ; we claim that  $r = ((A_r, <_r), B_r, C_r, F_r)$  is a condition, and is stronger than each  $r_n$ .

The quadruple  $r$  clearly has properties (1)–(4). It is also easy to see that for every  $n$ ,  $<_r$  agrees with  $<_n$  and  $\perp_r$  agrees with  $\perp_n$  on  $A_n$ . To verify (5), let  $\{a_n\}_n \in C_r$ . There is an  $m$  such that  $\{a_n\}_n \in C_k$  for all  $k \geq m$ , and therefore  $\{a_n\}_n$  has no lower bound in any  $A_k$ . Thus  $\{a_n\}_n$  has no lower bound in  $A_r$ . Finally, to verify (6), let  $F_r(\vec{a}, (p, q)) = m$  and let  $(p', q')$  be stronger than  $(p, q)$ . Since (\*) holds in  $r_n$  where  $n$  is large enough so that  $\vec{a} \in C_n$  and  $(p, q), (p', q') \in B_n$ , (\*) holds in  $r$  as well.

Therefore  $r$  is a condition and for every  $n$ ,  $r$  is stronger than  $r_n$ .

**Lemma 2.**  *$R$  has the  $\aleph_2$ -chain condition.*

*Proof.* If  $W$  is a set of conditions of size  $\aleph_2$ , then a  $\Delta$ -system argument (using CH) yields two conditions  $r, s \in W$  such that if  $r = ((A_r, <_r), B_r, C_r, F_r)$  and  $s = ((A_s, <_s), B_s, C_s, F_s)$ , then there is a  $D$  (the root of the  $\Delta$ -system) such that  $D = A_r \cap A_s$ ,  $\sup D < \min(A_r - D)$ ,  $\sup A_r < \min(A_s - D)$ ,  $<_r$  and  $<_s$  agree on  $D$ ,  $\perp_r$  and  $\perp_s$  agree on  $D$ ,  $B_r \cap (D \times \text{Col}(D)) = B_s \cap (D \times \text{Col}(D))$ ,  $C_r \cap D^\omega = C_s \cap D^\omega$ , and  $F_r(\vec{a}, (p, q)) = F_s(\vec{a}, (p, q))$  whenever  $\vec{a} \in C_r \cap D^\omega$  and  $(p, q) \in B_r \cap (D \times \text{Col}(D))$ .

Moreover, there exists a mapping  $\pi$  of  $A_s$  onto  $A_r$  that is an isomorphism between  $s$  and  $r$  and is the identity on  $D$ .

Let  $t = ((A_t, <_t), B_t, C_t, F_t)$  where  $A_t = A_r \cup A_s$ ,  $B_t = B_r \cup B_s$ ,  $C_t = C_r \cup C_s$ ,  $<_t = <_r \cup <_s$ , and  $F_t$  will be defined below such that  $F_t \supseteq F_r \cup F_s$ . We claim that  $t$  is a condition, and is stronger than both  $r$  and  $s$ ; thus  $r$  and  $s$  are compatible. Properties (1)–(4) are easy to verify. It is also easy to see that  $<_t$  agrees with  $<_r$  on  $A_r$  and with  $<_s$  on  $A_s$ , and  $\perp_t$  agrees with  $\perp_r$  on  $A_r$  and with  $\perp_s$  on  $A_s$ .

Note that if  $a \in A_r - D$  and  $b \in A_s - D$  then  $a \perp_t b$ . Thus if  $\{a_n\}_n$  is in  $C_r$

but not in  $C_s$  (or vice versa) then  $\{a_n\}_n$  has no lower bound in  $A_r \cup A_s$ , and so (5) holds.

In order to deal with (6), we first verify it for the values of  $F_t$  inherited from either  $r$  or  $s$ . Thus let  $\vec{a} \in C_r$ ,  $(p, q) \in B_r$ ,  $m = F_r(\vec{a}, (p, q))$  and let  $(p', q') \in B_t$  be stronger than  $(p, q)$ . (The argument for  $s$  in place of  $r$  is completely analogous.) If  $(p', q') \in B_r$  then (\*) holds in  $r$  and therefore in  $t$ . Thus assume that  $(p', q') \in B_s$ .

Since  $p' \in A_s$  and  $p' <_t p$ , it follows that  $p \in D$ , and since  $\text{range}(q) \subseteq \text{range}(q') \subseteq A_s$ , we have  $(p, q) \in B_s$ . Now if  $\vec{a} \in C_s$  then  $F_s(\vec{a}, (p, q)) = F_r(\vec{a}, (p, q))$  and so  $p'$  satisfies (\*) in  $s$  and hence in  $t$ .

If  $\vec{a} \notin C_s$  and  $p' \notin A_r$  then  $p' \perp_t \vec{a}$  and again  $p'$  satisfies (\*).

The remaining case is when  $p' \in D$  and  $(p, q) \in B_r \cap B_s$ . Since  $(p', \pi q') = (\pi p', \pi q')$  is stronger than  $(p, q) = (\pi p, \pi q)$ ,  $p'$  satisfies (\*) in  $r$  and therefore in  $t$ .

To complete the verification of (6) we define  $F_t(\vec{a}, (p, q))$  for those  $\vec{a}$  and  $(p, q)$  that come from the two different conditions. Let  $\vec{a} \in C_r - C_s$  and  $(p, q) \in B_s - B_r$  (the other case being analogous) be such that  $p \geq a_0$ . We let  $F_t(\vec{a}, (p, q))$  be the least  $m$  such that  $a_m \notin D$ .

Let  $(p', q') \in B_t$  be stronger than  $(p, q)$ ; we shall show that  $p' \not\leq_t a_m$ . This is clear if  $p' \in D$ , by (3). If  $p' \notin D$  then we claim that  $p'$  cannot be in  $A_r$ ; then it follows that  $p' \perp_t a_m$ . To prove the claim, note that  $\text{range}(q) \not\subseteq A_r$  (because  $(p, q) \notin B_r$ ) and hence  $\text{range}(q') \subseteq A_s$ . By (4),  $p' \in A_s$  and so  $p' \notin A_r$ .

Therefore  $t$  is a condition and is stronger than both  $r$  and  $s$ .

Let  $G$  be a generic filter on  $R$ . In  $V_G$ , we let  $P = \bigcup\{A_r : r \in G\}$ ,  $<_P = \bigcup\{<_r : r \in G\}$ , and  $Q = \bigcup\{B_r : r \in G\}$ .  $(P, <_P)$  is a partial ordering and  $Q \subset P \times \text{Col}$ . We shall prove that  $Q$  is  $\sigma$ -closed and is dense in  $P \times \text{Col}$ , and that the complete Boolean algebra  $B(P)$  does not have a dense  $\sigma$ -closed subset.

**Lemma 3.**  $P = \omega_2$ .

*Proof.* We prove that for every  $s$  and every  $p \in \omega_2$  there exists an  $r <_R s$  such that  $p \in A_r$ . But this is straightforward: let  $A_r = A_s \cup \{p\}$ ,  $B_r = B_s$ ,  $C_r = C_s$ ,

$F_r = F_s$  and  $\langle_r = \langle_s$ ; properties (1)–(11) are easily verified. (Note that  $p \perp_r a$  for all  $a \in A_s$ .)

**Lemma 4.**  *$Q$  is dense in  $P \times Col$ .*

*Proof.* Let  $s$  be a condition and let  $p_0 \in A_s$  and  $q_0 \in Col$ . We shall find an  $r <_R s$ ,  $p \in A_r$  and  $q \supset q_0$  such that  $p <_r p_0$  and  $(p, q) \in B_r$ : Let  $p$  be an ordinal greater than  $\sup A_s$ , let  $q \in Col$  be such that  $q \supset q_0$  and  $p \in \text{range}(q)$ , and let  $A_r = A_s \cup \text{range}(q)$ ,  $B_r = B_s \cup \{(p, q)\}$ ,  $C_r = C_s$ , and let  $\langle_r$  be the partial order of  $A_r$  that extends  $\langle_s$  by making  $p <_r p_0$ . Finally, let  $F_r(\vec{a}, (p, q)) = 0$  for all  $\vec{a} \in C_r$ .

To see that  $r = ((A_r, \langle_r), B_r, C_r, F_r)$  is a condition, note that for every  $\vec{a} \in C_r$ ,  $p$  is not a lower bound of  $\vec{a}$  (because  $p_0$  isn't) and hence  $p \perp_r \vec{a}$ . This implies both (5) and (6). Since adding  $p$  does not affect the relation  $\perp$  on  $A_s$ , we have (8) and so  $r$  is stronger than  $s$ .

Next we prove that  $Q$  is  $\sigma$ -closed.

**Lemma 5.** *If  $u = \{(p_n, q_n)\}_{n=0}^\infty$  is a descending chain in  $Q$  then  $u$  has a lower bound.*

*Proof.* Let  $\dot{u}$  be a name for a descending chain and let  $s$  be a condition. By extending  $s$   $\omega$  times if necessary ( $R$  is  $\sigma$ -closed), we may assume that there is a sequence  $u = \{(p_n, q_n)\}_{n=0}^\infty$  in  $\omega_2 \times Col$  such that  $s$  forces  $\dot{u} = u$ , such that for every  $n$ ,  $p_n \in A_s$ ,  $(p_n, q_n) \in B_s$ , that  $p_0 >_s p_1 >_s \cdots >_s p_n > \cdots$  is a descending chain in  $(A_s, \langle_s)$  and that  $q_0 \subset q_1 \subset \cdots \subset q_n \subset \cdots$ .

Let  $p$  be an ordinal greater than  $\sup A_s$ , let  $q \supseteq \bigcup_{n=0}^\infty q_n$  be such that  $p \in \text{range}(q) \subseteq A_s \cup \{p\}$ , let  $A_r = A_s \cup \{p\}$ ,  $B_r = B_s \cup \{(p, q)\}$ ,  $C_r = C_s$ , and let  $\langle_r$  be the partial order of  $A_r$  that extends  $\langle_s$  by making  $p$  a lower bound of  $\{p_n\}_{n=0}^\infty$ . Finally, let  $F_r(\vec{a}, (p, q)) = 0$  for all  $\vec{a} \in C_r$  and  $r = ((A_r, \langle_r), B_r, C_r, F_r)$ .

We shall show that for every  $\vec{a} \in C_s$ ,  $p$  is not a lower bound of  $\vec{a}$ . This implies that  $p \perp_r \vec{a}$  and (5) and (6) follow. Since making  $p$  a lower bound of  $\{p_n\}_n$  does not affect the relation  $\perp$  on  $A_s$ , we'll have (8) and hence  $r <_R s$ . In  $r$ ,  $(p, q)$  is a lower bound of  $u$ .

Thus let  $\vec{a} = \{a_k\}_k \in C_s$ . We claim that

$$\exists k \forall n p_n \not\prec_s a_k.$$

This implies that  $p \not\prec_r a_k$  and hence  $p$  is not a lower bound of  $\vec{a}$ .

If  $p_n < a_0$  for all  $n$  then we let  $k = 0$  because then  $p_n \not\prec_s a_0$  for all  $n$ .

Otherwise let  $N$  be the least  $N$  such that  $p_N \geq a_0$ , and let  $m = F_s(\vec{a}, (p_N, q_N))$ .

Either  $p_n \not\prec_s a_m$  for all  $n$  and we are done (with  $k = m$ ) or else  $p_M <_s a_m$  for some  $M \geq N$ . By (\*) there exists some  $k$  such that  $p_M \perp_s a_k$  and hence  $p_n \not\prec_s a_k$  for all  $n$ .

Finally, we shall prove that  $B(P)$  is not countably closed.

**Lemma 6.** *The complete Boolean algebra  $B(P)$  does not have a dense  $\sigma$ -closed subset.*

*Proof.* Assume that  $B(P)$  does have a dense  $\sigma$ -closed subset  $D$ . For  $a, b \in P$ , we define

$$a \prec b \quad \text{if} \quad a <_P b \quad \text{and} \quad \exists d \in D \quad \text{such that} \quad a <_{B(P)} d <_{B(P)} b.$$

The relation  $\prec$  is a partial ordering of  $P$ ,  $(P, \prec)$  is  $\sigma$ -closed,  $a \prec b$  implies  $a <_P b$  and for every  $a \in P$  there is some  $b \in P$  such that  $b \prec a$ .

Toward a contradiction, let  $s$  be a condition and assume that  $s$  forces the preceding statement. For each  $\alpha < \omega_2$ , there exist a condition  $s_\alpha$  stronger than  $s$ , and a descending chain  $\{c_n^\alpha\}_n$  in  $A_{s_\alpha}$  such that  $c_0^\alpha \geq \alpha$  and that for every  $n$ ,  $s_\alpha \Vdash c_{n+1}^\alpha \prec c_n^\alpha$ .

By a  $\Delta$ -system argument we find among these a countable sequence  $r_n = s_{\alpha_n} = ((A_n, <_n), B_n, C_n, F_n)$  and a set  $E$  such that for every  $m$  and  $n$  with  $m < n$  we have  $E = A_m \cap A_n$ ,  $\sup E < \min(A_m - E)$ ,  $\sup A_m < \min(A_n - E)$ ,  $<_m$  and  $<_n$  agree on  $E$ ,  $\perp_m$  and  $\perp_n$  agree on  $E$ ,  $B_m \cap (E \times \text{Col}(E)) = B_n \cap (E \times \text{Col}(E))$ ,  $C_m \cap E^\omega = C_n \cap E^\omega$ , and  $F_m(\vec{a}, (p, q)) = F_n(\vec{a}, (p, q))$  whenever  $\vec{a} \in C_m \cap E^\omega$  and  $(p, q) \in B_m \cap (E \times \text{Col}(E))$ . Moreover, there exists a mapping  $\pi_{mn}$  of  $A_m$  onto  $A_n$  that is an isomorphism between  $(r_m, \{c_k^{\alpha_m}\}_k)$  and  $(r_n, \{c_k^{\alpha_n}\}_k)$  and is the

identity on  $E$ . We also let  $\pi_{nm} = \pi_{mn}^{-1}$ ,  $\pi_{mm} = id$  and assume that the  $\pi_{mn}$  form a commutative system. Note that for every  $n$  and  $k$ ,  $c_k^{\alpha_n} \notin E$ .

For each  $n$  and  $k$ , let  $a_k^n = c_{2k}^{\alpha_n}$  and  $b_k^n = c_{2k+1}^{\alpha_n}$ . Let  $\vec{u} = \{u_n\}_n$  be the ‘‘diagonal sequence’’

$$u_{2n} = a_n^n, \quad u_{2n+1} = b_n^n.$$

We shall find a condition  $t = ((A_t, <_t), B_t, C_t, F_t)$  stronger than all  $r_n$  such that the diagonal sequence  $\vec{u}$  is a descending chain and belongs to  $C_t$ . Since  $t \Vdash b_n^n \prec a_n^n$  for every  $n$ , it forces that  $(P, \prec)$  is not  $\sigma$ -closed. This will complete the proof.

To construct  $t$  we first let  $A_t = \bigcup_{n=0}^{\infty} A_n$  and  $B_t = \bigcup_{n=0}^{\infty} B_n$ . Let  $<_t$  be the minimal partial ordering extending  $\bigcup_{n=0}^{\infty} <_n$  such that for every  $n$ ,  $a_{n+1}^{n+1} <_t b_n^n$ . Before proceeding to define  $C_t$  and  $F_t$  we shall prove some properties of  $(A_t, <_t)$ .

**Lemma 7.** (i) *Let  $m < n$  and let  $y \in A_m - E$  and  $x \in A_n - E$ . If  $x <_t y$  then  $x \leq_n a_n^n$  and  $b_m^m \leq_m y$ . If  $x$  and  $y$  are compatible in  $<_t$  then  $b_m^m \leq_m y$ .*

(ii) *For all  $m$  and  $n$ , if  $x \in A_n$  and  $y \in A_m$  and if  $x <_t y$  then  $x <_n \pi_{mn}y$  (and  $\pi_{nm}x <_m y$ ). In particular, if  $x, y \in A_n$  then  $x <_t y$  if and only if  $x <_n y$ .*

(iii) *For all  $m$  and  $n$ , if  $x \in A_n$  and  $y \in A_m$  and if  $x$  and  $y$  are compatible in  $<_t$  then  $x$  and  $\pi_{mn}y$  are compatible in  $<_n$  (and  $\pi_{nm}x$  and  $y$  are compatible in  $<_m$ ). In particular, if  $x, y \in A_n$  then  $x \perp_t y$  if and only if  $x \perp_n y$ .*

*Proof.* (i) The first statement is an obvious consequence of the definition of  $<_t$ , and the second follows because any  $z$  such that  $z \leq_t x$  is in some  $A_k - E$  where  $k \geq n$ .

(ii) Let  $x \in A_n$  and  $y \in A_m$  and let  $x <_t y$ . First assume that  $y \notin E$  (and so  $x \notin E$ .) Necessarily,  $m \leq n$  and if  $m = n$  then clearly  $x <_n y$ . Thus consider  $m < n$ . By (i)  $x \leq_n a_n^n <_n b_m^m = \pi_{mn}(b_m^m) \leq_n \pi_{mn}y$ .

Now assume that  $y \in E$  and proceed by induction on  $x$ . If  $x \in E$  then  $x <_n y$ . If  $x \notin E$  then either  $x <_n y$  or there exists some  $z \notin E$  such that  $x <_t z <_t y$ , and by the induction hypothesis  $z <_k \pi_{mk}y$  (where  $z \in A_k$ ). Applying the preceding paragraph to  $x$  and  $z$  we get  $\pi_{nk}x <_k z$  and hence  $\pi_{nk}x <_k \pi_{mk}y$ . The statement now follows.

(iii) Let  $x \in A_n$  and  $y \in A_m$  and let  $z \in A_k$  be such that  $z <_t x$  and  $z <_t y$ . By (ii) we have  $\pi_{kn}z <_n x$  and  $\pi_{km}z <_m y$ . Hence  $\pi_{kn}z = \pi_{mn}\pi_{km}z <_n \pi_{mn}y$ . The second statement follows from this and from the second statement of (ii).

Lemma 7 guarantees that  $t$  will be stronger than every  $r_n$ . Another consequence is that if  $\vec{a} \in C_n$  then  $\vec{a}$  has no lower bound in  $<_t$ : if  $x \in A_m$  were a lower bound then  $\pi_{mn}x$  would be a lower bound in  $<_n$ .

Let  $C_t = \bigcup_{n=0}^{\infty} C_n \cup \{\vec{u}\}$ . Every sequence in  $C_t$  is a descending chain in  $<_t$  without a lower bound (clearly,  $\vec{u}$  has no lower bound).

**Lemma 8.** *For all  $k$  and  $n$ , if  $(p, q) \in B_k - B_n$  and if  $(p', q') \in B_t$  is stronger than  $(p, q)$  then  $(p', q') \in B_k - B_n$ .*

*Proof.* Since  $(p, q) \notin B_n$ , we have either  $\text{range}(q) \not\subseteq E$  or  $p \notin E$ , in which case  $p \in \text{range}(q)$  by (4) and again  $\text{range}(q) \not\subseteq E$ . Since  $q \subseteq q'$  it must be the case that  $(p', q') \in B_k - B_n$ .

We shall now define  $F_t$  so that  $F_t \supset \bigcup_{n=0}^{\infty} F_n$  and verify (6). This will complete the proof.

First we let  $F_t(\vec{a}, (p, q)) = F_n(\vec{a}, (p, q))$  whenever the right-hand side is defined; we have to show that (6) holds in  $t$ . Let  $m = F_n(\vec{a}, (p, q))$  and let  $(p', q') \in B_k$  be stronger than  $(p, q)$ . It follows from Lemma 8 that  $(p, q) \in B_k$ . Now  $(\pi_{kn}p', \pi_{kn}q')$  is stronger than  $(\pi_{kn}p, \pi_{kn}q) = (p, q)$  and (\*) holds for  $\pi_{kn}p'$  in  $r_n$ . If  $p' <_t a_m$  then by Lemma 7  $\pi_{kn}p' <_n a_m$  and hence  $\pi_{kn}p' \perp_n \vec{a}$ . By Lemma 7 again,  $p' \perp_t \vec{a}$ .

Next, let  $\vec{a}$  and  $(p, q)$  be such that  $\vec{a} \in C_n - C_k$ ,  $(p, q) \in B_k - B_n$  and  $p \geq a_0$ . If  $k < n$ , we have  $\pi_{kn}p \geq p \geq a_0$  and we let  $F_t(\vec{a}, (p, q)) = F_n(\vec{a}, (\pi_{kn}p, \pi_{kn}q))$ . To verify (6), let  $m = F_t(\vec{a}, (p, q))$  and let  $(p', q') \in B_t$  be stronger than  $(p, q)$ . By Lemma 8  $(p', q') \in B_k$ , and  $(\pi_{kn}p', \pi_{kn}q')$  is stronger (in  $r_n$ ) than  $(\pi_{kn}p, \pi_{kn}q)$ . If  $p' <_t a_m$  then by Lemma 7  $\pi_{kn}p' <_n a_m$  and so  $\pi_{kn}p' \perp_n \vec{a}$ . By Lemma 7 again,  $p' \perp_t \vec{a}$ .

If  $k > n$ , we let  $F_t(\vec{a}, (p, q))$  be the least  $m$  such that  $a_m \notin E$  and that  $b_n^n \not\leq_n a_m$  (such  $m$  exists as  $\vec{a}$  does not have a lower bound in  $A_n$ ). To verify (6), let  $(p', q') \in B_t$  be stronger than  $(p, q)$ . If  $p' \in E$  then  $p' \not<_t a_m$  and if  $p' \notin E$  then by Lemma 7(i)

$p' \perp_t a_m$ . In either case, (6) is satisfied.

Finally, we define  $F_t(\vec{u}, (p, q))$ . Thus let  $(p, q) \in B_t$  be such that  $p \geq u_0$ . Since  $u_0 = a_0^0 \notin E$ , we have  $p \notin E$ . Let  $n$  be the  $n$  such that  $p \in A_n$ . We let  $F_t(\vec{u}, (p, q)) = 2n + 2$ . That is, the chosen  $u_m$  is  $u_{2n+2} = a_{n+1}^{n+1}$ . To verify (6), let  $(p', q') \in B_t$  be stronger than  $(p, q)$ . Since  $p \in A_n - E$ , by Lemma 8 we have  $(p', q') \in B_n$  and therefore  $p' \in A_n - E$ . But  $a_{n+1}^{n+1} \in A_{n+1} - E$  and so  $p' \not\leq_t a_{n+1}^{n+1}$ . Therefore (6) holds.

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