

## Automorphism towers and automorphism groups of fields without Choice

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*Dedicated to Professor Rüdiger Göbel on his 70th birthday*

ABSTRACT. This paper can be viewed as a continuation of the article by Kaplan and Shelah that dealt with the automorphism tower problem without Choice. Here we deal with the inequality  $\tau_\kappa^{\text{nlg}} \leq \tau_\kappa$  without Choice and introduce a new proof to a theorem of Fried and Kollár that any group can be represented as an automorphism group of a field. The proof uses a simple construction: working more in graph theory, and less in algebra.

### 1. Introduction and preliminaries

**Background.** Although this paper hardly mentions automorphism towers, it is the main motivation for it. So we shall start by giving the story behind them.

Given any centerless group  $G$ ,  $G \cong \text{Inn}(G) \leq \text{Aut}(G)$  so we can embed  $G$  into its automorphism group. Also, an easy exercise shows that  $\text{Aut}(G)$  is also without center, so we can do this again, and again:

DEFINITION 1.1. For a centerless group  $G$ , we define *the automorphism tower*  $\langle G^\alpha \mid \alpha \in \text{ord} \rangle$  by

- $G^0 = G$ .
- $G^{\alpha+1} = \text{Aut}(G^\alpha)$ .
- $G^\delta = \cup \{G^\alpha \mid \alpha < \delta\}$  for  $\delta$  limit.

REMARK 1.2. The union in limit stages can be understood as the direct limit. But we shall think of the tower as an increasing continuous sequence of groups.

The natural question that arises, is whether this process stabilizes, and when. We define

DEFINITION 1.3. For such a group, define  $\tau_G = \min \{ \alpha \mid G^{\alpha+1} = G^\alpha \}$ .

In 1939, Weilandt proved in [Wie39] that for finite  $G$ ,  $\tau_G$  is finite. What about infinite  $G$ ? There exist examples of centerless infinite groups such that

1991 *Mathematics Subject Classification.* Primary 20A15; Secondary 12F12, 20E36.

*Key words and phrases.* Automorphism tower, axiom of choice, automorphism groups of fields.

The first author's research was partially supported by the SFB 878 grant.

The second author would like to thank the United States-Israel Binational Science Foundation for partial support of this research, Grant no. 2002323. Publication 913.

this process does not stop in any finite stage. For example — the infinite dihedral group  $D_\infty = \langle x, y \mid x^2 = y^2 = 1 \rangle$  satisfies  $\text{Aut}(D_\infty) \cong D_\infty$  while the automorphism replacing  $x$  with  $y$  is not in  $\text{Inn}(D_\infty)$ . The question remained open until the works of Faber [Fab78] and Thomas [Tho85, Tho98] (who was not aware of Faber's work), that showed  $\tau_G < (2^{|G|})^+$ .

DEFINITION 1.4. For a cardinal  $\kappa$  we define  $\tau_\kappa$  as the smallest ordinal such that  $\tau_\kappa > \tau_G$  for all centerless groups  $G$  of cardinality  $\leq \kappa$ , or in other words

$$\tau_\kappa = \bigcup \{ \tau_G + 1 \mid G \text{ is centerless and } |G| \leq \kappa \}.$$

Since  $(2^\kappa)^+$  is regular we can immediately conclude  $\tau_\kappa < (2^\kappa)^+$ .

This paper is concerned with a Choiceless universe, i.e. we do not assume the axiom of Choice. As a consequence, the previous definition is generalized to

DEFINITION 1.5. For a set  $k$ , we define  $\tau_{|k|}$  to be the smallest ordinal  $\alpha$  such that  $\alpha > \tau_G$  for all groups  $G$  with power  $\leq |k|$ .

Note that when we write  $|X| \leq |Y|$  as in the definition above, we mean that there is an injective function from  $X$  to  $Y$ . Below we provide a short glossary.

A helpful and close notion is that of *the normalizer tower*  $\langle \text{nor}_G^\alpha(H) \mid \alpha \in \mathbf{ord} \rangle$  of a subgroup  $H$  of  $G$  in  $G$ .

DEFINITION 1.6. Let

- $\text{nor}_G^0(H) = H$ .
- $\text{nor}_G^{\alpha+1}(H) = \text{nor}_G(\text{nor}_G^\alpha(H))$ .
- $\text{nor}_G^\delta(H) = \bigcup \{ \text{nor}_G^\alpha(H) \mid \alpha < \delta \}$  for  $\delta$  limit.

And we let the normalizer length be  $\tau_{G,H}^{\text{nlg}} = \min \{ \alpha \mid \text{nor}_G^{\alpha+1}(H) = \text{nor}_G^\alpha(H) \}$  (sometimes we just write  $\tau_{G,H}$ ).

Analogously to  $\tau_\kappa$ , we define

DEFINITION 1.7. For a cardinal  $\kappa$ , let  $\tau_\kappa^{\text{nlg}}$  be the smallest ordinal such that  $\tau_\kappa^{\text{nlg}} > \tau_{\text{Aut}(\mathfrak{A}),H}$ , for every structure  $\mathfrak{A}$  of cardinality  $\leq \kappa$  and  $H \leq \text{Aut}(\mathfrak{A})$  of cardinality  $\leq \kappa$ .

In general (i.e. without assuming Choice), for a set  $k$ , we define  $\tau_{|k|}^{\text{nlg}}$  as the smallest ordinal  $\alpha$ , such that for every structure  $\mathfrak{A}$  of power  $||\mathfrak{A}|| \leq |k|$ ,  $\tau_{\text{Aut}(\mathfrak{A}),H} < \alpha$  for every subgroup  $H \leq \text{Aut}(\mathfrak{A}) = G$  of power  $|H| \leq |k|$ . In other words,  $\tau_{|k|}^{\text{nlg}} = \sup \{ \tau_{G,H} + 1 \mid \text{for such } G, H \}$ .

In [JST99, Lemma 1.8], Just, Shelah and Thomas proved the following inequality

$$\tau_\kappa \geq \tau_\kappa^{\text{nlg}}.$$

In fact it was essentially already proved by Thomas in [Tho85].

In [KS09] we dealt with an upper bound of  $\tau_\kappa$  without assuming Choice. Here we prove  $\tau_\kappa \geq \tau_\kappa^{\text{nlg}}$  without Choice, and also provide a Choiceless variant of  $\tau_{|k|} \geq \tau_{|k|}^{\text{nlg}}$ .

It is worth mentioning some previous results regarding  $\tau_\kappa$  that were proved using this inequality.

In [Tho85], Thomas proved that  $\tau_\kappa \geq \kappa^+$ . It is easy to conclude from Main Theorem A below that this result still holds without Choice. We will elaborate in the end of this section (See Corollary 2.5).

In [JST99] the authors found that for uncountable  $\kappa$  one cannot find an explicit upper bound for  $\tau_\kappa$  better than  $(2^\kappa)^+$  in  $ZFC$  (using set theoretic forcing). In [She07], Shelah proved that if  $\kappa$  is strong limit singular of uncountable cofinality then  $\tau_\kappa > 2^\kappa$  (using results from PCF theory). In the proofs the authors construct normalizer towers to find lower bound for  $\tau_\kappa$ , but we did not check how much Choice was used.

It remains an open question whether or not there exists a countable centerless group  $G$  such that  $\tau_G \geq \omega_1$ .

**Description of paper.** As mentioned before, we wish to prove  $\tau_\kappa \geq \tau_\kappa^{\text{nlg}}$  without Choice. So we started by reading what was done in [JST99] (which is also described in detail in [Tho]).

The proof contains three parts:

- (1) Given some structure, code it in a graph (i.e. find a graph with the same cardinality and automorphism group).
- (2) Given a graph code it in a field. Now we have a field  $K$  with some subgroup  $H \leq \text{Aut}(K)$  such that  $|K| = |H| = \kappa$ .
- (3) Use some lemmas from group theory and properties of  $PSL(2, K)$  to find a centerless group whose automorphism tower coincides with the normalizer tower of  $H$  in  $\text{Aut}(K)$ .

Our first intention was to mimic this proof, and to prove some version of  $\tau_{|k|} \geq \tau_{|k|}^{\text{nlg}}$  (see definitions 1.7 and 1.1 above). To explain what we did prove, we need some notation:

DEFINITION 1.8. Let  $X$  be a set.

- (1)  $X^{<\omega}$  is the set of all finite sequences of members of  $X$ .
- (2)  $[X]^{<\aleph_0} = \{a \subseteq X \mid |a| < \aleph_0\}$ .
- (3)  $X^{\langle <\omega \rangle} = [X^{<\omega}]^{<\aleph_0}$ , i.e. the set of all finite subsets of finite sequences of elements of  $X$ .

Our methods cannot tackle  $\tau_{|k|} \geq \tau_{|k|}^{\text{nlg}}$  without Choice, since one often needs to code finite sequences. The natural way to overcome this is to replace  $k$  with  $k^{<\omega}$ , so that we get  $\tau_{|k^{<\omega}|} \geq \tau_{|k^{<\omega}|}^{\text{nlg}}$ . However, we managed to prove a slightly different version:

MAIN THEOREM A. For any set  $k$ ,  $\tau_{|k^{\langle <\omega \rangle}|} \geq \tau_{|k^{\langle <\omega \rangle}|}^{\text{nlg}'}$ .

Where  $\tau_{|k|}^{\text{nlg}'}$  is a variant of  $\tau_{|k|}^{\text{nlg}}$ . See Definition 2.2 below.

With Choice there is no difference, and moreover, we get as a corollary the original inequality for a cardinal  $\kappa$  (see Corollary 2.4 below). It is a matter of taste whether replacing  $k^{<\omega}$  and nlg by  $k^{\langle <\omega \rangle}$  and nlg' matters. Still, one may ask whether  $\tau_{|k^{<\omega}|}^{\text{nlg}} \leq \tau_{|k^{<\omega}|}$  or even  $\tau_{|k|} \geq \tau_{|k|}^{\text{nlg}}$  holds without Choice.

Part (1) was easy enough. However, it needs a passage to a structure with countable language. This stage uses Choice. In order to fix this, we just bypassed the problem all together and replaced  $\tau_{|k|}^{\text{nlg}}$  by  $\tau_{|k|}^{\text{nlg}'}$ .

Part (3) was easy as well: An algebraic lemma which obviously did not need Choice (Lemma 4.1); And two lemmas regarding  $PSL(2, K)$  — Lemma 4.4 and

Lemma 4.8. The latter is a theorem of Van der Waerden and Schreier which described  $\text{Aut}(PSL(2, K))$ . There is a simple model theoretic argument that shows that these lemmas do not require Choice (Lemma 4.5).

However, part (2) seemed to be somewhat harder. In [JST99], the authors referred to the work of Fried and Kollár [FK82]. In [Tho], the author gives a less technical proof that the construction in [FK82] works. The proof, in both cases, was a little bit complicated, and we were suspicious that Choice was used in it. After some time we realized that it is most likely not used, but by then we already came up with a proof of our own, in which the construction of the field is much simpler, and thought that it is worth presenting. So, for part (2) we prove:

**MAIN THEOREM B.** *Let  $\Gamma = \langle X, E \rangle$  be a connected graph. Then for any choice of characteristic there exists a field  $K_\Gamma$  of that characteristic such that  $|K_\Gamma| \leq |X^{<\omega}|$  and  $\text{Aut}(K_\Gamma) \cong \text{Aut}(\Gamma)$ .*

The proof of Main Theorem B is given in Section 6. Here we will give a brief outline of the construction.

The plan was this: work a little bit on the graph, so that the algebra would be easier. First code the given graph as a graph with the following properties: its edges are colorable with some finite number  $N$  of colors, and the subgraphs induced by any particular color is a union of disjoint stars. This is done in Lemma 6.4.

Now the construction of the field is as follows: first let  $\langle p_0, p_1, \dots, p_N \rangle$  be a list of distinct odd primes. Start with  $\mathbb{Q}$  (or any prime field), and add the set of vertices  $X$  as transcendental elements over it. For each one, add  $p_0^n$  roots to it for all  $n < \omega$ . Now, for each edge,  $e = \{s, t\}$ , colored with the color  $l < N$ , adjoin  $p_{l+1}^n$  roots for all  $n < \omega$  to  $(s + t + 1)$ . This is it. The reader is invited to compare to [FK82].

This construction can be done without Choice.

In the proof we use a generalized form of a lemma by P. Pröhle that appears in [Prö84]. In their original paper, Fried and Kollár could construct  $K_\Gamma$  with the restriction that  $\text{char}(K_\Gamma) \neq 2$  and Pröhle removed this restriction. His “third lemma” from [Prö84] seemed to be perfect for our situation. However, we needed to generalize it in order to suit our purposes (and prove the generalization). This is Lemma 6.8.

The proof of Lemma 6.8 can be found in full detail in [KS11] which is an online copy of the present paper, but with the added proof of this technical lemma. We felt that the details of the proof can be omitted since it is technical and similar to the proof in [Prö84].

**Acknowledgment.** We would like to thank the referee for many useful remarks and to Haran Pilpel for drawing a graph with certain properties in record time.

**A note about reading this paper.** If the reader is not interested in Choice, but still wants to see the proof of Main Theorem A and Main Theorem B, he should ignore all the computations of cardinalities, since they become trivial. Also, with Choice, the construction of the field is somewhat easier — in our construction, we took the polynomial ring  $\mathbb{Q}[Y]$  (where  $Y$  is a set containing the vertices) and then the quotient by an ideal. Then we had to show the ideal is prime in order to take the field of fractions. But with Choice we can construct the field by adding roots from the algebraic closure. See also Remark 6.14.

### A small glossary.

- $|X| \leq |Y|$  means: There is an injective function from  $X$  to  $Y$ .
- $|X| = |Y|$  means: There is a bijection from  $X$  onto  $Y$ .
- For a structure  $\mathfrak{A}$ ,  $|\mathfrak{A}|$  is its universe and  $\|\mathfrak{A}\|$  is its cardinality.
- $\mathbb{V}$  is the universe and  $\mathbb{L}$  is Gödel's constructible universe.

### 2. A variant of $\tau_{|k|}^{\text{nlg}}$ and some corollaries of Main Theorem A

DEFINITION 2.1. A structure  $\mathfrak{A}$  is called rigid if  $\text{Aut}(\mathfrak{A}) = 1$ , i.e. it has no non-trivial automorphism.

DEFINITION 2.2. For a set  $k$ , we define  $\tau_{|k|}^{\text{nlg}'}$  as the smallest ordinal  $\alpha$  which is greater than  $\tau_{\text{Aut}(\mathfrak{A}), H}$  where  $\mathfrak{A}, H$  are as in Definition 1.7 and in addition the vocabulary (language)  $L$  of  $\mathfrak{A}$  satisfies

- (1) There is some rigid structure with universe  $L$  and a countable vocabulary (for instance,  $L$  is well-orderable); and
- (2)  $|L| \leq \left| |\mathfrak{A}|^{<\omega} \right|$ .

REMARK 2.3. If  $\kappa$  is a cardinal number (i.e. an  $\aleph$ ), then  $\tau_{\kappa}^{\text{nlg}} = \tau_{|\kappa^{<\omega}|}^{\text{nlg}'}$  and  $\tau_{\kappa} = \tau_{|\kappa^{<\omega}|}$ . This is true since  $|\kappa^{<\omega}| = |\kappa|$ , and because given any  $\mathfrak{A}$  as in the definition, we may assume that  $|\mathfrak{A}| \subseteq \kappa$  and that  $L$  is  $|\mathfrak{A}|^{<\omega} \subseteq \kappa^{<\omega}$  which is well-orderable (see [KS09, Observation 2.3]).

Hence, by Main Theorem A

COROLLARY 2.4. (ZF) For a cardinal  $\kappa$ ,  $\tau_{\kappa}^{\text{nlg}} \leq \tau_{\kappa}$ .

The following is another easy conclusion of Main Theorem A

COROLLARY 2.5. (ZF) for any cardinal  $\kappa$ ,  $\tau_{\kappa} \geq \kappa^+$ . Moreover, letting  $v_{k^{<\omega}}$  be the smallest nonzero ordinal  $\alpha$  such that there is no injective function  $f: \alpha \rightarrow k^{<\omega}$ , then  $\tau_{|k^{<\omega}|} \geq v_{k^{<\omega}}$  for any set  $k$ .

PROOF. By [Tho85], we know that this result is true with Choice. Moreover, he proves that  $\tau_{\kappa}^{\text{nlg}} \geq \kappa^+$  (see Lemma in the proof of Theorem 2 there). Let  $\alpha < v_{k^{<\omega}}$  be some ordinal. We know that  $\mathbb{L} \models \tau_{|\alpha|}^{\text{nlg}} \geq |\alpha|^+ > \alpha$  and that  $|\alpha| \leq k^{<\omega}$ .

For a moment we work in  $\mathbb{L}$ . So there is a group  $G$  (the automorphism group of some structure) and a subgroup  $H \leq G$  such that  $|H| \leq |\alpha|$  and  $\alpha \leq \tau_{G, H}$ . We may assume that  $|G| \leq |\alpha|$ . For one reason, this is the way it is constructed in [Tho85]. However, we give a self-contained explanation:

Let  $L$  be the language  $\{P, Q, <, R\} \cup L_{\text{Groups}}$  where  $P, Q$  are predicates,  $<, R$  are binary relation symbols and  $L_{\text{Groups}}$  is the language of groups. Consider the  $L$ -structure  $\mathfrak{G}$  with universe the disjoint union of  $G$  and  $\alpha$  where  $P^{\mathfrak{G}} = G$ ,  $Q^{\mathfrak{G}} = \alpha$ , with the group structure on  $P$ , the order on  $Q$  and  $R^{\mathfrak{G}}(x, \beta)$  holds iff  $x \in \text{nor}_G^{\beta}(H)$ . Let  $\mathfrak{G}' \prec \mathfrak{G}$  be an elementary substructure of size  $\leq |\alpha|$  such that  $H \subseteq P^{\mathfrak{G}'}$ ,  $\alpha \subseteq Q^{\mathfrak{G}'}$  (so  $\alpha = Q^{\mathfrak{G}'}$ ), and let  $G' = P^{\mathfrak{G}'}$ . As a group  $G'$  is a subgroup of  $G$  containing  $H$  of size  $\leq |\alpha|$  and for all  $\beta < \alpha$ ,  $\text{nor}_{G'}^{\beta}(H) \neq \text{nor}_{G'}^{\beta+1}(H)$ , and in particular  $\alpha \leq \tau_{G', H}$ .

Now we go back to  $\mathbb{V}$ , so  $|G| \leq |\alpha| \leq |k^{<\omega}|$  by assumption. By [KS09, Claim 2.8],  $\alpha \leq \tau_{G,H}^I = \tau_{G,H}^V$ . Let  $\mathfrak{A}$  be the structure with universe  $G$  and for each  $g \in G$  a unary function  $f_g$  taking  $x$  to  $x \cdot g$ . Then  $\text{Aut}(\mathfrak{A}) \cong G$ . So we conclude that  $\tau_{k^{<\omega}}^{\text{nlg}'} \geq \alpha$  (because  $G$  is well-orderable as in Remark 2.3 above). By Main Theorem A,  $\tau_{|k^{<\omega}|} \geq \alpha$ .  $\square$

### 3. Coding structures as graphs

The next lemma allows us to present any automorphism group of an (almost) arbitrary structure as the automorphism group of a graph.

LEMMA 3.1. *Let  $\mathfrak{A}$  be a structure for the vocabulary (=language)  $L$  such that*

- (1) *There is some rigid structure on  $L$  with vocabulary  $L'$  such that  $|L'| \leq \aleph_0$ .*
- (2)  $|L| \leq \left| |\mathfrak{A}|^{<\omega} \right|$ .

Then there is a structure  $\mathfrak{B}$  with vocabulary  $L_{\mathfrak{B}}$  such that

- $|\mathfrak{B}| \leq |\mathfrak{A}| + |L|$  (so  $\leq \left| |\mathfrak{A}|^{<\omega} \right|$ )
- $\text{Aut}(\mathfrak{B}) \cong \text{Aut}(\mathfrak{A})$
- $|L_{\mathfrak{B}}| = \aleph_0$

PROOF. We may assume that both  $L$  and  $L'$  are relational languages.

Define  $\mathfrak{B}$  by:

- $|\mathfrak{B}| = |\mathfrak{A}| \times \{0\} \cup L \times \{1\}$ .
- The vocabulary is  $L_{\mathfrak{B}} = \{R_n \mid n \in \omega\} \cup L' \cup \{P\}$  where  $P$  is a unary predicate and each  $R_n$  is an  $n+1$  place relation.

Where:

- $Q^{\mathfrak{B}} = Q^L$  on  $L \times \{1\}$  for each  $Q \in L'$ .
- $R_n^{\mathfrak{B}} = \left\{ ((a_0, 0), \dots, (a_{n-1}, 0), (R, 1)) \mid \begin{array}{l} R \in L \text{ is an } n \text{ place relation and} \\ (a_0, \dots, a_{n-1}) \in R^{\mathfrak{A}} \end{array} \right\}$
- $P^{\mathfrak{B}} = L \times \{1\}$ .

It is easy to see that  $\mathfrak{B}$  is as desired.  $\square$

This is well known:

THEOREM 3.2. *Let  $\mathfrak{A}$  be a structure for the first order language  $L$  which is as in the conditions of 3.1. Then there is a connected graph  $\Gamma = \langle X_{\Gamma}, E_{\Gamma} \rangle$  such that  $\text{Aut}(\Gamma) \cong \text{Aut}(\mathfrak{A})$ , and  $|X_{\Gamma}| \leq \left| |\mathfrak{A}|^{<\aleph_0} \right|$ .*

PROOF. For details see e.g. [Tho, Lemma 4.2.2] or [Hod93, Theorem 5.5.1]. From the construction (which does not use Choice) described there, one can deduce the part regarding the cardinality. The proof uses the fact that we can reduce to structures with countable languages, but this is exactly Lemma 3.1.  $\square$

### 4. Some group theory

LEMMA 4.1. *Let  $S$  be a simple non-abelian group, and let  $G$  be a group such that  $\text{Inn}(S) \leq G \leq \text{Aut}(S)$ . Then the automorphism tower of  $G$  is naturally isomorphic to the normalizer tower of  $G$  in  $\text{Aut}(S)$ .*

The proof of this lemma can be found in [Tho, Theorem 4.1.4] (and, of course, it does not use Choice).

So we need a simple group. Recall

DEFINITION 4.2. Let  $K$  be a field,  $n < \omega$ , then:

- $GL(n, K)$  is the group of invertible  $n \times n$  matrices over  $K$ .
- $PGL(n, K) = GL(n, K) / Z(GL(n, K))$  (Here,  $Z(GL(n, K))$  is the group  $K^\times \cdot I$  where  $I$  is the identity matrix).
- $SL(n, K) = \{x \in GL(n, K) \mid \det(x) = 1\}$ .
- $PSL(n, K) = SL(n, K) / Z(SL(n, K))$  (The denominator is just  $Z(GL(n, K)) \cap SL(n, K)$ ).

FACT 4.3.  $PSL(n, K)$  is a normal subgroup of  $PGL(n, K)$ .

LEMMA 4.4.  $PSL(2, K)$  is simple for any field  $K$  such that  $|K| \geq 3$ .

The proof of this lemma can be found in many books, e.g. [Rot95]. It is also true in  $ZF$ , by the following Lemma and Claim:

LEMMA 4.5. Suppose  $P$  is a claim, such that  $ZFC \vdash P$ , and  $\psi$  is a first order sentence (in some language) such that  $ZF \vdash 'P$  is true iff  $\psi$  does not have a model'. Then  $ZF \vdash P$ .

PROOF. If we have a model  $\mathbb{V}$  of  $ZF$ , such that  $\mathbb{V} \models \neg P$ , then  $\psi$  has a model so cannot prove contradiction (there is no use of Choice here). Hence  $\psi$  is consistent in  $\mathbb{L} = \mathbb{L}^{\mathbb{V}}$  as well. (If  $\psi$  was not consistent in  $\mathbb{L}$ , then a proof of a contradiction from  $\psi$  would exist in  $\mathbb{V}$  as well). Hence, by Gödel Completeness Theorem in  $ZFC$ ,  $\mathbb{L} \models \neg P$ , but  $\mathbb{L} \models ZFC$  — a contradiction.  $\square$

CLAIM 4.6. There is a first order sentence  $\psi$  such that  $\psi$  has a model iff there is a field  $K$ ,  $|K| \geq 3$  such that  $PSL(2, K)$  is not simple.

PROOF. Let  $L$  be the language of fields with an extra 4-ary relation  $H$ , i.e.  $L = \{+, \cdot, 0, 1, H\}$ . Let the sentence  $\psi$  say that the universe is a field  $K$  of size  $\geq 3$  and that  $H \subseteq K^4$  is a normal subgroup of  $SL(2, K)$  (after some choice of coordinates), and that  $H$  contains  $Z(SL(2, K))$  and also some element outside  $Z(SL(2, K))$ .  $\square$

We close this section by showing one final algebraic fact holds over  $ZF$ . Recall:

DEFINITION 4.7. Given any two groups  $N$  and  $H$  and a group homomorphism  $\varphi : H \rightarrow \text{Aut}(N)$ , we denote by  $N \rtimes_{\varphi} H$  (or simply  $N \rtimes H$  if  $\varphi$  is known) the semi-direct product of  $N$  and  $H$  with respect to  $\varphi$ .

Note that for a field  $K$ , there are canonical homomorphisms  $\text{Aut}(K) \rightarrow \text{Aut}(PSL(2, K))$  and  $\text{Aut}(K) \rightarrow \text{Aut}(PGL(2, K))$ .

FACT 4.8. (Van der Waerden, Schreier [vdWS28]) Let  $K$  be a field. Then every automorphism of  $PSL(2, K)$  is induced via conjugation by a unique element of  $PGL(2, K) := PGL(2, K) \rtimes \text{Aut}(K)$ . Hence  $\text{Aut}(PSL(2, K)) \cong PGL(2, K)$ .

This means that if  $\varphi \in \text{Aut}(PSL(2, K))$  then there are unique  $\alpha \in \text{Aut}(K)$  and  $g \in PGL(2, K)$  such that for every  $x \in PSL(2, K)$ ,  $\varphi(x) = g\alpha(x)g^{-1}$ .

We again use the model theoretic argument of Lemma 4.5 to give a proof of this fact in  $ZF$ :

CLAIM 4.9.

- (1) There is a first order sentence  $\psi$  such that  $\psi$  has a model iff there is a field  $K$ , and an automorphism  $\varphi \in \text{Aut}(PSL(2, K))$  such that  $\varphi$  is not in  $PGL(2, K)$ . (This implies the existence of  $(\alpha, g)$  required by the fact).
- (2) There is a first order sentence  $\psi'$  such that  $\psi'$  has a model iff there is a field  $K$ , and some  $1 \neq g \in PGL(2, K)$ ,  $\alpha \in \text{Aut}(K)$ , such that for every  $x \in PSL(2, K)$ ,  $\alpha(x) = gxg^{-1}$ . (This implies the uniqueness of  $(\alpha, g)$  required by the fact).

PROOF. (1): Let  $K$  be a field. Recall that  $x_t = \begin{pmatrix} 1 & t \\ 0 & 1 \end{pmatrix}$  and  $z_t = \begin{pmatrix} 1 & 0 \\ t & 1 \end{pmatrix}$  generate  $SL(2, K)$ . Let  $g \in PGL(2, K)$ ,  $\sigma \in \text{Aut}(PSL(2, K))$ .

Then  $\alpha \in \text{Aut}(K)$  satisfies  $\sigma(x) = g\alpha(x)g^{-1}$  iff the map  $x \mapsto g^{-1}\sigma(x)g$  takes  $\bar{x}_t$  to  $\bar{x}_{\alpha(t)}$  and  $\bar{z}_t$  to  $\bar{z}_{\alpha(t)}$ . Let  $L$  be the language of fields augmented with 4-place function symbols  $\{\sigma_i \mid i < 4\}$ .  $\psi$  says that the universe  $K$  is a field, and that  $\sigma$  is an automorphism of  $PSL(2, K)$  ( $SL(2, K)$  is a definable subset of  $K^4$ , as is  $Z(SL(2, K))$ ), such that for all  $g \in PGL(2, K)$ , the maps  $t \mapsto g^{-1}\sigma(\bar{x}_t)g$  and  $t \mapsto g^{-1}\sigma(\bar{z}_t)g$  do not induce a well defined automorphism of  $K$ .

(2): Let  $L$  be the language of fields.  $\psi'$  says that the universe  $K$  is a field and that there is some nontrivial  $g \in PGL(2, K)$  such that the maps  $t \mapsto g^{-1}\bar{x}_t g$  and  $t \mapsto g^{-1}\bar{z}_t g$  are induced by an automorphism  $\alpha$  of  $K$ .  $\square$

## 5. Proof of Main Theorem A from Main Theorem B

From Main Theorem B which is proved in the next section, we can now deduce

MAIN THEOREM A. For any set  $k$ ,  $\tau_{|k^{< \omega}|}^{\text{nlg}'} \leq \tau_{|k^{< \omega}|}$ .

PROOF. (essentially the same proof as in [JST99]). We are given a structure  $\mathfrak{A}$ , with language  $L$  such that on the set  $L$  there is a rigid structure with countable vocabulary, and  $|\mathfrak{A}| \leq |k^{< \omega}|$ . By Theorem 3.2 and Main Theorem B we may assume that  $\mathfrak{A}$  is an infinite field,  $K$ . We are also given a subgroup  $H \leq \text{Aut}(K)$ ,  $|H| \leq |k^{< \omega}|$ .

Let  $G = PGL(2, K) \rtimes H$ . Obviously  $|G| \leq |k^{< \omega}|$ .

$G$  is centerless, because by Fact 4.8, the centralizer of  $PSL(2, K)$  in  $PGL(2, K)$  is trivial, and  $PSL(2, K) \leq G$ . So  $PSL(2, K) \leq G \leq PGL(2, K)$ . By Lemmas 4.1, 4.4, and 4.8,  $G^\alpha$  is isomorphic to  $\text{nor}_{PGL(2, K)}^\alpha(G)$ .

Now, by induction on  $\alpha$ , one has  $\text{nor}_{PGL(2, K)}^\alpha(G) = PGL(2, K) \rtimes \text{nor}_{\text{Aut}(K)}^\alpha(H)$  and we are done.  $\square$

## 6. Coding graphs as fields

In the introduction we mentioned that the following theorem of Fried and Kollár [FK82] was used in [JST99]:

THEOREM 6.1. (Fried and Kollár) (ZFC) For every connected graph  $\Gamma$  there is a field  $K$  such that  $\text{Aut}(\Gamma) \cong \text{Aut}(K)$ , and  $|K| = |\Gamma| + \aleph_0$ .

Here we will offer a different proof of the Choiceless version, namely

MAIN THEOREM B. Let  $\Gamma = \langle X, E \rangle$  be a connected graph. Then there exists a field  $K_\Gamma$  of any characteristic such that  $|K_\Gamma| \leq |X^{< \omega}|$  and  $\text{Aut}(K_\Gamma) \cong \text{Aut}(\Gamma)$ .

**COROLLARY 6.2.** *If  $G$  is a group and there is some rigid structure with countable vocabulary on it, then there is a field  $K$  such that  $\text{Aut}(K) \cong G$ , and  $|K| \leq |G^{<\omega}|$ .*

**PROOF.** (of corollary) Let  $\mathfrak{A}$  be the structure with universe  $G$  and for each  $g \in G$  a unary function  $f_g$  taking  $x$  to  $x \cdot g$  so that  $\text{Aut}(\mathfrak{A}) \cong G$ . Now apply 3.2 and Main Theorem B.  $\square$

**6.1. Coding graphs as colored graphs.** We start by working a bit on the graph, to make the algebra easier.

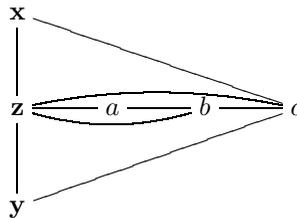
**DEFINITION 6.3.** A graph  $G = \langle X, E \rangle$  is called a star if there is a vertex  $v$  (the center) such that  $E \subseteq \{\{v, u\} \mid u \in V - \{v\}\}$ .

**LEMMA 6.4.** *There is some number  $N$  such that for every connected graph  $\Gamma = \langle X_\Gamma, E_\Gamma \rangle$ , there is a connected graph  $\Gamma^+ = \langle X_{\Gamma^+}, E_{\Gamma^+} \rangle$  with the following properties:*

- (1)  $\text{Aut}(\Gamma) \cong \text{Aut}(\Gamma^+)$ .
- (2) *There is a coloring  $C : E_{\Gamma^+} \rightarrow N$  of the edges of  $\Gamma^+$  in  $N$  colors such that for all  $l < N$  the  $l$ -th colored subgraph is a disjoint union of stars.*
- (3) *Every  $\varphi \in \text{Aut}(\Gamma^+)$  preserves the coloring.*
- (4)  $|X_{\Gamma^+}| \leq |X_\Gamma^{<\omega}|$ , in fact  $|X_\Gamma| \leq |X_{\Gamma^+}| \leq |X_\Gamma| + 4|E_\Gamma|$ .

**PROOF.** The idea is to replace each edge  $\{x, y\}$  by a copy of the graph  $G$  described below.

Recall that the valency of a vertex is the number of edges incident to the vertex, and will be denoted by  $\text{val}(x)$ . Let  $G = \langle X_G, E_G \rangle$  be the following auxiliary graph:



Note the following properties of  $G$ :

- It has only 2 automorphisms:  $\text{id}$  and  $\sigma$ , where  $\sigma$  switches  $\mathbf{x}$  and  $\mathbf{y}$ , but fixes all other vertices:  $\mathbf{z}, b, c$  are characterized by their valency and  $a$  is the only vertex with valency 2 which is adjacent to  $b, \mathbf{z}$ .
- $\mathbf{z}$  is adjacent to all the vertices, its valency is unique and is not divisible by  $\text{val}(\mathbf{x})$ .
- $\mathbf{x}$  and  $\mathbf{y}$  are not adjacent.

Description of  $\Gamma^+$ :

The set of vertices is

$$X_{\Gamma^+} = \{(1, x) \mid x \in X_\Gamma\} \cup \{(2, u, w) \mid u \in E_\Gamma, w \in X_G - \{\mathbf{x}, \mathbf{y}\}\}.$$

And the edges are:

- $(2, u, w)$  and  $(2, u', w')$  are adjacent iff  $u = u'$  and  $\{w, w'\} \in E_G$ .
- $(1, x)$  and  $(2, u, w)$  are adjacent iff  $x \in u$  and  $\{\mathbf{x}, w\} \in E_G$  (iff  $\{\mathbf{y}, w\} \in E_G$ ).

- That is all.

So, for each edge  $\{x, y\} = u \in E_\Gamma$  there is an induced subgraph  $\Gamma_u^+$  of  $\Gamma^+$ , whose vertices are  $\{(1, x), (1, y)\} \cup \{(2, u, w) \mid w \neq \mathbf{x}, \mathbf{y}\}$ , and  $\Gamma_{\{x, y\}}^+ \cong G$  (by sending  $\mathbf{x}$  to  $(1, x)$ ,  $\mathbf{y}$  to  $(1, y)$  and  $w \neq \mathbf{x}, \mathbf{y}$  to  $(2, u, w)$ ).

Let  $G'$  be the subgraph of  $G$  induced by removing  $\mathbf{y}$ , let  $N = |E_{G'}|$  (so  $N = 7$ ), and denote  $E_{G'} = \{e_0, \dots, e_{N-1}\}$ . Let  $f : \Gamma^+ \rightarrow G'$  be a homomorphism of graphs defined as follows:  $f(1, x) = \mathbf{x}$ ,  $f(2, u, w) = w$ . The coloring  $C : E_{\Gamma^+} \rightarrow N$  is defined by  $C(e) = i$  iff  $f(e) = e_i$ .

Let us now show (2). For each  $i < N$ , let  $\Gamma_i^+ = \langle X_i, E_i \rangle$  be the subgraph induced by the color  $i$ . If  $\mathbf{x} \notin e_i$ , then  $\Gamma_i^+$  is a union of disjoint edges by the definitions (and an edge is a star). If  $\mathbf{x} \in e_i$ , then  $\Gamma_i^+$  is a disjoint union of  $|X_\Gamma|$  stars, with centers  $\{(1, x) \mid x \in X_\Gamma\}$ , each having  $\text{val}_\Gamma(x)$  edges.

For (1), note that  $\text{val}_{\Gamma^+}(1, x) = \text{val}_G(\mathbf{x}) \cdot \text{val}_\Gamma(x)$  (or  $\infty$ , if  $\text{val}_\Gamma(x) \geq \aleph_0$ ), while  $\text{val}_{\Gamma^+}(2, u, w) = \text{val}_G(w)$ , hence  $\text{val}_{\Gamma^+}(2, u, \mathbf{z})$  is not divisible by  $\text{val}_{\Gamma^+}(1, x)$ .

Hence if  $\varphi \in \text{Aut}(\Gamma^+)$  then  $\varphi(2, u, \mathbf{z}) = (2, u', \mathbf{z})$  for some  $u' \in E_\Gamma$ . Since  $\mathbf{z}$  is adjacent to all the vertices in  $G$ ,  $\Gamma_{\{x, y\}}^+$  consists of all the vertices  $(2, u, \mathbf{z})$  is adjacent to and itself. So  $\varphi \upharpoonright \Gamma_u^+$  is an isomorphism onto  $\Gamma_{u'}^+$ . Since  $\text{Aut}(G) = \{\text{id}, \sigma\}$ , for all  $w \neq \mathbf{x}, \mathbf{y}$ ,  $\varphi(2, u, w) = (2, u', w)$ . This allows us to define  $\psi_\varphi = \psi \in \text{Aut}(\Gamma)$  by  $\psi(x) = x'$  where  $\varphi(1, x) = (1, x')$ . It is now easy to see that  $\varphi \mapsto \psi_\varphi$  is an isomorphism from  $\text{Aut}(\Gamma^+)$  onto  $\text{Aut}(\Gamma)$ .

(3) and (4) should be clear.  $\square$

**6.2. Coding colored graphs as fields.** Now we may assume that our graph is as in 6.4, and we start constructing the field.

We use the somewhat nonstandard notation of  $r$  as the characteristic of a field, so that  $\mathbb{F}_r$  is the prime field with  $r$  elements.

**DEFINITION 6.5.** Let  $F \subseteq K$  be a field extension.  $F$  is said to be relatively algebraically closed in  $K$  if every  $x \in K \setminus F$  is transcendental over  $F$ .

**DEFINITION 6.6.** Let  $p$  be a prime. An element  $x$  in a field  $F$  is called  $p$ -high, if there is a sequence  $\langle x_i \mid i < \omega \rangle$  of elements in  $F$ , such that  $x_0 = x$ , and  $x_{i+1}^p = x_i$ . With Choice this means that  $x$  has a  $p^n$ -th root for all  $n < \omega$ .

**EXAMPLE 6.7.** If  $F = \mathbb{Q}$ , then for  $p$  odd, the only  $p$ -high element in  $F$  are  $1, -1, 0$ . If  $F = \mathbb{F}_r$  for some prime  $r$ , then for every  $p$  such that  $(p, r-1) = 1$  (i.e. the map  $x \mapsto x^p$  is onto), every element in  $F$  is  $p$ -high.

This next lemma is the technical key. Its proof may use Choice, and this is OK, because we use it for finite  $\Gamma$  (see Remark 6.10 below).

**LEMMA 6.8.** (taken from [Prö84, The third lemma] with some adjustments) Let  $r$  be a prime number or 0,  $p$  a prime number different from  $r$  and let  $\{p_0, \dots, p_{n-1}\}$  be a set of pairwise distinct primes, different from  $p, r$ . Let  $F$  be a field of characteristic  $r$ . For  $k < n$ , let  $V_k$  be some set such that  $k \neq l \Rightarrow V_k \cap V_l = \emptyset$ , and let  $V = \bigcup_{k < n} V_k$ .

For each  $v \in V$ , let  $T_v \in F[X]$  be polynomials such that:

- none of them is constant.
- none of them is divisible by  $X$ .
- they are separable polynomials.
- they are pairwise relatively prime (i.e. no nontrivial common divisor).

Suppose that  $K$  is an extension of  $F$  generated by the set  $\{z_i \mid i < \omega\} \cup \{t_i^v \mid v \in V, i < \omega\}$  from the algebraic closure of  $F(z_0)$  where:

- $z_0$  is transcendental over  $F$ .
- $(z_{i+1})^p = z_i$  for all  $i < \omega$ .
- For  $v \in V$ ,  $t_0^v = T_v(z_0)$
- if  $v \in V_k$  then  $(t_{i+1}^v)^{p^k} = t_i^v$ .

Then we have the following properties:

- (1)  $F$  is relatively algebraically closed in  $K$ .
- (2) An equivalent definition of  $K$  is the following one: Suppose  $F$  is the field of fractions of an integral domain  $S$ . Then  $K$  is the field of fractions of the integral domain  $R/I$  (in particular  $I$  is prime) where  $R = S[Y_i, S_l^v \mid i, l < \omega, v \in V]$  (i.e. the ring generated freely by  $S$  and these elements) and  $I \leq R$  is the ideal generated by the equations:
  - (a)  $Y_{i+1}^p = Y_i$  for  $i < \omega$ .
  - (b)  $S_0^v = T_v(Y_0)$  for  $v \in V$ .
  - (c) If  $v \in V_k$ , then  $(S_{l+1}^v)^{p^k} = S_l^v$  for  $k < n, l < \omega$ .
- (3) Each  $q$ -high element of  $K$  belongs to  $F$  whenever  $q$  is a prime different from  $p$  and  $\langle p_k \mid k < n \rangle$ .
- (4) Each  $p$ -high element of  $K$  is of the form  $c \cdot (z_i)^m$ , where  $c$  is a  $p$ -high element of  $F$ ,  $i < \omega$  and  $m$  is an integer.
- (5) If  $p'$  is a prime different from  $p$  then  $z_0$  does not have a  $p'$  root.
- (6) If  $V$  is finite then  $|K| \leq |F^{(<\omega)}|$ . Furthermore, the injection witnessing this is definable from the parameters given when constructing  $K$  (i.e. the function  $v \mapsto T_v$ , etc).

The proof may be found in [KS11].

The rest of the section is devoted to proving

**THEOREM 6.9.** *Let  $\Gamma = \langle X, E, C \rangle$  be an  $N$ -colored graph as in Lemma 6.4. Then there exists a field  $K_\Gamma$  such that  $|K_\Gamma| \leq |X^{(<\omega)}|$  and  $\text{Aut}(K_\Gamma) \cong \text{Aut}(\Gamma)$ . Furthermore,  $X \subseteq K_\Gamma$  and  $\pi \mapsto \pi \upharpoonright X$  is an isomorphism from  $\text{Aut}(K_\Gamma)$  onto  $\text{Aut}(\Gamma)$ . We can choose  $K_\Gamma$  to be of any characteristic.*

So Main Theorem B immediately follows from this and Lemma 6.4.

The construction of  $K_\Gamma$ : Let  $L$  be the field  $\mathbb{Q}$  or  $\mathbb{F}_r$  for some prime  $r$ . Let  $\langle p_i \mid i \leq N \rangle$  list odd prime numbers which are different than  $r$ , and do not divide  $r - 1$  (so that in  $L$  there are no  $p_i$ -roots of unity). Let  $R$  be the ring  $L[Y_\Gamma]$  where  $Y_\Gamma = \{x_s^i \mid i < \omega, s \in X\} \cup \{x_e^i \mid i < \omega, e \in E\}$ <sup>1</sup> is an algebraically independent set. Let  $I_\Gamma \subseteq R$  be the ideal generated by the equations:

- $(x_s^{i+1})^{p_0} = x_s^i$  for all  $s \in X$  and  $i < \omega$ .
- If  $e = \{s, t\}$  then  $x_e^0 = x_s^0 + x_t^0 + 1$  for all  $s, t \in X$  and  $e \in E$ .
- If  $C(e) = l$  then  $(x_e^{i+1})^{p^{l+1}} = x_e^i$  for all  $e \in E$ .

Now let  $R_\Gamma$  be the ring  $R/I_\Gamma$ .

**REMARK 6.10.**

- (1) If  $\Gamma, \Gamma'$  are  $N$ -colored graphs, and  $\Gamma \cong \Gamma'$  (and the isomorphism respects the coloring) then  $R_\Gamma \cong R_{\Gamma'}$ .

<sup>1</sup>The  $i$  s are indices not exponents! Later we will use parentheses in order not to confuse a superscript with an exponent.

- (2) Hence we may use Choice when proving properties regarding  $R_\Gamma$  (and later  $K_\Gamma$ ) when  $\Gamma$  is finite because we may assume  $\Gamma \in \mathbb{L}$  (hence also  $R_\Gamma \in \mathbb{L}$  etc). In that case we may use Lemma 6.8 even if there is Choice in the proof.

PROPOSITION 6.11.  $I_\Gamma$  is prime, so we let  $K_\Gamma$  be the field of fractions of  $R_\Gamma$ .

The proof uses the following remark (when it makes sense)

REMARK 6.12. If  $\Gamma_0 \subseteq \Gamma_1$  are finite where  $\Gamma_i = \langle X_i, E_i, C_i \rangle$  for  $i < 2$  and  $X_1 = X_0 \cup \{t\}$ ,  $t \notin X_0$ , then the field extension  $K_{\Gamma_0} \subseteq K_{\Gamma_1}$  is as in Lemma 6.8, where

- $F$  is the field  $K_{\Gamma_0}$ ;  $r$  is its characteristic;  $p$  is  $p_0$ ;  $\{p_0, \dots, p_{n-1}\}$  is  $\{p_{l+1} \mid l < N\}$ ;  $V_k$  is the set of edges  $\{t, s\} \in E$  of color  $k$ ; for  $s \in X_0$  such that  $v = \{t, s\} \in E$ ,  $T_v$  is the polynomial  $X + x_s^0 + 1$ ;  $z_i$  is  $x_t^i$  and for  $v = e = \{t, s\}$ ,  $t_e^v$  is  $x_e^i$ .

PROOF. (of proposition) We may assume  $\Gamma$  is finite, so the proof is by induction on  $|X|$ . Suppose that  $\Gamma_0 \subseteq \Gamma_1$  where  $\Gamma_i = \langle X_i, E_i, C_i \rangle$  for  $i < 2$  and that  $X_1 = X_0 \cup \{t\}$ ,  $t \notin X_0$ . By induction,  $I_{\Gamma_0}$  is prime, so  $R = R_{\Gamma_0}$  is an integral domain.

Let  $Y_t = \{x_t^i \mid i < \omega\} \cup \{x_e^i \mid i < \omega, t \in e \in E_1\}$ ;  $I_t \subseteq R[Y_t]$  be the ideal generated by the equations related to  $t$  and  $\{e \in E_1 \mid t \in e\}$ .

By Lemma 6.8, clause (2),  $I_t$  is prime.

Consider the canonical projection  $\pi : L[Y_{\Gamma_1}] \rightarrow R[Y_t]$  so that  $\pi(I_{\Gamma_1}) = I_t$  and  $\langle I_{\Gamma_0} \rangle = \ker(\pi)$ . Hence,  $\pi$  induces an isomorphism  $L[Y_{\Gamma_1}]/I_{\Gamma_1} \rightarrow R[Y_t]/I_t$  and we are done since the right hand side is an integral domain.  $\square$

DEFINITION 6.13. (ZFC) Let  $F$  be a field and let  $p$  be a natural number. Let  $S$  be a set of elements from  $F$ . Then  $F(S, p)$  denotes the field which is obtained by adjoining the elements  $\{s(l) \mid s \in S, l < \omega\}$  from the algebraic closure of  $F$  where:

- $s(0) = s$ .
- $s(l+1)^p = s(l)$ ,  $l < \omega$ .

REMARK 6.14. Choice is a priori needed in this definition because the construction implicitly assumes the existence of an algebraic closure, and some ordering of  $S$  and of the  $p$ -roots of the  $s(l)$ s.

DEFINITION 6.15. Let  $K_{-1} = L(Y)(Y, p_0)$ , where  $Y = \{x_t^0 \mid t \in X\}$ , and  $L(Y)$  denotes the purely transcendental extension of  $L$ , and for  $l < N$ ,  $K_l = K_{l-1}(E_l, p_{l+1})$ , where  $E_l = \{x_s^0 + x_t^0 + 1 \mid \{s, t\} = e \in E, C(e) = l\}$ .

LEMMA 6.16.

- (1) For  $\Gamma$  finite<sup>2</sup>,  $K_\Gamma$  is canonically isomorphic to  $K_{N-1}$ .
- (2) If  $\Gamma_0 \subseteq \Gamma_1$  then  $K_{\Gamma_0} \subseteq K_{\Gamma_1}$ .

PROOF. (1) follows from Lemma 6.8, (2) by induction on the size of  $\Gamma$ , similarly to the proof of Proposition 6.11. (2) follows from (1) for finite  $\Gamma$ , which is enough.  $\square$

From now on, fix some  $\Gamma$ .

<sup>2</sup>The assumption that  $\Gamma$  is finite is only to insure that  $K_{N-1}$  is well defined, with Choice this assumption is not needed.

DEFINITION 6.17. For each  $Y \subseteq X$ , let  $\Gamma_Y$  be the induced subgraph generated by  $Y$  (i.e.  $\Gamma_Y = \langle Y, E \upharpoonright Y \rangle$ ) and let  $R_Y = R_{\Gamma_Y}$ ,  $K_Y = K_{\Gamma_Y}$ .

Some properties of  $K_\Gamma$ :

LEMMA 6.18. For  $\Gamma$  as in Lemma 6.4,

- (1) For any prime  $p$ , if  $a \in K_Y$  for some  $Y \subseteq X$  and is  $p$ -high in  $K_\Gamma$  then  $a$  is already  $p$ -high in  $K_Y$ .
- (2) For each  $i < \omega$ , the set  $\{x_s^i \mid s \in X\}$  is algebraically independent over  $L$ .
- (3) If  $X_1 \subseteq X_2$  then  $K_{X_1}$  is relatively algebraically closed in  $K_{X_2}$  (in particular  $L$  is r.a.c in  $K_\Gamma$ ).

PROOF. (1) and (2) follows from (3). For (3), we may assume  $X_1, X_2$  are finite, and then it is enough to prove it for the case  $X_2 = X_1 \cup \{t\}, t \notin X_1$ . Now use Remark 6.12, and clause (1) of Lemma 6.8.  $\square$

Now we shall define the isomorphism from  $\text{Aut}(\Gamma)$  to  $\text{Aut}(K_\Gamma)$ :

PROPOSITION 6.19. For  $\Gamma$  as in 6.4, there is a canonical injective homomorphism  $\sigma : \text{Aut}(\Gamma) \rightarrow \text{Aut}(K_\Gamma)$  defined by  $\sigma(\varphi)(x_t^i) = x_{\varphi(t)}^i$ , and  $\sigma(\varphi)(x_e^i) = x_{\varphi(e)}^i$ , for  $\varphi \in \text{Aut}(\Gamma)$  and all  $t \in X, e \in E$ .

PROOF.  $\sigma$  is well defined because of clause (3) of Lemma 6.4.  $\sigma$  is obviously a homomorphism. It is injective: If  $\sigma(\varphi) = \text{id}$ , while  $\varphi(s) = t \neq s$ , then  $x_s^0 = \sigma(\varphi)(x_s^0) = x_t^0$  — a contradiction to clause (2) of Lemma 6.18.  $\square$

Our aim is to prove that  $\sigma$  is onto. We start with:

CLAIM 6.20. Suppose that  $a \in K_\Gamma$  is  $p$ -high, then:

- (1) If  $p = p_0$  then  $a$  can be written in the form  $\varepsilon \cdot \prod \{(x_s^{n_s})^{m_s} \mid s \in X_0\}$  for some finite  $X_0 \subseteq X$ , some choice of  $m_s \in \mathbb{Z}, n_s < \omega$  for  $s \in X_0$  and a  $p_0$ -high element  $\varepsilon \in L$ .
- (2) If  $p = p_{l+1}$  for some  $l < N$  then  $a$  can be written in the form  $\varepsilon \cdot \prod \{(x_e^{n_e})^{m_e} \mid e \in E_0\}$  for some finite  $E_0 \subseteq E$  such that  $C \upharpoonright E_0 = l$ , some choice of  $n_e < \omega, m_e \in \mathbb{Z}$  for  $e \in E_0$  and a  $p_{l+1}$ -high element  $\varepsilon \in L$ .

PROOF. By Lemma 6.18, clause (1), there is some  $X_0 \subseteq X$  such that  $a$  is  $p$ -high in  $K_{X_0}$ . The proof is by induction on  $|X_0|$ . The base of the induction —  $X_0 = \emptyset$  — is clear. For the induction step, we prove that if  $X_0 \subseteq X_1$  are finite and  $X_1 = X_0 \cup \{t\}, t \notin X_0$ , and the claim is true for  $X_0$ , then every  $a \in K_{X_1}$  which is  $p$ -high has the desired form.

For clause (1), Remark 6.12 implies that we can use Lemma 6.8, clause (4).

For (2), we shall use the assumption on the coloring.

Case 1. There is no edge  $e_0 \ni t$  in  $\Gamma_{X_1}$  such that  $C(e_0) = l$ . In that case, we use clause (3) of Lemma 6.8, and conclude that  $a \in K_{X_0}$ .

Case 2. There is an edge  $e_0 \ni t$  in  $\Gamma_{X_1}$  with  $C(e_0) = l$ , but only one such edge. If  $e_0 = \{s, t\}, s \in X_0$  then  $x_{e_0}^0 = x_s^0 + x_t^0 + 1 \in K_{X_1}$  is transcendental over  $K_{X_0}$  (because  $x_t^0$  is). In addition  $x_t^0 = x_{e_0}^0 - x_s^0 - 1$  and for all vertices  $r \in X_0$  such that  $e_r = \{t, r\}$  is an edge (of some other color),  $x_{e_r}^0 = x_{e_0}^0 - x_s^0 + x_r^0$ . The polynomials  $X - x_s^0 - 1, X - x_s^0 + x_r^0$  satisfy the conditions of Lemma 6.8, and so, by clause (4),  $a$  is of the form  $(x_{e_0}^i)^m \cdot c$  for  $c$  which is  $p_{l+1}$ -high in  $K_{\Gamma_0}$  and we are done (we do

not use the lemma in the same way as in Remark 6.12 — here  $z_0$  is played by  $x_{e_0}^0$ , but it is the same idea).

*Case 3.* There is more than one edge  $e_0 \ni t$  in  $\Gamma_{X_1}$  with color  $l$ . Then  $t$  is the center of a star in the subgraph of  $\Gamma_1$  induced by that color. Assume that  $s_1, \dots, s_k \in X_0$  list the vertices such that  $C(s_i, t) = l$ , ( $k \geq 2$ ). Let  $X^- = X_0 \setminus \{s_1, \dots, s_k\}$ , and  $X' = X^- \cup \{t\}$ . Note that  $|X'| < |X_1|$ , so by the induction hypothesis, the claim is true for  $K_{X'}$ .  $\Gamma_{X_1}$  is built from  $\Gamma_{X'}$  by adding  $s_1, \dots, s_k$  and in each step we are in the previous case (because  $t$  was the center of a star), so we are done.  $\square$

LEMMA 6.21. *For all  $s \in X$ ,  $x_s^0$  does not have a  $p'$  root for  $p'$  a prime different from  $p_0$ .*

PROOF. Again, it is enough to prove this finite  $X_0 \subseteq X$ , and the proof is by induction on  $|X_0|$ , and follows from clause (5) of Lemma 6.8.  $\square$

This is the main proposition:

PROPOSITION 6.22. *Assume  $\varphi \in \text{Aut}(K_\Gamma)$  and that  $\{s_0, t_0\} \in E$  of color  $l$ . Then there is an edge  $\{s_1, t_1\} \in E$  of the same color such that  $\varphi(x_{s_0}^0) = x_{s_1}^0$  and  $\varphi(x_{t_0}^0) = x_{t_1}^0$ .*

PROOF. Let  $f_1 = \varphi(x_{s_0}^0)$ ,  $f_2 = \varphi(x_{t_0}^0)$ ,  $f = \varphi(x_{s_0}^0 + x_{t_0}^0 + 1) = f_1 + f_2 + 1$ . From Claim 6.20 it follows that

$$\begin{aligned} \bullet \quad & f_1 = \varepsilon_1 \cdot \prod \{(x_{s_s}^{i_s})^{m_s} \mid s \in X_0\}, \quad f_2 = \varepsilon_2 \cdot \prod \{(x_{t_t}^{i_t})^{m_t} \mid t \in Y_0\} \text{ and} \\ & f = \varepsilon_3 \cdot \prod \{(x_{e_e}^{i_e})^{m_e} \mid e \in E_0\}, \end{aligned}$$

where  $X_0, Y_0 \subseteq X$  and  $E_0 \subseteq E$  are finite nonempty;  $i_s < \omega$ ,  $m_s \in \mathbb{Z}$  for  $s \in X_0$ ;  $i_t < \omega$ ,  $m_t \in \mathbb{Z}$  for  $t \in Y_0$ ; and  $E_0$  is homogeneous of color  $l$  and  $i_e < \omega$ ,  $m_e \in \mathbb{Z}$  for  $e \in E_0$ . Let  $p = p_{l+1}$ , so  $f$  is  $p$ -high.

We can assume that unless  $i_s = 0$ ,  $p_0 \nmid m_s$  for  $s \in X_0 \cup Y_0$ , and that unless  $i_e = 0$ ,  $p \nmid m_e$  for  $e \in E_0$ .

Raising the equation  $f_1 + f_2 + 1 = f$  by  $p^k$  where  $k = \max\{i_e \mid e \in E_0\}$ , we have an equation of the form

$$\left( \varepsilon_1 \prod (x_s^{i_s})^{m_s} + \varepsilon_2 \prod (x_t^{i_t})^{m_t} + 1 \right)^{p^k} = \varepsilon_3^{p^k} \prod (x_r^0 + x_w^0 + 1)^{p^{k-i_{\{r,w\}} m_{\{r,w\}}}}.$$

Let  $i = \max\{i_t \mid t \in X_0 \cup Y_0\}$ . We can replace  $x_t^{i_t}$  by  $(x_t^i)^{p_0^{i-i_t}}$  and the same for  $x_s^{i_s}$ . Also replace  $x_r^0$  by  $(x_r^i)^{p_0^i}$  and the same for  $x_w^0$ . For  $t \in T := X_0 \cup Y_0 \cup \bigcup E_0$ , let  $y_t = x_t^i$ , then we get

$$\begin{aligned} \left( \varepsilon_1 \prod (y_s)^{p_0^{i-i_s} m_s} + \varepsilon_2 \prod (y_t)^{p_0^{i-i_t} m_t} + 1 \right)^{p^k} \\ = \varepsilon_3^{p^k} \prod \left( (y_r)^{p_0^i} + (y_w)^{p_0^i} + 1 \right)^{p^{k-i_{\{r,w\}} m_{\{r,w\}}}}. \end{aligned}$$

By Lemma 6.18, these elements are algebraically independent so this is an equation in the field of rational functions  $L(y_t \mid t \in T)$ .

The next step is to see that the exponents ( $m_t$  and  $m_{\{r,w\}}$ ) are non-negative. For that we use valuations.

Recall that for any field,  $F$  and any irreducible  $g \in F[X]$  there is a unique discrete (i.e. with value group  $\mathbb{Z}$ ) valuation on the field of rational functions  $F(t)$  defined by  $v(g(t)) = 1$ ,  $v \upharpoonright F^\times = 0$ . In this case,  $v \upharpoonright F[t] \geq 0$  and  $v(m(t)) > 0$  iff  $g|m$  for  $m \in F[X]$ . This is the  $g$ -adic valuation.

Suppose  $m_{t_0}$  is negative for some  $t \in X_0 \cup Y_0$ . Consider the discrete valuation  $v$  on the field  $L(y_t | t \in T)$  defined by  $v(y_{t_0}) = 1$ ,  $v \upharpoonright L(y_t | t \neq t_0)^\times = 0$ . Then on the left hand side we get  $v(LHS) < 0$  while on the right hand side,  $v(RHS) = 0$  — contradiction.

Suppose  $m_{\{r,w\}} < 0$  for some  $\{r,w\} \in E_0$ . Consider the valuation  $v$  on the field  $L(y_t | t \in T)$  defined by  $v(g(y_r)) = 1$ ,  $v \upharpoonright L(y_t | t \neq r)^\times = 0$  where  $g$  is any irreducible polynomial dividing  $X^{p_0^i} + (y_w)^{p_0^i} + 1$ . So  $v\left((y_r)^{p_0^i} + (y_w)^{p_0^i} + 1\right) > 0$ , while  $g$  does not divide  $\left(X^{p_0^i} + (y_{w'})^{p_0^i} + 1\right)$  for  $w \neq w'$  (they relatively prime) so  $v(RHS) < 0$ . On the other hand, since  $v(y_r) = 0$ ,  $v(RHS) \geq 0$  — contradiction.

Hence we can consider this equation as one in the polynomial ring  $L[y_t | t \in T]$ . Moreover, since these elements are algebraically independent, each one appearing in the left hand side must appear in the right hand side and vice versa, i.e.  $T = X_0 \cup Y_0 = \bigcup E_0$ .

By examining the free factor,  $\varepsilon_3^k = 1$ .

By substituting  $y_r$  and  $y_w$  with 0 for some  $r, w$ , we can show that  $E_0 = \{\{r, w\}\}$  (so  $k = i_{\{r,w\}}$ ) and that there are no mixed monomials in the left hand side, i.e. we get an equation of the form

$$\left(\varepsilon_1 (y_r)^{p_0^{i-r} m_r} + \varepsilon_2 (y_w)^{p_0^{i-w} m_w} + 1\right)^{p^k} = \left((y_r)^{p_0^i} + (y_w)^{p_0^i} + 1\right)^{m_{\{r,w\}}}.$$

Suppose  $i = i_r$  and  $i \neq 0$ , then  $p_0 \nmid m_r$ , by examining the degree of  $y_r$ , we get a contradiction, so  $i = 0$  and by choice of  $i$ ,  $i_w = 0$  as well. In the same way we can deduce that  $k = 0$ . From this it follows that  $\varepsilon_1 = \varepsilon_2 = 1$  and  $m_r = m_w$ . So we have

$$f_1 + f_2 + 1 = (x_r^0)^{m_r} + (x_w^0)^{m_w} + 1 = (x_r^0 + x_w^0 + 1)^{m_{\{r,w\}}} = f.$$

So  $\{r, w\}$  is an edge of color  $l$ ,  $m := m_{\{r,w\}} = m_w = m_r$ , and  $m = 1$  or a power of  $r$  (the characteristic).

So finally we have that  $\varphi(x_{t_0}^0)$  is a power of  $m$  which is a power of  $r$ . This implies that  $x_{t_0}^0$  itself has an  $m$ -root. But if  $m > 1$ , this is a contradiction, because  $x_{t_0}^0$  has no  $r$ -roots by Lemma 6.21.

This concludes the proof of the proposition.  $\square$

**COROLLARY 6.23.** *The map  $\sigma : \text{Aut}(\Gamma) \rightarrow \text{Aut}(K_\Gamma)$  is a bijection.*

**PROOF.** Recall that all that is left is to show that  $\sigma$  is onto (by Proposition 6.19).

Let  $\varphi \in \text{Aut}(K_\Gamma)$ . Let  $t \in X$  and suppose  $\{t, t_0\} \in E$ . By Proposition 6.22,  $\varphi(x_t^0) = x_{t'}^0$  for the some  $t' \in X$ . Since the graph  $\Gamma$  is connected, we can define  $\varepsilon \in \text{Aut}(\Gamma)$  by  $\varepsilon(t) = t'$  (note that  $t'$  does not depend on the choice of  $t_0$ ). Proposition 6.22 implies that  $\varepsilon$  is indeed an automorphism.

Since there are no  $p_i$ -roots of unity in  $L$  for all the primes we chose, it follows then that  $\varphi(x_t^i) = x_{\varepsilon(t)}^i$  and that  $\varphi(x_{\varepsilon(t)}^i) = x_{\varepsilon^i(t)}^i$ , and hence  $\varphi = \sigma(\varepsilon)$ .  $\square$

We still have to prove that  $|K_\Gamma| \leq |X^{(<\omega)}|$ .

LEMMA 6.24. *If  $X_i \subseteq X$  ( $i = 1, 2$ ) are two subsets of the vertices set then  $K_{X_1} \cap K_{X_2} = K_{X_1 \cap X_2}$ .*

PROOF. We may assume that  $X_1, X_2$  are finite. Assume  $x \in K_{X_1} \cap K_{X_2}$  and that  $|X_1|$  is minimal with respect to  $x \in K_{X_1}$ . If  $X_1 \subseteq X_2$  then we are done. If not, let  $t \in X_1 \setminus X_2$  be some vertex, and let  $X' = X_1 \setminus \{t\}$ . So  $x \notin K_{X'}$ , and  $x$  is transcendental over  $K_{X'}$  while  $x_t^0$  is algebraic over  $K_{X'}(x)$ . Let  $X'_2 = X' \cup X_2$ ,  $X_3 = X'_2 \cup \{t\}$ . We have  $x \in K_{X_2} \subseteq K_{X'_2}$ , and  $x_t^0 \in K_{X_3}$  is transcendental over  $K_{X'_2}$ . This is a contradiction, because  $x_t^0$  is algebraic over  $K_{X'}(x) \subseteq K_{X'_2}$ . Hence there is no such  $t$  i.e.  $X_1 \subseteq X_2$ .  $\square$

And now it is easy to define an injective map  $\Psi : K_\Gamma \rightarrow X^{(<\omega)}$ . Define by induction on  $n$  injective functions  $\Psi_Y : K_Y \rightarrow X^{(<\omega)}$  for  $|Y| \leq n$  such that  $Y_1 \subseteq Y_2$  implies  $\Psi_{Y_1} \subseteq \Psi_{Y_2}$ . This is enough, since by the lemma above,  $\bigcup \{\Psi_Y \mid Y \subseteq X, |Y| < \omega\}$  is an injection from  $K_\Gamma$  to  $X^{(<\omega)}$ .

For the construction of  $\Psi_Y : K_Y \rightarrow X^{(<\omega)}$ , the idea is that given  $x \in K_Y$  such that  $x \notin K_{Y'}$  for any  $Y' \subsetneq Y$  we can code  $x$  using the set  $Y$  and the set of codes that Lemma 6.8, clause (6) gives us for any choice of  $Y' \subsetneq Y$  of size  $|Y| - 1$ .

This (and Lemma 6.8, clause (6)) was the reason we chose  $X^{(<\omega)}$  and not  $X^{<\omega}$ : in order to code  $x \in K_\Gamma$ , we need first to code the minimal set  $Y$  such that  $x \in K_Y$ , and then  $x$  can be coded in  $|Y|$  different ways, depending on the choice of  $|Y'|$  as above. However, there is no well ordering of  $Y$ , so we have no way of ordering these codes. For instance, the code of  $x_t^0 + x_s^0$  for  $s, t \in X$ , should be  $\{\langle s \rangle, \langle t \rangle, \dots\}$ .

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