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## Non-trivial homeomorphisms of $\beta N \setminus N$ without the continuum hypothesis

by

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**Abstract.** The problem of constructing non-trivial homeomorphisms of  $\beta N \setminus N$  without assuming the continuum hypothesis is examined.

In [3] Shelah showed that it is consistent that all automorphisms of  $\mathcal{P}(\omega)/Finite$ , or, equivalently, all autohomeomorphisms of  $\beta N \setminus N$ , are trivial in the sense that they are induced by almost-permutations of the integers (an almost-permutation of  $\omega$  is an injective function from  $\omega$  to  $\omega$  whose domain and range are both cofinite). In [2] W. Rudin showed that the continuum hypothesis implies that there is a non-trivial autohomeomorphism by showing that there are in fact  $2^{2^{\aleph_0}}$  such homeomorphisms. It is the purpose of this paper to examine the question of how to construct non-trivial autohomeomorphisms in the absence of the continuum hypothesis. The reader should be warned that  $\beta N \setminus N$  and  $\mathcal{P}(\omega)/Finite$  will be used almost interchangeably. As well, subsets of the integers will routinely be confused with clopen sets in  $\beta N \setminus N$ .

At this point the reader may be wondering why the argument assuming  $2^{\aleph_0} = \aleph_1$  does not generalize to  $MA_{\aleph_1}$  and make the rest of this paper pointless. The reason, of course, is that an induction of length greater than  $\omega_1$  may run into a Hausdorff gap and stop. In fact it will be shown in [4] that PFA implies that all autohomeomorphisms of  $\beta N \setminus N$  are trivial and so this is consistent with  $MA_{\aleph_1}$ . This raises the following unanswered question:

**QUESTION.** Is it consistent with  $MA_{\aleph_1}$  that there is a non-trivial autohomeomorphism of  $\beta N \setminus N$ ?

The first result towards obtaining non-trivial autohomeomorphisms of  $\beta N \setminus N$  without the continuum hypothesis is due to Frolik [1]. He showed that the set of fixed points of any 1-1 continuous function from an extremally disconnected space to itself form a clopen set. To see how this can be used to construct non-trivial autohomeomorphisms of  $\beta N \setminus N$  consider the following lemma.

**LEMMA 1.** Suppose that  $\mathcal{I}$  is an ideal on  $\omega$  generated by an  $\subseteq^*$ -ascending sequence  $\{A_\alpha: \alpha \in \kappa\}$ . Suppose further that  $f_\alpha$  is an almost-permutation of  $A_\alpha$  for

each  $\alpha \in \kappa$  and that if  $\alpha \in \beta$  then  $f_\alpha \subseteq^* f_\beta$ . Then these functions  $\{f_\alpha: \alpha \in \kappa\}$  induce an isomorphism,  $\Phi$ , of the subalgebra of  $\mathcal{P}(\omega)/\text{Finite}$  generated by  $\mathcal{F}$  defined by

$$\Phi(X) = \begin{cases} f_\alpha'' X & \text{if there is some } \alpha \text{ such that } X \subseteq^* A_\alpha, \\ \omega \setminus f_\alpha''(\omega \setminus X) & \text{if there is some } \alpha \text{ such that } \omega \setminus X \subseteq^* A_\alpha. \end{cases}$$

*Proof.* Easy.

Now suppose that  $\{A_\alpha: \alpha \in \omega_1\}$  is a neighbourhood base of clopen sets for a  $p$ -point,  $\mathcal{F}$ , in  $\beta N \setminus N$ . It is easy to construct permutations  $f_\alpha$  of  $\omega \setminus A_\alpha$  such that if  $\alpha \in \beta$  then  $f_\alpha \subseteq^* f_\beta$  and such that homeomorphism of the clopen set  $A_\alpha \setminus A_{\alpha+1}$  induced by  $f_{\alpha+1}$  is not the identity. By Lemma 1 these permutations induce an autohomeomorphism  $\Phi$  of  $\beta N \setminus N$  (since the subalgebra generated by  $\mathcal{F}$  is all of  $\mathcal{P}(\omega)/\text{Finite}$ ). Note that  $\mathcal{F}$  is an isolated fixed point of  $\Phi$ . If  $\Phi$  was trivial then either it or its inverse would extend to a continuous 1-1 function from  $\beta N$  to  $\beta(N \setminus a)$  where  $a$  is finite. But then  $\mathcal{F}$  would still be an isolated fixed point contradicting the fact that, since  $\beta N$  is extremely disconnected, the set of fixed points must be clopen. Consequently, in any model where there is a  $p$ -point of character  $\aleph_1$  (and there are many such models where  $2^{\aleph_0} > \aleph_1$ ) there is a non-trivial autohomeomorphism of  $\beta N \setminus N$ . A similar proof due to Baumgartner appears in [5]. Recently van Douwen (unpublished) has shown that a point of character  $\aleph_1$  is sufficient to imply the existence of a non-trivial autohomeomorphism of  $\beta N \setminus N$ .

This raises the question of whether it is necessary to have points of small character in order to have non-trivial autohomeomorphisms. It will be shown that in the model obtained by adding  $\aleph_2$  Cohen reals to a model of the continuum hypothesis there is a non-trivial autohomeomorphism whereas it is known that in this model every point in  $\beta N \setminus N$  has character  $\aleph_2$ . The result for  $p$ -points of character  $\aleph_1$  will also be extended. It will be shown that if  $X$  is a closed  $p$ -set of character  $\aleph_1$  in  $\beta N \setminus N$  then the quotient space obtained by shrinking  $X$  to a point has a non-trivial autohomeomorphism. This is equivalent to saying that if  $\mathcal{S}$  is a  $p$ -ideal of character  $\aleph_1$  then there is an automorphism of the Boolean algebra generated by  $\mathcal{S}$  which is not induced by any function from  $\omega$  to  $\omega$ . Moreover, it will be shown that this automorphism is absolutely non-trivial in the sense that even in any  $\omega_1$ -preserving extension of the set-theoretic universe it is not induced by a function from  $\omega$  to  $\omega$ . The automorphisms constructed by using the method of Baumgartner and Frolik need not have this property.

The significance of this absolute non-triviality becomes apparent upon considering the method used in [3] to construct a model where all automorphisms of  $\mathcal{P}(\omega)/\text{Finite}$  are trivial. The construction consists of trapping non-trivial automorphisms and adding subsets of  $\omega$  to which it is impossible to extend the automorphism. One might wonder whether it is possible to obtain such a model by adding generic permutations which turn a non-trivial automorphism into a trivial one. The absolute non-triviality of the automorphisms to be constructed shows that this is not possible. The final point worth noting in this regard is the connection with uniformization properties. In [3] page 58 it is shown that  $MA_{\aleph_1}$  implies that if  $\{A_\alpha: \alpha \in \omega_1\}$  is a certain

type of almost disjoint family and  $f_\alpha: A_\alpha \rightarrow \omega$  are functions then there is  $F: \omega \rightarrow \omega$  such that for all  $\alpha \in \omega_1$   $F \upharpoonright A_\alpha \equiv^* f_\alpha$ . One might reasonably conjecture that a similar uniformization property is true for towers  $\{A_\alpha: \alpha \in \omega_1\}$  and functions  $f_\alpha: A_\alpha \rightarrow \omega$  so long as  $\alpha \in \beta$  implies  $f_\alpha \subseteq^* f_\beta$ . Again the example to be constructed shows that this is not possible.

**THEOREM 1.** *If  $\aleph_2$  Cohen reals are added to a model where  $2^{\aleph_0} = \aleph_1$  then there is a non-trivial automorphism of  $\mathcal{P}(\omega)/\text{Finite}$  in the resulting model.*

*Proof.* Let  $C_\gamma$  represent the partial order for adding  $\omega_1 \gamma$  Cohen reals. If  $G$  is  $C_{\omega_2}$  generic over a model of the continuum hypothesis then let  $V_\gamma = V[G \cap C_\gamma]$ . It will be shown by induction on  $\gamma \in \omega_2$  that there is an automorphism of  $(\mathcal{P}(\omega)/\text{Finite}) \cap \mathcal{V}_{\gamma+1}$ , which will be referred to as  $F_\gamma$ . Moreover, the automorphism  $F_\gamma$  will be constructed so that if  $\delta \in \gamma$  then  $F_\delta \subseteq F_\gamma$ . The desired automorphism will be  $\bigcup \{F_\gamma: \gamma \in \omega_2\}$ . The fact that is non-trivial will follow from the fact that each automorphism  $F_\gamma$  will be constructed so that there is no permutation of  $\omega$  in  $V_{\gamma+1}$  which induces  $F_\gamma$ .

To see how to construct the automorphisms notice that since  $V_1$  is a model of the continuum hypothesis it is easy to construct  $F_0$  inductively to satisfy the induction hypothesis. Now suppose that  $F_\gamma$  has been constructed on  $(\mathcal{P}(\omega)/\text{Finite}) \cap V_{\gamma+1}$ . To construct  $F_{\gamma+1}$  on  $(\mathcal{P}(\omega)/\text{Finite}) \cap V_{\gamma+2}$  proceed by induction on  $(\mathcal{P}(\omega)/\text{Finite}) \cap \mathcal{V}_{\gamma+2} \setminus V_{\gamma+1}$ . Let  $\{\pi_\xi: \xi \in \omega_1\}$  enumerate all almost-permutations of  $\omega$  and  $\{X_\xi: \xi \in \omega_1\}$  enumerate  $\mathcal{P}(\omega)/\text{Finite}$  in  $V_{\gamma+2}$ . Suppose that  $B$  is a countable subalgebra of  $\mathcal{P}(\omega)/\text{Finite}$  and that the automorphism  $F_{\gamma+1}$  has been defined on the algebra generated by  $((\mathcal{P}(\omega)/\text{Finite}) \cap V_{\gamma+1}) \cup B$  (call this algebra  $B'$ ). Let  $X_\xi \in (\mathcal{P}(\omega)/\text{Finite}) \setminus B'$  and let  $X$  be an equivalence class representative of  $X_\xi$ . Define

$$\mathcal{F}(X) = \{b \in B'; X \subseteq^* b\}, \quad \mathcal{S}(X) = \{b \in B'; X \supseteq^* b\}.$$

It will first be shown that  $\mathcal{F}(X)$  and  $\mathcal{S}(X)$  are countably generated and then this fact will be used to extend the automorphism.

Since  $B$  is countable, it suffices to show that  $\{b \in B' \cap V_{\gamma+1}; X \subseteq^* b\}$  and  $\{b \in B' \cap V_{\gamma+1}; X \supseteq^* b\}$  are countably generated. Let  $C$  be a countable completely embedded subalgebra of  $C_{\gamma+1}$  such that  $X$  has a  $C$  name. For each  $q \in C$  let  $X(q) = \{n \in \omega; q \text{ forces } "n \in N"\}$ . It will be shown that  $\{X(q); q \in C\}$  generate  $\{b \in B' \cap V_{\gamma+1}; X \subseteq^* b\}$  (a similar proof will work for  $\{b \in B' \cap V_{\gamma+1}; X \supseteq^* b\}$ ). Now suppose that  $p \in C_{\gamma+1}$  is a condition forcing that  $X \subseteq b \cup k$  for some integer  $k$ . Then let  $p^*$  be the projection of  $p$  on  $C$ . To see that  $b \subseteq X(p^*) \cup k$  suppose that  $m \in b \setminus k$ . If  $m \notin X(p^*)$  then there is some  $q \supseteq p^*$  such that  $q$  forces " $m \notin X$ ". Moreover, it may be assumed that  $q \in C$  since  $C$  is completely embedded in  $C_{\gamma+1}$ . Hence  $q \cup p$  is a condition which forces contradictory statements.

Now to extend  $F_{\gamma+1}$  it suffices to define  $F_{\gamma+1}(X_\xi) = Y_\xi$  where  $Y$  is an equivalence class representative of  $Y_\xi$  and the following conditions are satisfied:

1.  $Y_\xi \notin B'$  and  $Y \neq \pi_\xi(X)$ ;
2.  $F_{\gamma+1}(b) \supseteq^* Y_\xi \supseteq^* F_{\gamma+1}(c)$  for every  $b \in \mathcal{F}(X)$  and  $c \in \mathcal{S}(X)$ ;

3. if neither  $b$  nor  $\omega \setminus b$  is in  $\mathcal{F}(X) \cup \mathcal{S}(X)$  then  $Y \cap b$  and  $b \setminus Y$  are both infinite.

It is left to the reader to verify that extending  $F_{\gamma+1}$  to the algebra generated by  $B' \cup \{Y_i\}$  in the natural way is an isomorphism onto the algebra generated by  $(F_{\gamma+1}''B') \cup \{Y_i\}$ .

So it suffices to see why  $Y$  can be found. Since  $V_{\gamma+2}$  is obtained from  $V_{\gamma+1}$  by adding  $\omega_1$  Cohen reals and  $B$  is countable it is possible to find a Cohen generic real of  $B'$ . Notice that the partial order for splitting the gap formed by  $(F_{\gamma+1}''\mathcal{F}(X))$  and  $F_{\gamma+1}''\mathcal{S}(X)$  is a countable partial order since the gap is countably generated. Hence the Cohen real can be thought of as filling this gap. It is easy to check that such a Cohen real satisfies properties (1), (2) and (3). So if the induction is carried out and  $F_{\gamma+1}$  and  $F_{\gamma+1}^{-1}$  are dealt with alternately so as to make the limit function surjective then  $F_{\gamma+1}$  will have been defined as wanted.

The only thing left to consider is the limit stages. Limits of cofinality  $\omega_1$  take care of themselves. The limits of cofinality  $\omega$  are handled almost the same as the successor case; the only difference is that if  $\gamma$  has cofinality  $\omega$  then  $F_\gamma$  has been defined on

$$(\mathcal{P}(\omega)/\text{Finite}) \cap (\cup \{V_\delta; \delta \in \gamma\}).$$

Instead of defining  $F_{\gamma+1}$  on  $V_\gamma \setminus \cup \{V_\delta; \delta \in \gamma\}$  we define it on  $V_{\gamma+1} \setminus \cup \{V_\delta; \delta \in \gamma\}$  so that there will be enough Cohen reals to make the argument work.

**THEOREM 2.** *If  $\mathcal{S}$  is the dual of a  $p$ -filter of character  $\aleph_1$  then there is an automorphism of the Boolean algebra generated by  $\mathcal{S}$  which is not induced by any function from  $\omega$  to  $\omega$ . Moreover, this is upward absolute with respect to models preserving  $\omega_1$ . (The automorphism itself does extend canonically to the Boolean algebra generated by  $\mathcal{S}$  in any extensions of the universe.)*

**Proof.** Let  $\{A_\gamma; \gamma \in \omega\}$  be  $\ast$ -ascending. It may be assumed that  $A_\gamma \subseteq A_{\gamma+1}$  for each  $\gamma$ . Furthermore, if  $\gamma = \lambda + k$  where  $\lambda$  is a limit then it may be assumed that  $k \in A_\gamma$ . Hence if  $\nu$  is a limit then  $A_\nu \subseteq \cup \{A_{\sigma(n)}; n \in \omega\}$  for some increasing sequence  $\sigma$  approaching  $\nu$ . Transfinite induction on  $\omega_1$  will be used to define bijections  $f_\alpha: A_\alpha \rightarrow A_\alpha$  such that the following conditions are satisfied:

4.  $f_\alpha \circ f_\alpha = \text{id}$ ;
5. if  $\alpha \in \beta$  then  $f_\alpha \subseteq \ast f_\beta$ ;
6. if  $P_X(f, g)$  represents the statement

$$“\forall m \in \omega \setminus X \text{ (if } m \in \text{dom}(f) \cap \text{dom}(g) \text{ then } f(m) = g(m))”$$

then  $\{\beta \in \alpha; P_n(f_\beta, f_\alpha)\}$  is finite for each  $n$ . Condition (6) is reminiscent of Hausdorff's construction of an  $(\omega_1, \omega_1^*)$  gap. Finally, if  $f$  is a function and  $X \in [\omega]^{<\omega}$  then let  $f \upharpoonright X$  represent the restriction of  $f$  to  $\text{dom}(f) \setminus X$ . Lemma 1 and condition (5) assure that the functions  $\{f_\alpha; \alpha \in \omega_1\}$  will generate an automorphism  $\Phi$  of the Boolean algebra generated by  $\mathcal{S}$ .

Note that (6) implies that there is no  $g: \omega \rightarrow \omega$  such that  $g \supseteq \ast f_\beta$  for all  $\beta$  because otherwise there is some  $k$  such that

$$\{\beta \in \omega_1; g \supseteq f_\beta \upharpoonright k\}$$

is uncountable. Let  $\gamma$  belong to  $\{\beta \in \omega_1; g \supseteq f_\beta \upharpoonright k\}$  with infinitely many predecessors in this set. Now if

$$\alpha \in \{\beta \in \gamma; g \supseteq f_\beta \upharpoonright k\}$$

then  $(f_\alpha \upharpoonright k) \cup (f_\gamma \upharpoonright k) \subseteq g$  and hence it follows that  $P_k(f_\alpha, f_\gamma)$  holds contradicting (6). It follows that  $\Phi$  is non-trivial.

It will now be shown how to carry out the induction. Suppose that  $\{f_\alpha; \alpha \in \gamma\}$  have been constructed. If  $\gamma = \beta + 1$  then let

$$f_\gamma(m) = \begin{cases} m & \text{if } m \in A_\gamma \setminus A_\beta, \\ f_\beta(m) & \text{if } m \in A_\beta. \end{cases}$$

Notice that since  $A_{\beta+1} \supseteq A_\beta$  it follows that  $f_\gamma \supseteq f_\beta$  and hence (5) holds. To see that (6) holds note that for each  $k \in \omega$

$$\{\mu \in \beta; P_k(f_\mu, f_\gamma)\} \subseteq \{\mu \in \beta; P_k(f_\mu, f_\beta)\}$$

because  $f_\gamma \supseteq f_\beta$ .

If  $\gamma$  is a limit then let  $\{\sigma(n); n \in \omega\}$  be an increasing sequence cofinal in  $\gamma$  such that  $\cup \{A_{\sigma(n)}; n \in \omega\} \supseteq A_\gamma$ . Define by induction on  $\omega$  finite sets  $K_i \subseteq \omega$  and finite functions  $h_i$  from  $A_{\sigma(i)} \cap K_i$  to  $A_\gamma$  such that:

7.  $K_i$  is closed under  $f_{\sigma(i)}$ ;
8.  $i \subseteq K_i$ ;
9. if  $F_i = \cup \{h_j \cup (f_{\sigma(j)} \upharpoonright K_j); j \leq i\}$  then  $F_i$  is injective
10.  $\text{domain}(F_i) = \text{range}(F_i) \supseteq \cup \{A_{\sigma(j)} \cap A_\gamma; j \leq i\}$ ;
11.  $h_i \circ h_i = \text{id}$ ;
12. if  $\beta \in \sigma(i+1) \setminus \sigma(i)$  and both  $P_{K_i}(f_\beta, F_i)$  and  $P_{K_{i+1}}(f_\beta, f_{\sigma(i+1)})$  hold then  $((K_{i+1} \setminus K_i) \cap A_{\sigma(i+1)} \cap A_\beta) \setminus \text{range}(F_i) \neq \emptyset$ ;
13. if  $\beta \in \sigma(i+1) \setminus \sigma(i)$ ,  $P_{K_{i+1}}(f_\beta, f_{\sigma(i+1)})$  holds and  $j$  belongs to

$$((K_{i+1} \setminus K_i) \cap A_{\sigma(i+1)} \cap A_\beta) \setminus \text{range}(F_i)$$

then  $h_{i+1}(j) \neq f_\beta(j)$ ;

14. there is some  $j \in K_{i+1} \setminus K_i$  such that  $h_{i+1}(j) \neq f_{\sigma(i+1)}(j)$ .

If the induction can be carried out then let  $f_\gamma = \cup \{F_i; i \in \omega\}$ . The fact that  $f_\gamma$  is a bijection follows from (9) and (10) and the choice of the sequence  $\{\sigma(n); n \in \omega\}$ . The fact that (4) is satisfied follows from (11) and (5) follows from the construction. To see that (6) holds let  $n \in \omega$ . Then  $n \subseteq K_n$  and it suffices to show that

$$\{\beta \in \gamma; P_{K_n}(f_\beta, f_\gamma)\}$$

is finite. But note that

$$\{\beta \in \sigma(n); P_{K_n}(f_\beta, f_\gamma)\} \subseteq \{\beta \in \sigma(n); P_{K_n}(f_\beta, f_{\sigma(n)} \upharpoonright K_n)\} \\ = \{\beta \in \sigma(n); P_{K_n}(f_\beta, f_{\sigma(n)})\}$$

and the last set is finite by the induction hypothesis.

Hence, it suffices to show that if  $\beta \geq \sigma(n) + 1$  then  $P_{K_n}(f_\beta, f_\gamma)$  fails. First note that (14) ensures that  $P_{K_n}(f_\gamma, f_{\sigma(m)})$  fails for every  $m \geq n + 1$  and so it may be assumed that  $\sigma(m) < \beta < \sigma(m + 1)$ , where  $m \geq n$ . But now it follows that  $P_{K_{m+1}}(f_\beta, f_{\sigma(m+1)})$  holds if  $P_{K_n}(f_\beta, f_\gamma)$  does since  $K_n \subseteq K_{m+1}$ . Hence (12) ensures that either  $P_{K_m}(f_\beta, F_m)$  fails or  $((K_{m+1} \setminus K_m) \cap A_{\sigma(m+1)}) \cap A_\beta \setminus \text{range}(F_m) \neq \emptyset$ . But the first possibility implies that  $P_{K_m}(f_\beta, f_\gamma)$  fails and so the second must hold. Now an application of (13) yields that  $P_{K_m}(f_\beta, h_{m+1})$  fails, and hence so does  $P_{K_n}(f_\beta, f_\gamma)$ .

It now suffices to show that the induction can be carried out. To this end suppose that  $h_i$  and  $K_i$  have been constructed so that properties (8) to (14) hold. (To begin the induction simply choose  $K_0$  so that  $A_{\sigma(0)} \setminus K_0 \subseteq A_\gamma$  and  $(f_{\sigma(0)} \upharpoonright K_0)''(A_{\sigma(0)} \setminus K_0) = A_{\sigma(0)} \setminus K_0$  and  $h_0: K_0 \cap A_{\sigma(0)} \rightarrow A_\gamma \setminus A_{\sigma(0)}$  consists of 2-cycles.) Choose  $T$  to be finite and such that  $F_i \upharpoonright T \subseteq f_{\sigma(i+1)}$ ,  $T$  is closed under both  $F_i$  and  $f_{\sigma(i+1)}$ ,  $K_i \cup i \subseteq T$  and  $(f_{\sigma(i+1)} \upharpoonright T) \cup F_i$  is injective. Let

$$B = \{\beta \in (\sigma(i+1) \setminus \sigma(i)); P_T(f_\beta, f_{\sigma(i+1)})\}.$$

Then  $B$  is finite. Now choose  $K_{i+1} \supseteq T$  such that for each  $\beta \in B$  there is

$$j \in ((K_{i+1} \setminus T) \cap A_\beta \cap A_{\sigma(i+1)}) \setminus \text{range}(F_i)$$

and  $K_{i+1}$  is closed under  $f_{\sigma(i+1)}$ . As well, it may be assumed that  $(K_{i+1} \setminus K_i) \cap A_{\sigma(i+1)} \neq \emptyset$ . Now let  $C = \{\beta \in (\sigma(i+1) \setminus \sigma(i)); P_{K_{i+1}}(f_\beta, f_{\sigma(i+1)})\}$ . Once more the induction hypothesis implies that  $C$  is finite. Now define

$$h_{i+1}: K_{i+1} \rightarrow A_\gamma \setminus (\text{range}(F_i) \cup A_{\sigma(i+1)})$$

such that  $h_{i+1}$  consists of 2-cycles and  $h_{i+1}(j) \neq f_\beta(j)$  and  $h_{i+1}(j) \neq f_{\sigma(i+1)}$  for all  $\beta \in C$  and  $j \in K_{i+1}$ .

Except for (12) it is easy to verify that properties (7) to (14) all hold. To see that (12) holds suppose that  $\beta \in \sigma(i+1) \setminus \sigma(i)$ , and both  $P_{K_i}(f_\beta, F_i)$  and  $P_{K_{i+1}}(f_\beta, f_{\sigma(i+1)})$  hold. There are two possibilities. The first one is that  $P_T(f_\beta, f_{\sigma(i+1)})$  holds. In this case  $\beta \in B$  and so the choice of  $K_{i+1}$  ensures that (12) is satisfied. If, on the other hand,  $P_T(f_\beta, f_{\sigma(i+1)})$  fails then, since  $P_{K_{i+1}}(f_\beta, f_{\sigma(i+1)})$  does hold, there is some  $j \in K_{i+1} \setminus T$  such that  $f_\beta(j) \neq f_{\sigma(i+1)}(j)$ . Hence

$$j \in (K_{i+1} \setminus K_i) \cap A_{\sigma(i+1)} \cap A_\beta.$$

It suffices to show that  $j \notin \text{range}(F_i)$ . But since  $\text{range}(F_i \upharpoonright T) = \text{domain}(F_i \upharpoonright T)$  and  $F_i \upharpoonright T \subseteq f_{\sigma(i+1)}$  it follows that, if  $j \in \text{domain}(F_i)$ , then  $F_i(j) = f_{\sigma(i+1)}(j) \neq f_\beta(j)$ . This contradicts that  $P_{K_i}(f_\beta, F_i)$  holds.

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