

# THE SPECTRUM PROBLEM III: UNIVERSAL THEORIES<sup>†</sup>

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*Dedicated to Professor Abraham Robinson*

## ABSTRACT

We solve the classification problem and essentially the spectrum problem for universal theories (see [6] for discussion of the meaning of this). We first solve it for  $T$  such that if  $M_1, M_2$  elementarily extend  $M_0$  and are independent over it, then over  $M_0 \cup M_1$  there is a prime model. This generalizes [2]. This was subsequently used and generalized for countable first order theories. (This will appear in [5].) But note that there the theory is countable and in the case of structure the model is prime over a non-forking tree of models; here the model is generated by the union (and the  $T$  not necessarily countable). The universal-ity is used in

**THEOREM.** *If  $T$  is stable and complete then either (A) for every  $M_l < M$  ( $l = 0, 1, 2$ ) models of  $T$ , if  $M_0 \subseteq M_1 M_2$ ,  $\{M_1, M_2\}$  is independent over  $M_0$  (i.e.  $\text{tp}(M_1, M_2)$  is finitely satisfiable in  $M_0$ ), then the submodel of  $M$  which  $M_1 \cup M_2$  generates is an elementary submodel of  $M$ , or (B) there is an unstable theory extending the universal part of  $T$  (we can replace universal by  $\Sigma_2$  and slightly more).*

**CONCLUSION.** For any universal  $T$ : *Either (a) for every model  $M$  of  $T$  there is a tree  $I$  with  $\leq \omega$  levels and submodels  $N_\eta$  ( $\eta \in I$ ) of power  $\leq 2^{|T|}$  (by [5], just  $\leq |T|$ ) such that (i)  $M$  is generated by  $\bigcup_{\eta \in I} N_\eta$ , (ii)  $\eta < \nu \Rightarrow N_\eta \subseteq N_\nu$ , (iii) if  $\nu$  is an immediate successor of  $\eta$  then  $\text{tp}(N_\eta, \bigcup\{N_\rho : \rho \in I, \nu \not\leq \rho\})$  is finitely satisfiable in  $N_\eta$  (note that asking this just for quantifier-free formulas is enough). Or (b) for every cardinal  $\lambda > |T|$ , there are  $2^\lambda$  non-isomorphic models for power  $\lambda$ .*

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## TABLE OF CONTENTS

<b>§0. Canonization .....</b>	<b>230</b>
We quote some results from the new edition of the author's book which may be used here.	
<b>§1. On Strong Elementary Submodels .....</b>	<b>231</b>
We present here some exercises on Ch. V (of [1]) in order to deal with models rather than with $\aleph_*$ -saturated models. This is parallel to the role of [3, §1].	
<b>§2. A Prime Atomic Model Over Stable Amalgamation is Enough .....</b>	<b>238</b>
We assume that if $\{M_1, M_2\}$ are independent on $M_0$ , $M_0 < M_i$ then on $M_1 \cup M_2$ there is a prime atomic model (this holds for $T$ totally transcendental, but is much weaker). Then we redo the decomposition lemmas (of [2, §3]) using the first section. The price is that the "building blocks" $N_\alpha$ are less uniquely determined.	
<b>§3. Universal Theories .....</b>	<b>241</b>
This is the heart of the matter. We prove that if every completion of the universal part of $T$ is stable, then the condition from §2 holds. Moreover, the prime model is just the closure of $M_1 \cup M_2$ . We also get the result for slightly more general theories.	
<b>§4. Examples .....</b>	<b>249</b>
We give examples for the possible spectrum.	
<b>§5. Counterexample .....</b>	<b>254</b>
We show that in the main result of §3, we cannot bound the complexity of the formula defining the order.	
<b>References .....</b>	<b>256</b>

*Convention.* We use some convention and definitions freely;  $\mathfrak{C}$  is the "monster" quite saturated model (of  $T$ ).  $\text{Cb}(p)$  is the canonical base (see III §7).

When referring to [1], we do not write [1] usually (no confusion arises as we have to use a chapter number (latin numeral)).

In Sections 0–3,  $\text{Th}(\mathfrak{C})$  is assumed to be stable.  $<$  is elementary submodel.

## §0. Canonization

This section can be avoided, if you avoid 1.5A (but not 1.5), 1.8 and 2.3(1) which are not used later.

Here we quote some facts from the new edition of the author's book [5] using the definition from there (see Definition V 4.5). So  $\mathfrak{C}$  is a quite saturated model of a complete first-order stable  $T$ .

0.1. CLAIM. ( $\mathfrak{C}^{\text{eq}}$ ) Suppose  $r$  is regular and  $\text{stp}(\bar{a}, A)$  is not orthogonal to  $r$ . Then:

(1) There is  $E \in FE^m(\text{acl } A)$  ( $m = l(\bar{a})$ ) such that  $\text{stp}(\bar{c}/E, \bar{b})$  is  $\text{cl}^2(r)$ -simple but not orthogonal to  $r$ .

- (2) Moreover  $\text{stp}(\bar{c}/E, A)$  contains a formula  $\vartheta(y, \bar{b})$  that is  $\text{cl}^2(r)$ -simple.  
 (3) Also, for every  $\bar{b}'$  realizing  $\text{tp}(\bar{b}, \emptyset)$ ,  $\vartheta(y, \bar{b}')$  is  $\text{cl}^3(r)$ -simple.

0.1A. FACT. There is  $\varphi(x, \bar{c})$  which is  $\text{cl}^3(r)$ -regular but has completions not orthogonal to  $r$ .

Let  $\varphi(x, \bar{c})$  be any formula which has completions not orthogonal to  $r$ , but has no extension of smaller rank  $(R^1(-, L, \infty))$  with this property.

Clearly such  $\varphi(x, \bar{c})$  exists. Now any such  $\varphi(x, \bar{c})$  is  $\text{cl}^2(r)$ -regular — this follows by

$\oplus$  for every complete stationary  $p$ ,  $\varphi(x, \bar{c}) \in p$ : if  $R(p, L, \infty) < R(\varphi(x, \bar{c}), L, \infty)$  then  $p$  is orthogonal to  $r$ ; hence if  $p$  forks over  $\bar{c}$  then  $p$  is orthogonal to  $r$ .

0.1B. FACT. Every complete stationary  $\mathcal{P}$ -regular  $p$  not orthogonal to  $r \in \mathcal{P}$  is regular.

0.1C. FACT. There is  $\varphi(x, \bar{c})$  as in Fact 0.1A, such that for every  $\bar{c}'$  realizing  $\text{tp}(\bar{c}, \emptyset)$ ,  $\varphi(x, \bar{c}')$  is  $\text{cl}^2(r)$ -regular.

0.2. CLAIM. ( $\mathfrak{C}^{\text{eq}}$ ) If  $\text{tp}(\bar{a}, A)$  is not orthogonal to some trivial regular type  $r$  then for some  $b \in \text{act}(A \cup \bar{a}) - \text{acl}(A)$ ,  $\text{tp}(b, A)$  is  $\text{cl}^3(\{r\})$ -simple and  $w_r(b, A) = 1$ . If  $\text{tp}(\bar{a}, A)$  is regular we can replace simple by regular, and if  $T$  is superstable then some  $\varphi \in \text{tp}(b, A)$  is  $\text{cl}^3(\{r\})$ -simple [regular].

## §1. On Strong Elementary Submodels

HYPOTHESIS.  $T$  is superstable.

1.1. DEFINITION. (1) We say that  $M \subseteq N$  if  $M \subseteq N$  and for every  $\bar{a} \in M$ ,  $\bar{b} \in N$  there is  $\bar{b}' \in M$  realizing  $\text{tp}(\bar{b}, \bar{a})$ . We define  $M \subseteq_s A$ ,  $B \subseteq_s A$  similarly.

(2) We say that  $M \subseteq_a N$  if  $M \subseteq N$ , and for every  $\bar{a} \in M$ ,  $\bar{b} \in N$  there is  $\bar{b}' \in M$  realizing  $\text{stp}(\bar{b}, \bar{a})$ . We define  $M \subseteq_a A$ ,  $B \subseteq_a A$  similarly.

1.2. CLAIM. (1) If  $A \subseteq B \subseteq C$ ,  $A \subseteq_s C$  then  $A \subseteq_s B$ .

(2) Let  $A \subseteq_s B_l$  for  $l < n$ ,  $\bar{a} \in A$ ,  $\bar{b}_l \in B_l$ ,  $\models \varphi[\bar{b}_0, \dots, \bar{b}_{n-1}, \bar{a}]$  and  $\{B_l : l < n\}$  is independent over  $A$ . Then there are  $\bar{b}'_0, \dots, \bar{b}'_{n-1} \in A$ , such that:  $\models \varphi[\bar{b}'_0, \dots, \bar{b}'_{n-1}, \bar{a}]$  and  $\bar{b}'_l$  realizes  $\text{tp}(\bar{b}_l, \bar{a})$  (and in fact it realizes  $\text{tp}(\bar{b}_l, \bar{a} \cup \bigcup_{m < l} \bar{b}'_m)$ ) and  $\{\bar{b}'_0, \dots, \bar{b}'_{n-1}\}$  is independent over some  $\bar{a}'$ ,  $\bar{a} \subseteq \bar{a}' \subseteq A$ .

(3) If  $N \subseteq N_1 \subseteq M$ ,  $N \subseteq_a A$ ,  $M$  is  $\mathbf{F}_{\kappa_0}^1$ -atomic over  $N_1 \cup A$ ,  $\{A, N_1\}$  is independent over  $N$ . Then  $N_1 \subseteq_s M$  (in fact  $N_1 \subseteq_a M$  holds,  $M$  a set suffices).

(4) If  $A \subseteq B \subseteq C$ ,  $A \subseteq_a C$  then  $A \subseteq_a B$ .

(5) If  $A \subseteq_a B_i$  ( $i < \alpha$ ) and  $\{B_i : i < \alpha\}$  is independent over  $A$ , then  $A \subseteq_a \bigcup_{i < \alpha} B_i$ .

(6) If  $A \subseteq_a B$  then  $A \subseteq_s B$ .

PROOF. (1), (4), (5), (6) are easy.

(2) We can find a finite  $\bar{a}'$ ,  $\bar{a} \subseteq \bar{a}' \subseteq A$ , such that  $\text{tp}(\bar{b}_0 \wedge \dots \wedge \bar{b}_{n-1}, A)$  does not fork over  $\bar{a}'$ . Now we define by induction on  $i < \omega$ ,  $\langle \bar{b}_i^l : l < n \rangle$ , such that  $\bar{b}_i^l \in A$ . For a given  $i$  we define  $\bar{b}_i^l$  by induction on  $l$ :  $\bar{b}_i^1$  realizes  $\text{tp}(\bar{b}_i, \bar{a}' \cup \bigcup_{j < i} \bigcup_{m < n} \bar{b}_m^j \cup \bigcup_{m < i} \bar{b}_m^i)$  and  $\bar{b}_i^1 \in A$ . We can easily prove by induction on  $k < \omega$  that  $\{\bar{b}_i^l : ni + l < k\}$  is independent over  $\bar{a}'$ . As in the proof of [1; Ch. II, 2.17, p. 38, Ch. III, 2.13, pp. 98] we finish.

(3) Let  $\bar{a} \in N_1$ ,  $\bar{b} \in M$  and we should find  $\bar{b}' \in N_1$  realizing  $\text{stp}(\bar{b}, \bar{a})$ . As  $M$  is  $\mathbf{F}_{\aleph_0}^c$ -atomic over  $N_1 \cup A$ , there are  $\bar{b}_1 \in N_1$ ,  $\bar{b}_2 \in A$  such that  $\models \psi[\bar{b}, \bar{b}_1, \bar{b}_2]$  and  $\psi(\bar{x}, \bar{b}_1 \bar{b}_2) \vdash \text{tp}(\bar{b}, N_1 \cup A)$ ; w.l.o.g.  $\bar{a} \subseteq b_1$ . For some  $\bar{b}_0 \in N$ ,  $\text{tp}(\bar{b}_i, N)$  does not fork over  $\bar{b}_0$  for  $i = 1, 2$  and remember  $\{\bar{b}_1, \bar{b}_2\}$  is independent over  $N$ .

Now choose  $\bar{b}'_2 \in N$  which realizes  $\text{stp}(\bar{b}_2, \bar{b}_0)$  (possible as  $N \subseteq_a A$ ). As  $\text{tp}(\bar{b}_1, A)$  does not fork over  $\bar{b}_0$ ,  $\bar{b}'_2$  realizes  $\text{stp}(\bar{b}_2, \bar{b}_0 \cup \bar{b}_1)$ . So  $\bar{b}_0 \wedge \bar{b}_1 \wedge \bar{b}_2$ ,  $\bar{b}_0 \wedge \bar{b}_1 \wedge \bar{b}'_2$  realizes the same type, hence  $\models (\exists \bar{x}) \psi(\bar{x}, \bar{b}_1, \bar{b}'_2)$ , and letting  $\bar{b}'$  realize  $\psi(\bar{x}, \bar{b}_1, \bar{b}'_2)$ ,  $\bar{b}_0 \wedge \bar{b}_1 \wedge \bar{b}_2 \wedge \bar{b}$ ,  $\bar{b}_0 \wedge \bar{b}_1 \wedge \bar{b}'_2 \wedge \bar{b}'$  realizes the same type. But we can choose  $\bar{b}' \in N_1$ . So there is  $\bar{b}' \in N_1$  realizing  $\text{tp}(\bar{b}, \bar{a})$  (as  $\bar{a} \subseteq \bar{b}_1$ ). In fact  $\bar{b}'$  realizes  $\text{stp}(\bar{b}, \bar{a})$ . For every  $E \in FE(\bar{b}_1)$ ,  $E$  is a formula over  $N_1$  (as it is almost over  $\bar{b}_1$ ) hence  $\text{tp}(\bar{b} \wedge \bar{b}_2, N_1) \vdash E(\bar{x}, \bar{y}; \bar{b}, \bar{b}_2)$ . But for every  $\Theta(\bar{x}, \bar{y}) \in \text{tp}(\bar{b} \wedge \bar{b}_2, N_1)$ ,  $\text{tp}(\bar{b}, N_1 \cup \bar{b}_2) \vdash \Theta(\bar{x}, \bar{b}_2)$ , hence  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2) \vdash \Theta(\bar{x}, \bar{b}_2)$ . We can conclude that  $\psi(\bar{x}, \bar{b}_1, \bar{b}'_2) \vdash E(\bar{x}, \bar{b}'_2; \bar{b}, \bar{b}_2)$  [as  $\Theta(\bar{x}, \bar{y}) \stackrel{\text{def}}{=} E(\bar{x}, \bar{y}; \bar{b}, \bar{b}_2)$  is almost over  $\bar{b}_1$  and  $\bar{b}'_2$  realizes  $\text{stp}(\bar{b}_2, \bar{b}_1 \wedge \bar{b}_0)$ , clearly  $\psi(\bar{x}, \bar{b}_1, \bar{b}'_2) \vdash \Theta(\bar{x}, \bar{b}'_2)$  hence  $\psi(\bar{x}, \bar{b}_1, \bar{b}'_2) \vdash E(\bar{x}, \bar{b}'_2; \bar{b}, \bar{b}_2)$ ], hence  $\models E(\bar{b}', \bar{b}'_2; \bar{b}, \bar{b}_2)$ . As this holds for every  $E \in FE(\bar{b}_1)$  clearly  $\bar{b}' \wedge \bar{b}'_2$  realizes  $\text{stp}(\bar{b} \wedge \bar{b}_2, \bar{b}_1)$ , hence  $\bar{b}'$  realizes  $\text{stp}(\bar{b}, \bar{a})$ .

1.3. LEMMA. Suppose  $N \subseteq A$ ,  $N \subseteq_a M$ ,  $\{M, A\}$  is independent over  $N$ :

(1) If  $\text{tp}(\bar{a}, A)$  is orthogonal to  $N$  or is  $\mathbf{F}_{\aleph_0}^c$ -isolated (for  $\mathbf{F}_{\aleph_0}^c$ -isolation see Definition IV 2.7), then  $\text{tp}(\bar{a}, A) \vdash \text{tp}(\bar{a}, A \cup M)$ .

(2) If for each  $i < \alpha$ ,  $\text{tp}(\bar{a}_i, A \cup \bigcup_{j < i} \bar{a}_j)$  is orthogonal to  $N$  or  $\mathbf{F}_{\aleph_0}^c$ -isolated (for  $\mathbf{F}_{\aleph_0}^c$ -isolation see Definition IV, 2.7) then

$$\text{tp}_*(\{\bar{a}_i : i < \alpha\}, A) \vdash \text{tp}_*(\{\bar{a}_i : i < \alpha\}, M \cup A).$$

(3) In (1), (2) we can replace  $M$  by any  $B$  (if  $N \subseteq_a B$ ).

PROOF. (1) Trivial (for the second case note  $\text{dcl}(M \cup A) \cap \text{acl}(A) = \text{dcl}(A)$ ).

(2) Follows from (1) by induction on  $\alpha$ .

(3) The same proof.

1.4. CLAIM. (1) (in  $\mathcal{C}^{eq}$ ) Let  $x \in \{t, a\}$ . If  $p$  is an  $m$ -type over  $A$ ,  $r$  is a regular type,  $\text{stp}(\bar{a}, A)$  is not orthogonal to  $r$ ,  $r$  extends  $p$ ,  $p$  is an  $\mathbf{F}_\mu^x$ -type,  $A = |N|$ ,  $N$  is  $\mathbf{F}_\mu^x$ -compact, then  $\text{stp}_*(B, A)$  is not orthogonal to  $r$  where  $B = \text{Cb}(\text{tp}(\bar{a}, A \cup p(\mathcal{C}^{eq})))$ .

(2) Suppose  $\bar{b} \in M$ ,  $\bar{a} \notin M$ ,  $\text{tp}(\bar{a}, M)$  is not orthogonal to some type  $q$  to which  $\psi(\bar{x}, \bar{b})$  belongs. Then for every model  $N$  including  $M \cup \bar{a}$ ,  $\psi(N, \bar{b}) \neq \psi(M, \bar{b})$ .

REMARK. The most interesting cases of (1) are  $\mathbf{F}_\mu^x = \mathbf{F}_{\kappa_0}^t$ ,  $p$  finite (so  $N$  is just a model) and  $\kappa = \kappa_r(T)$ ,  $|\text{Dom } p| < \kappa$  (even for any stable  $T$ ).

PROOF. (1) We shall prove later:

1.4A. FACT. For every  $\bar{d}$

$\text{tp}(\bar{a}, \bar{d} \cup A \cup p(\mathcal{C}^{eq}))$  does not fork over  $A \cup B$   
whenever  $\text{tp}(\bar{a}, A \cup \bar{d})$  does not fork over  $A$ .

Let  $\lambda > |T| + |A| + \|M\|$  be regular, let  $N_0$  be an  $\mathbf{F}_\lambda^a$ -saturated model,  $A \subseteq N_0$ ; and w.l.o.g.  $r$  is a type over  $N_0$  and  $\text{tp}_*(N_0, A \cup \bar{a})$  does not fork over  $A$ .

By 1.4A,  $\text{tp}(\bar{a}, N_0 \cup p(\mathcal{C}^{eq}))$  does not fork over  $N \cup B$  hence over  $N_0 \cup B$ . Let  $p \subseteq r_0 \in S^m(N_0)$ ,  $r_0$  regular not orthogonal to  $r$  ( $r_0$  exists, of course). Let  $N_1$  be  $\mathbf{F}_\lambda^a$ -prime over  $N_0 \cup \bar{a}$ , hence it is  $\mathbf{F}_\lambda^a$ -atomic over  $N_0 \cup \bar{a}$  hence over  $N_0 \cup \bar{a} \cup B$  (note that  $B \subseteq N_1$ , as  $B \subseteq \text{acl}(N \cup \bar{a})$ ). As  $\text{tp}_*(\bar{a}, N_0 \cup p(\mathcal{C}^{eq}))$  does not fork over  $N_0 \cup B$  clearly  $\text{tp}_*(\bar{a}, N_0 \cup p(N_1))$  does not fork over  $N_0 \cup B$  hence (by III, 0.1)  $\text{tp}_*(p(N_1), N_0 \cup B \cup \bar{a})$  does not fork over  $N_0 \cup B$ . By IV, 4.3 it is easy to see that  $p(N_1)$  is a  $\mathbf{F}_\lambda^a$ -atomic over  $N_0 \cup B$  ( $B \subseteq N_1$  as noted above). So if 1.4(1)'s conclusion fails, as for every  $\bar{b} \in B$ , as  $\text{tp}(\bar{b}, N_0)$  is a stationarization of  $\text{stp}(\bar{b}, A)$ , clearly it is orthogonal to  $r$ .

Note  $\text{tp}(\bar{b}, N_0)$  is orthogonal to  $r_0$ .

So  $\text{stp}_*(B, N_0)$  is orthogonal to  $r_0$  and  $p(N_1)$  is  $\mathbf{F}_\lambda^a$ -atomic over  $N_0 \cup B$ .

Hence if  $N_2$  is  $\mathbf{F}_\lambda^a$ -prime over  $N_0 \cup B$ ,  $r_0$  is not realized in  $N_2$  and for every  $\bar{c} \in p(N_1)$ ,  $\text{tp}(\bar{c}, N_0 \cup B)$  is  $\mathbf{F}_\lambda^a$ -isolated hence is realized in  $N_2$ . We can conclude that (as  $p$  is over  $N_0$ )  $N_1$  does not realize  $r_0$ . But as  $\text{tp}(\bar{a}, N_0)$  is not orthogonal to  $r_0$ ,  $r_0$  is realized in  $N_1$ , and  $p \subseteq r_0$ , contradiction.

PROOF OF FACT 1.4A. Why does it hold? As we have assumed  $\text{tp}(\bar{a}, A \cup \bar{d})$  does not fork over  $A$ , also  $\text{tp}(\bar{d}, A \cup \bar{a})$  does not fork over  $A$ , hence [as

$B \subseteq \text{acl}(A \cup \bar{a})$ ]  $\text{tp}(\bar{d}, A \cup \bar{a} \cup B)$  does not fork over  $A$ , hence [as  $A$  is a model] is finitely satisfiable in  $A$ . Suppose the conclusion of 1.4A fails, i.e., for some  $\bar{e} \in A \cup \bar{d} \cup p(\mathbb{U}^{\text{eq}})$ , and  $\varphi, \models \varphi[\bar{a}, \bar{e}]$ , but  $\varphi(x, \bar{e})$  forks over  $A \cup B$ . W.l.o.g. for some finite  $\Delta$ ,  $R(\varphi(\bar{x}, \bar{e}), \Delta, \aleph_0) < R(\text{tp}(\bar{a}, A), \Delta, \aleph_0)$ .

We know that for some large enough finite  $\kappa$ ,  $R[\varphi(\bar{x}, \bar{e}), \Delta, \kappa] = R[\varphi(\bar{x}, \bar{e}), \Delta, \aleph_0]$ ,  $R[\text{tp}(\bar{a}, A \cup B), \Delta, \kappa] = R[\text{tp}(\bar{a}, A \cup B), \Delta, \aleph_0]$ . W.l.o.g. for every  $\bar{e}'$ ,  $R[\varphi(\bar{x}, \bar{e}'), \Delta, \kappa] < R[\varphi(\text{tp}(\bar{a}, A \cup B), \Delta, \kappa)]$ . W.l.o.g.  $\bar{e} = \bar{d} \wedge \bar{a}' \wedge \bar{e}_1 \wedge \dots \wedge \bar{e}_n$  where  $\bar{a}' \subseteq A$ ,  $\bar{e}_i \in p(\mathbb{U}^{\text{eq}})$ ; so  $\models \varphi(\bar{a}, \bar{d}, \bar{a}', \bar{e}_1, \dots, \bar{e}_n)$  and  $p$  w.l.o.g. is a singleton or closed under conjunction. So  $\bar{d}$  realizes the type

$$q = \left\{ (\exists \bar{z}_1, \dots, \bar{z}_m) \left[ \varphi(\bar{a}, \bar{x}, \bar{a}', \bar{a}_1, \dots, \bar{z}_m) \wedge \bigwedge_{i=1}^m \Theta(\bar{z}_i) \right] : \Theta(\bar{z}) \in p \right\}.$$

As  $q$  does not fork over  $A$ ,  $q$  an  $F_\mu^*$ -type, some  $\bar{d}' \in A$  realizes  $q$ . So for some  $\bar{e}'_i \in p(\mathbb{U}^{\text{eq}}) \models \varphi[\bar{a}, \bar{d}', \bar{a}', \bar{e}'_1, \dots]$ . But this is a contradiction as

$$\begin{aligned} R[\text{tp}(\bar{a}, A \cup p(\mathbb{U}^{\text{eq}}), \Delta, \aleph_0] &= R[\text{tp}(\bar{a}, A \cup B), \Delta, \aleph_0] \\ &= R[\text{tp}(\bar{a}, A \cup B), \Delta, \kappa] > R[\varphi(x, \bar{d}', \bar{a}', \bar{e}'_1, \dots), \Delta, \kappa]. \end{aligned}$$

PROOF OF 1.4(2). By (1) with  $(p = \{\varphi(\bar{x}, \bar{b})\}, A = M)$ . We know that  $\text{Cb}(\text{tp}(\bar{a}, M \cup \varphi(\mathbb{U}, \bar{b})))$  is not contained in  $M$ , but it is contained in  $\text{acl}(M \cup \varphi(N, \bar{b}))$ . Hence  $\varphi(N, \bar{b}) \not\subseteq M$ .

1.5. LEMMA. Assume  $N \subseteq_a M'$ ,  $N \subseteq M' \subseteq M$  and  $m < \omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M'$ ,  $l(\bar{c}) = m$  and  $R[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal (under this condition). If  $\text{tp}(\bar{c}, M')$  is not orthogonal to  $N$  then there is  $\bar{c}' \in M$ ,  $l(\bar{c}') = l(\bar{c})$ ,  $\bar{c}' \notin M'$ ,  $\text{tp}(\bar{c}', M')$  does not fork over  $M$  and  $R[\text{tp}(\bar{c}', M'), L, \infty] = R[\text{tp}(\bar{c}, M), L, \infty]$ .

PROOF. Let  $\psi(\bar{x}, \bar{b}) \in \text{tp}(\bar{c}, M')$ ,  $R^m[\psi(\bar{x}, \bar{b}), L, \infty] = R^m[\text{tp}(\bar{c}, M'), L, \infty]$  hence  $\text{tp}(\bar{c}, M')$  does not fork over  $\bar{b}$ . We work in  $\mathbb{U}^{\text{eq}}$ . We can choose  $\bar{a} \in N$ , such that  $\text{tp}(\bar{b} \wedge \bar{c}, N)$  does not fork over  $\bar{a}$  and  $\text{stp}(\bar{c}, \bar{b})$  is not orthogonal to  $\bar{a}$  (as  $\text{tp}(\bar{c}, M')$ ,  $\text{stp}(\bar{c}, \bar{b})$  are parallel,  $\text{stp}(\bar{c}, \bar{b})$  is not orthogonal to some  $q \in \bigcup_n S^n(N)$  and w.l.o.g.  $q$  does not fork over  $\bar{a}$ ). Now as  $N \subseteq_a M'$  there is a sequence  $\bar{b}' \in N$  realizing  $\text{stp}(\bar{b}, \bar{a})$  and for some  $\bar{c}'$ ,  $\bar{b}' \wedge \bar{c}'$  realizes  $\text{stp}(\bar{b} \wedge \bar{c}, \bar{a})$ . By V, 3.5  $\text{stp}(\bar{c}, \bar{b} \cup \bar{a})$ ,  $\text{stp}(\bar{c}', \bar{b}' \cup \bar{a})$  are not orthogonal (note that  $\text{stp}(\bar{c}, \bar{b} \cup \bar{a})$ ,  $\text{stp}(\bar{c}, \bar{b})$  and  $\text{stp}(\bar{c}, M')$  are parallel). So  $\text{tp}(\bar{c}, M')$  is not orthogonal to some type to which  $\psi(\bar{x}, \bar{b}')$  belongs, hence by 1.4(2) there is  $\bar{c}'' \in M$ ,  $\bar{c}'' \notin M'$ ,  $\models \psi[\bar{c}'', \bar{b}']$ . So

$$R[\text{tp}(\bar{c}'', M'), L, \infty] \leq R[\psi(\bar{x}, \bar{b}'), L, \infty] = R[\psi(\bar{x}, \bar{b}), L, \infty] = R[\text{tp}(\bar{c}, M'), L, \infty].$$

By the hypothesis that  $R[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal equality holds, hence  $\text{tp}(\bar{c}'', M')$  does not fork over  $N$ . So  $\bar{c}''$  is as required.

1.5A. LEMMA. Suppose  $N \subseteq M$  and  $N \subseteq M' \subseteq M$  and  $m < \omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M'$ ,  $l(\bar{c}) = m$ ,  $\text{tp}(\bar{c}, M')$  is not orthogonal to  $N$  and  $R^m[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal (under the previous constraints).

Then there is  $\bar{c}' \in M$ ,  $\bar{c}' \notin M'$ ,  $l(\bar{c}) = l(\bar{c}')$ ,  $\text{tp}(\bar{c}', M')$  does not fork over  $N$  and

$$R^m[\text{tp}(\bar{c}', M'), L, \infty] = R^m[\text{tp}(\bar{c}, M'), L, \infty].$$

REMARK. We can waive this lemma if in the decomposition theorems we omit 2.3(1).

PROOF. Let  $\varphi(\bar{x}, \bar{b}) \in \text{tp}(\bar{c}, M')$ ,  $R^m[\varphi(\bar{x}, \bar{b}), L, \infty] = R^m[\text{tp}(\bar{c}, M'), L, \infty]$ . We work in  $\mathcal{U}^{\text{eq}}$ . Let  $r^*$  be a regular type not orthogonal to  $N$  and not orthogonal to  $\text{tp}(\bar{c}, M')$ , with  $R^1[r^*, L, \infty]$  minimal. By 0.1(1) (maybe replacing  $\bar{b}$  by  $\bar{b}' \subseteq \text{acl } \bar{b}$ ) there is a formula  $E = E(\bar{x}, \bar{y}, \bar{b})$  where  $\bar{b} \in M'$ ,  $\text{tp}(\bar{c}, M')$  does not fork over  $\bar{b}$ , such that  $\text{stp}(\bar{c}/E, \bar{b})$  is  $\text{cl}^3(r^*)$ -simple not orthogonal to  $r$ . Moreover (see 0.1(2), (3)):

$\oplus$   $\text{stp}(\bar{c}/E, \bar{b})$  contains a formula  $\Theta(y, \bar{b})$  which is  $\text{cl}^3(r^*)$ -simple.

Moreover,  $\text{tp}(\bar{b}, \emptyset) = \text{tp}(\bar{b}', \emptyset)$  implies  $\Theta(x, \bar{b}')$  is  $\text{cl}^3(r^*)$ -simple.

Let  $\varphi_0(y, \bar{b}) = (\exists \bar{x})[\varphi(\bar{x}, \bar{b}) \wedge \bar{x}/E = y \wedge \Theta(y, \bar{b})]$ . We shall prove

(\*) there are  $\bar{b}' \in N$  realizing  $\text{tp}(\bar{b}, \emptyset)$ , and  $c' \in M - M'$ ,  $\models \varphi_0[c', \bar{b}']$ ,  $\text{tp}(c', M')$  not orthogonal to  $N$  (in fact, to  $r^*$ ).

If (\*) holds, there is  $\bar{c}'' \in M$  such that  $\varphi[\bar{c}'', \bar{b}'] \wedge \bar{c}''/E = c' \wedge \Theta(c', \bar{b})$ . Now  $\text{tp}(\bar{c}'', M')$  is not orthogonal to  $N$  as  $c' \in \text{acl}(N \cup \{\bar{c}''\})$ ,  $\text{tp}(c', M')$  is not orthogonal to  $N$ . Now  $R^m[\text{tp}(\bar{c}'', M'), L, \infty] \leq R^m[\varphi(\bar{x}, \bar{b}'), L, \infty] = R^m[\varphi(\bar{x}, \bar{b}), L, \infty]$ . Equality holds by the hypothesis " $R^m[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal". ( $\bar{c}'' \notin M'$  follows from:  $\text{tp}(\bar{c}'', M')$  is not orthogonal to  $N$ .) Hence  $\text{tp}(\bar{c}'', M')$  does not fork over  $\bar{b}' \subseteq N$ , so we get our conclusion.

Now we shall prove (\*). For notational simplicity let  $E$  be the equality, and  $\varphi(\bar{x}, \bar{b}) \vdash \Theta(\bar{x}, \bar{b})$ . Choose  $\bar{a} \in N$  such that  $\text{tp}(\bar{b}, N)$  does not fork over  $\bar{a}$ . We have assumed that  $r^*$  (hence  $\text{tp}(\bar{c}, M')$ ) is not orthogonal to  $N$ , hence it is not orthogonal to some  $r \in S^*(N)$ . Also, let  $\bar{d}_0$  realize  $r$ ,  $\text{tp}(\bar{d}_0, M)$  does not fork over  $N$ . W.l.o.g.  $r$  does not fork over  $\bar{a}$ .

Now there are  $\bar{b}_n, \bar{c}_n$  ( $n < \omega$ ) such that  $\bar{b}_n \wedge \bar{c}_n$  realizes  $\text{stp}(\bar{b} \wedge \bar{c}, N)$ ,  $\{\bar{b}_n \wedge \bar{c}_n : 0 < n < \omega\}$  is independent over  $(M \cup \bar{d}_0, N)$ ,  $\bar{b}_0 \wedge \bar{c}_0 = \bar{b} \wedge \bar{c}$  (so  $\{\bar{b}_n \wedge \bar{c}_n : n < \omega\}$  is independent over  $N$ ). Note that  $\Theta(x, \bar{b}_n)$  is  $\text{cl}^3(r)$ -simple.

For each  $n$  let  $\{\bar{c}_{n,i} : i < \omega\}$  be a family of sequences realizing  $\text{stp}(\bar{c}_n, \bar{b}_n \cup N)$ , independent over  $N \cup \bar{b}_n$ . Let  $\{\bar{d}^i : i < \omega\}$  be a family of sequences realizing

$\text{tp}(\bar{d}_0, N)$  independent over  $(N \cup \bigcup_n \bar{b}_n \cup \bigcup_{n,i} \bar{c}_{n,i}, N)$ . By V 2.7 for some  $k, l$   $\text{tp}(\bar{d}^0 \wedge \dots \wedge \bar{d}^l, N \cup \bar{c}_n)$ ,  $\text{tp}(\bar{c}_{n,0} \wedge \dots \wedge \bar{c}_{n,k}, N \cup \bar{c}_n)$  are not weakly orthogonal. Now w.l.o.g.  $l = 0$ .

Now by the proof of V 4.11 there is  $d_1$ ,  $d_1 \in \text{acl}(N \cup \bar{d}_0) - N$ ,  $\text{tp}(d_1, N)$  not orthogonal to  $r^*$  and some  $n < \omega$ , and there is a formula  $\psi(x, \bar{a})$  such that

(a)  $\models \psi[d_1, \bar{a}]$ ;

(b)  $d_1 \in \text{acl}(\bar{a} \cup \bigcup_{l < n} \bar{b}_l \cup \bigcup_{l < n, i < i(l)} \bar{c}'_{l,i})$  where  $\bar{c}'_{l,i}$  realizes  $\text{stp}(\bar{c}_l, \bar{b}_l)$ ,  $i(l) < \omega$  (for  $l < n$ ).

(c)  $\psi(\mathbb{U}^{\text{eq}}, \bar{a}) \subseteq \text{acl}(\bar{a} \cup \bigcup_{l < n} \bar{b}_l \cup \bigcup_{l < n} \varphi(\mathbb{U}^{\text{eq}}, \bar{b}_l))$ .

By 1.2(2) there are  $\bar{b}'_l \in N$  ( $l < n$ ) realizing  $\text{tp}(\bar{b}_l, \bar{a})$  such that

(d)  $\psi(\mathbb{U}^{\text{eq}}, \bar{a}) \subseteq \text{acl}(\bar{a} \cup \bigcup_{l < n} \bar{b}'_l \cup \bigcup_{l < n} \varphi(\mathbb{U}^{\text{eq}}, \bar{b}'_l))$ .

This is not exactly a first order property, but if it holds then some first order formula witnesses it, by compactness. Note that  $\Theta(\bar{x}, \bar{b}_l)$  is  $\text{cl}^1(r^*)$ -simple (as  $\text{tp}(b'_l, \emptyset) = \text{tp}(\bar{b}_l, \emptyset)$ ). Remember  $\text{tp}(d_1, N)$  is not orthogonal to  $r^*$ ,  $\psi(x, \bar{a}) \in \text{tp}(d_1, N)$ . Easily  $\text{tp}(\bar{c}, M' \cup \psi(\mathbb{U}^{\text{eq}}))$  does not fork over  $\bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}})$  hence  $\text{Cb}(\text{tp}(\bar{c}, M' \cup \psi(\mathbb{U}^{\text{eq}}))) = \text{Cb}(\text{tp}(\bar{c}, \bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}})))$  and the type of this set over  $M' \cup \psi(\mathbb{U}^{\text{eq}})$  does not fork over  $\bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}})$ .

Hence by 1.4(1) in  $\text{Cb}(\text{tp}(\bar{c}, \bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}})))$  there is an element  $d_2 \notin \text{acl}(N \cup \bar{b})$  such that  $\text{tp}(d_2, N \cup \bar{b})$  is not orthogonal to  $r^*$ . As  $\text{tp}(\bar{c}, M')$  does not fork over  $\bar{b}$ , clearly  $\text{tp}(d_2, M')$  does not fork over  $\bar{b} \cup N$ . So as  $d_2 \notin \text{acl}(\bar{b} \cup N)$  also  $d_2 \notin M'$  remembering that

$$\text{Cb}(\text{tp}(\bar{c}, \bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}}, \bar{a}))) \subseteq \text{acl}(\bar{b} \cup N \cup \psi(\mathbb{U}^{\text{eq}}, \bar{a})).$$

Clearly  $d_2 \in M$ .

Now there are  $\bar{c}'_{l,i} \in M$ ,  $d_2 \in \text{acl}(\bar{b} \cup N \cup \{\bar{c}'_{l,i} : l, i\})$ ,  $\models \varphi[\bar{c}'_{l,i} \bar{b}'_l]$  (this is by (d)). Now by  $\oplus$ , if  $\text{tp}(\bar{c}'_{l,i}, M')$  is orthogonal to  $r^*$ , then  $\text{tp}(\bar{c}'_{l,i}, M' \cup \{c_{k,j} : k < l < i \text{ or } k = l, j < i\})$  is also orthogonal to  $r^*$ . If this holds for every  $l, i$  then  $\text{tp}_*(\{\bar{c}_{l,i} : l, i\}, M')$  is orthogonal to  $r^*$ . Hence  $\text{tp}(d_2, M')$  is orthogonal to  $r^*$ , contradiction. Now if  $\text{tp}(\bar{c}'_{l,i}, M')$  is not orthogonal to  $r^*$ ,  $(*)$  holds, and we are finished.

1.6. LEMMA. Suppose  $N \subseteq M$  and  $N \subseteq M' \subseteq M$ , and  $m < \omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M'$ ,  $l(\bar{c}) = m$ ,  $\text{tp}(\bar{c}, M')$  is not orthogonal to  $N$  and  $R[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal.

If  $\text{tp}(\bar{c}, M')$  does not fork over  $N$  then it is regular.

REMARK. We can omit 1.6 if in §2 we waive the regularity, i.e., omit 2.2(b) and 2.4(3).



PROOF. Let  $\bar{b} \in N$ ,  $\text{tp}(\bar{c}, M')$  does not fork over  $\bar{b}$ . We can choose  $\psi$  such that  $\models \psi[\bar{c}, \bar{b}]$ ,  $R[\psi(\bar{x}, \bar{b}), L, \infty] = R[\text{tp}(\bar{c}, M'), L, \infty]$ . Let  $\bar{c}_n$  realize  $\text{stp}(\bar{c}, \bar{b})$  for  $n < \omega$ ,  $\{\bar{c}_n : n < \omega\}$  independent over  $(M, \bar{b})$ . If  $\text{stp}(\bar{c}, \bar{b})$  is not regular then there are  $m$  and  $\bar{c}_j^*$  ( $j < m$ ) and  $n$  such that  $\bar{c}_j^*$  realizes  $\text{stp}(\bar{c}, \bar{b})$ ,  $\text{tp}(\bar{c}_j^*, \bar{b} \cup \{\bar{c}_i : i \leq n\})$  forks over  $\bar{b}$  and  $\text{tp}(\bar{c}, \bar{b} \cup \{\bar{c}_i : i \leq n\} \cup \{\bar{c}_j^* : j < m\})$  forks over  $\bar{b}$ . For each  $j < m$  for some  $n_j \leq n$ ,  $\text{tp}(\bar{c}_{n_j}, \bar{b} \cup \{\bar{c}_i : i < n_j\} \cup \{\bar{c}_j^*\})$  fork over  $\bar{b}$  and remember  $\text{tp}(\bar{c}, \bar{b} \cup \{\bar{c}_i : i \leq n\} \cup \{\bar{c}_j^* : j < m\})$  forks over  $\bar{b}$ . By III 1.2(2), III 2.6(2), II 3.7 for some finite  $\Delta, k$

$$R[\text{tp}(\bar{c}_{n_j}, \bar{b} \cup \{\bar{c}_i : i < n_j\} \cup \{\bar{c}_j^*\}), \Delta, k] < R[\text{tp}(\bar{c}, \bar{b}), \Delta, k]$$

and

$$R[\text{tp}(\bar{c}, \bar{b}), \Delta, k] = R[\text{tp}(\bar{c}, \bar{b}), \Delta, \aleph_0] = R[\text{stp}(\bar{c}, \bar{b}), \Delta, \aleph_0]$$

and

$$R[\text{tp}(\bar{c}, \bar{b} \cup \{\bar{c}_i : i \leq n\} \cup \{\bar{c}_j^* : j < m\}), \Delta, k] < R[\text{tp}(\bar{c}, \bar{b}), \Delta, k].$$

These properties for fixed  $\bar{c}, \bar{b}$  are expressed by first-order formulas, i.e. there are formulas which  $\bar{b}, \bar{c}, \bar{c}_i, \bar{c}_j^*$  satisfy and imply this (see II 2.19). So by 1.2(2) we can define  $\bar{c}'_0, \dots, \bar{c}'_n \in N$  and then define  $\bar{c}_j^{**} \in M$  ( $j < m$ ) such that

- (i)  $\bar{c}'_0, \dots, \bar{c}'_n$  realizes  $\text{tp}(\bar{c}, \bar{b})$ .
- (ii)  $\{\bar{c}'_0, \dots, \bar{c}'_n\}$  is independent over  $\bar{b}$ .
- (iii)  $\psi(\bar{c}_j^{**}, \bar{b})$  for ( $j < m$ ).
- (iv)  $R^m[\text{tp}(\bar{c}_{n_j}, \bar{b} \cup \{\bar{c}'_i : i \leq n_j\} \cup \{\bar{c}_j^{**} : j < m\}), \Delta, k] < R^m[\text{stp}(\bar{c}, \bar{b}), \Delta, k]$ .
- (v)  $R^m[\text{tp}(\bar{c}, \bar{b} \cup \{\bar{c}'_i : i \leq n\} \cup \{\bar{c}_j^{**} : j < m\}), \Delta, k] < R^m[\text{stp}(\bar{c}, \bar{b}), \Delta, k]$ .

By (v),  $\text{tp}(\bar{c}, \bar{b} \cup \{\bar{c}'_i : i \leq n\} \cup \{\bar{c}_j^{**} : j < m\}) \subseteq \text{tp}(\bar{c}, M' \cup \{\bar{c}_j^{**} : j < m\})$  forks over  $\bar{b}$ . But  $\text{tp}(\bar{c}, M')$  does not fork over  $\bar{b}$ , hence  $\bar{c}_{j(0)}^{**} \notin M'$  for some  $j(0)$ . As  $\text{tp}(\bar{c}, M')$ ,  $\text{tp}(\bar{c}_{j(0)}^{**}, M')$  are not orthogonal, and the first does not fork over  $N$ , the second is not orthogonal to  $N$ . For notational simplicity assume  $n = n_{j(0)}$ .

By (iv) (and (i))  $\text{tp}(\bar{c}'_n, \bar{b} \cup \{\bar{c}'_i : i < n\} \cup \{\bar{c}_{j(0)}^{**}\})$  forks over  $\bar{b}$ .

(Note that as  $\bar{c}'_n, \bar{c}$  realizes the same type over  $\bar{b}$ ,

$$R^m[\text{tp}(\bar{c}'_n, \bar{b}), \Delta, k] = R^m[\text{stp}(\bar{c}'_n, \bar{b}), \Delta, k] = R^m[\text{tp}(\bar{c}, \bar{b}), \Delta, k].$$

Hence by (ii)  $\text{tp}(\bar{c}'_n, \bar{b} \cup \{\bar{c}'_i : i < n\} \cup \{\bar{c}_{j(0)}^{**}\})$  forks over  $\bar{b} \cup \{\bar{c}'_i : i < n\}$ . Hence  $\text{tp}(\bar{c}_{j(0)}^{**}, \bar{b} \cup \{\bar{c}'_i : i \leq n\})$  forks over  $\bar{b} \cup \{\bar{c}'_i : i < n\}$  and hence over  $\bar{b}$ . So

$$\begin{aligned} R[\text{tp}(\bar{c}_{j(0)}^{**}, M'), L, \infty] &\leq R[\text{tp}(\bar{c}_{j(0)}^{**}, \bar{b} \cup \{\bar{c}'_i : i \leq n\}), L, \infty] < R^m[\text{tp}(\bar{c}_{j(0)}^{**}, \bar{b}), L, \infty] \\ &\leq R^m[\psi(\bar{x}, \bar{b}), L, \infty]. \end{aligned}$$

This contradicts the minimality of the rank of  $\varphi(x, \bar{b})$ . Hence  $\text{stp}(\bar{c}, \bar{b})$  is regular, and we are finished.

1.7. CLAIM. (1) For any  $A \subseteq M$  there is  $N \subseteq_a M$ , such that  $A \subseteq N$ ,  $\|N\| \leq A + \lambda(T)$ .

(2) For any  $A \subseteq M$  there is  $N \subseteq_e M$  such that  $a \subseteq N$ ,  $\|N\| \leq |A| + |D(T)|$ .

PROOF. Trivial.

REMARK. We can replace  $M$  by  $B$ . The next lemma will not be used in the sequel.

1.8. LEMMA. ( $\mathfrak{C}^{\text{eq}}$ ) Suppose  $N \subseteq A \subseteq M$ ,  $N \subseteq_a M$ ,  $\bar{a} \in M$ ,  $\bar{a} \notin A$ ,  $r = \text{stp}(\bar{a}, A)$  is regular and trivial, but not orthogonal to  $N$ . Then there is a  $\bar{a}' \in M$  such that  $\text{stp}(\bar{a}', A)$  is regular but not orthogonal to  $\text{stp}(\bar{a}, A)$  and does not fork over  $N$ .

PROOF. W.l.o.g.  $A = \text{acl } A$ . By 0.2 there is  $b \in \text{acl}(A \cup \bar{a}) - \text{acl}(A)$  (hence  $b \in M$ ) such that:  $\text{tp}(b, A)$  is regular but not orthogonal to  $\text{stp}(\bar{a}, A)$  (and hence it is trivial, too) and some  $\varphi(x, \bar{c}_0) \in \text{tp}(b, A)$  is  $\text{cl}^3(r)$ -regular. W.l.o.g.  $\text{tp}(\bar{b}, A)$  does not fork over  $\bar{c}_0$ . Choose  $\bar{d} \in N$  such that  $\text{tp}(\bar{c}_0 \wedge \langle b \rangle, N)$  does not fork over  $\bar{d}$ . Now as  $N \subseteq_a M$  we can choose  $\bar{c}'_0$ , and  $b' \in N$  such that  $\text{stp}(\bar{c}'_0 \wedge \langle b' \rangle, \text{dbar}) \equiv \text{stp}(\bar{c}_0 \wedge \langle b \rangle, \bar{d})$ . By V. 3.4,  $\text{stp}(b, \bar{c}_0)$ ,  $\text{stp}(b', \bar{c}'_0)$  are not orthogonal. By [2] 5.11 (or more elaborately [5] X 7.1) there is  $b''$  realizing  $\text{stp}(b', \bar{c}'_0)$  such that  $\text{tp}(b, \bar{c}_0 \cup \bar{c}'_0 \cup \bar{d} \cup \{b''\})$  forks over  $\bar{c}_0 \cup \bar{c}'_0 \cup \bar{d}$ . Hence easily there is  $b^* \in M$  satisfying  $\varphi(x, \bar{c}'_0)$  such that  $\text{tp}(b, \bar{c}_0, \bar{c}'_0 \cup \bar{d} \cup \{b^*\})$  forks over  $\bar{c}_0 \cup \bar{c}'_0 \cup \bar{d}$ .

Hence  $\text{tp}(b, A \cup \{b^*\})$  forks over  $A$ , and thus  $\text{tp}(b^*, A)$  is not orthogonal to  $\text{tp}(b, A)$  (and  $\text{tp}(\bar{a}, A)$ ). As  $\varphi(x, \bar{c}_0)$  is  $\text{cl}^3(r)$ -regular,  $\varphi(x, \bar{c}'_0)$  is  $\text{cl}^3(r)$ -regular also. So as  $\text{tp}(b^*, A)$  is not orthogonal to  $r$ , it is regular, also it does not fork over  $\bar{c}'_0$  hence over  $N$ .

## §2. A prime atomic model over stable amalgamation is enough

HYPOTHESIS.  $T$  is superstable, and if  $\{M_1, M_2\}$  is independent over  $M$ ,  $M \subseteq M_l$  ( $l = 1, 2$ ) then there is a model  $N$ ,  $\mathbf{F}'_{\kappa_0}$ -prime and  $\mathbf{F}'_{\kappa_0}$ -atomic over  $M_1 \cup M_2$ .

2.1. LEMMA. Suppose  $\{N_\eta : \eta \in I\}$  is a non-forking tree (see III or [2] 3.2). Then there is a model  $N$ , which is  $\mathbf{F}'_{\kappa_0}$ -prime and  $\mathbf{F}'_{\kappa_0}$ -atomic over  $\bigcup_{\eta \in I} N_\eta$ .

PROOF. Let  $I = \{\eta_\alpha : \alpha < |I|\}$  be such that for every  $\alpha$  and  $k < l(\eta_\alpha)$ ,

$\eta_\alpha \upharpoonright k \in \{\eta_\beta : \beta < \alpha\}$ . We define by induction, on  $\alpha > 0$ ,  $M_\alpha$  such that

$$(*) \quad \left\{ \begin{array}{l} (1) \quad M_\alpha \text{ is } \mathbf{F}'_{\aleph_0}\text{-prime and } \mathbf{F}'_{\aleph_0}\text{-atomic over } \bigcup_{\beta < \alpha} N_{\eta_\beta} \\ \text{and even over } M_\gamma \cup \bigcup_{\beta < \alpha} N_{\eta_\beta} \text{ for each } \gamma < \alpha. \\ (2) \quad M_i \text{ (} i \leq \alpha \text{) is increasing continuous.} \end{array} \right.$$

For  $\alpha = 1$ , let  $M_\alpha = N_{\eta_0}$ ; for  $\alpha$  a limit take the union, and for  $\alpha = \beta + 1$  use the hypothesis ( $M, M_1, M_2, N$  there correspond to  $N_{\eta_\alpha} \upharpoonright (I(\eta_\alpha) - 1)$ ,  $N_{\eta_\alpha}$ ,  $M_\beta$ ,  $M_\alpha$  here). Why does this work? Note that  $(\bigcup_{\gamma < \alpha} N_{\eta_\gamma}, \bigcup_{\gamma < |I|} N_{\eta_\gamma})$  satisfies the Tarski-Vaught condition (see 3.2A below). Hence if  $M_\alpha$  is  $\mathbf{F}'_{\aleph_0}$ -prime and  $\mathbf{F}'_{\aleph_0}$ -atomic over  $\bigcup_{\beta < \alpha} N_{\eta_\beta}$ , then necessarily  $\{M_\alpha, \bigcup_{\gamma < |I|} N_{\eta_\gamma}\}$  is independent over  $\bigcup_{\beta < \alpha} N_{\eta_\beta}$ , and if  $\bigcup_{\alpha < |I|} N_{\eta_\gamma} \subseteq N$ ,  $F$  an embedding of  $M_\alpha$  into  $N$ ,  $F \upharpoonright \bigcup_{\beta < \alpha} N_{\eta_\beta} =$  the identity, then  $F \cup G$  is an elementary embedding, where  $G$  is the identity map of  $\bigcup_{\gamma < |I|} N_{\eta_\gamma}$ .

**2.2. THE ATOMIC DECOMPOSITION LEMMA** (in  $\mathfrak{C}^{\text{eq}}$ ). *Suppose  $T$  is superstable with the dop. Then for any pair of models  $N_1 \subseteq_a M$ , there are elements  $a_i \in M$  ( $i < \alpha$ ) and models  $M_i$  such that:*

- (a)  $N_1 \subseteq_a M_i \subseteq_a M$ ;
- (b)  $\text{tp}(a_i, N_1)$  is regular;
- (c)  $|M_i| = N_1 \cup \{a_i\} \cup \{b_{i,\alpha} : \alpha < \alpha_i\}$ , for every  $\alpha$ ,  $b_{i,\alpha} \notin A_{i,\alpha}$  and one of the following occurs (letting  $A_{i,\alpha} = N_1 \cup \{a_i\} \cup \{b_{i,\beta} : \beta < \alpha\}$ ):
  - (c1)  $\text{tp}(b_{i,\alpha}, A_{i,\alpha})$  is  $\mathbf{F}'_{\aleph_0}$ -isolated,
  - (c2)  $\text{tp}(b_{i,\alpha}, A_{i,\alpha})$  is orthogonal to  $N_1$ ;
- (d) for no  $b \in M - M_i$  is  $\text{tp}(b, M_i)$  orthogonal to  $N_1$ ;
- (e)  $M$  is  $\mathbf{F}'_{\aleph_0}$ -prime and  $\mathbf{F}'_{\aleph_0}$ -atomic over  $\bigcup_{i < \alpha} M_i$  (and  $\mathbf{F}'_{\aleph_0}$ -minimal);
- (f)  $\{M_i : i < \alpha\}$  is independent over  $N_1$ .

**PROOF.** Let  $I = \{a_i : i < \alpha^*\}$  be a maximal subset of  $M$  independent over  $N_1$ , of elements realizing regular types of  $N_1$  and for each  $i < \alpha^*$  define  $b_{i,\alpha}$ ,  $M_i$ ,  $\alpha_i$  as required in (c),  $b_{i,\alpha} \in M - A_{i,\alpha}$  and  $\alpha_i$  is maximal. So (b), (c) hold trivially.

Why is  $|N_1| \cup \{a_i\} \cup \{b_{i,\alpha} : \alpha < \alpha_i\}$  the universe of a submodel (elementary, of course)? See IV 2.21. Now (d) follows from the choice of  $\alpha_i$ .

Clearly (f) follows by 1.3(2) and (a) by 1.2(3) provided that (e) holds. Apply 2.1. Let  $M'$  be  $\mathbf{F}'_{\aleph_0}$ -prime  $\mathbf{F}'_{\aleph_0}$ -atomic model over  $\bigcup_{i < \alpha} M_i$ . So w.l.o.g. (3)  $M' \subseteq M$ .

The only missing point is  $M' = M$ .

If not, there is  $c \in M - M'$ ,  $R[\text{tp}(c, M'), L, \infty]$  is minimal. Then by 1.5:  $\text{tp}(c, M')$  is orthogonal to  $N_1$ , or  $\text{tp}(c, M')$  does not fork over  $N_1$ , hence by 1.6

$\text{tp}(c, N_1)$  is regular. The latter case contradicts the maximality of **I**. In the former case, we can find  $N_1^*, M_i^*$  such that:  $N_1 \subseteq N_1^*, M_i \subseteq M_i^*$  are  $\mathbf{F}_{\aleph_0}^a$ -saturated,  $\text{tp}(N_1^*, \bigcup_{i < \alpha} M_i)$  does not fork over  $N_1$ ,  $\text{tp}(M_i^*, \bigcup_{j \neq i} M_j^* \cup N_1^*)$  does not fork over  $N_1^* \cup M_i$ . By 3.3 (next section) the pair  $(\bigcup_{i < \alpha} M_i, \bigcup_{i < \alpha} M_i^*)$  satisfies the Tarski-Vaught condition.

Let  $\text{tp}(c, M')$  not fork over some  $\bar{a} \in M'$ , then  $\text{tp}(\bar{a}, \bigcup_{i < \alpha} M_i)$  is  $\mathbf{F}_{\aleph_0}^a$ -isolated, hence  $\text{tp}(\bar{a}, \bigcup_{i < \alpha} M_i^*)$  is  $\mathbf{F}_{\aleph_0}^a$ -isolated. By [2] §2 (as  $T$  does not have the dop) for some  $i$ ,  $\text{tp}(c, M')$  (equivalently,  $\text{stp}(c, \bar{a})$ ) is not orthogonal to  $M_i^*$ , hence (as  $\{M_i^*, \bar{a}\}$  is independent over  $M_i$ ), the type is not orthogonal to  $M_i$ . For notational simplicity let  $i = 0$ . As  $M_0 \subseteq_a M'$  (by 1.2(3), as mentioned above), we can apply 1.5 and find  $c' \in M - M'$ ,  $\text{tp}(c', M')$  does not fork over  $M_0$ . If  $\text{tp}(c', M_0)$  is not orthogonal to  $N_1$ , we can get a contradiction to the maximality of **I**. If  $\text{tp}(c', M_0)$  is orthogonal to  $N_1$  we get a contradiction to the choice of  $M_0$ .

**2.2A. ASSERTION.** We can add the demand: for each  $i$ ,  $\text{tp}(b_{i,\alpha}, A_{i,\alpha})$  is orthogonal to every trivial regular type not orthogonal to  $N_1$ .

**PROOF.** The only problem is when  $A_{i,\alpha_i}$  is not the universe of an  $M_i$ . Let  $\alpha = \alpha_i$ . If not there is a formula  $\varphi(x, \bar{b})$ ,  $\bar{b} \in A_{i,\alpha}$ ,  $\models (\exists x)\varphi(x, \bar{b})$  but for no  $c \in A_{i,\alpha}$ ,  $\models \varphi(c, \bar{b})$ . Choose such  $\varphi(x, \bar{b})$  with minimal  $R[\varphi(x, \bar{b}), L, \infty]$ , hence every  $q$ ,  $\varphi(x, \bar{b}) \in q \in S^1(A_{i,\alpha})$  is  $\mathbf{F}_{\aleph_0}^c$ -isolated. Let  $c \in M$ ,  $\models \varphi[c, \bar{b}]$ , so  $\text{tp}(c, A_{i,\alpha})$  is  $\mathbf{F}_{\aleph_0}^c$ -isolated. So necessarily some trivial regular  $r$  is not orthogonal to  $\text{stp}(c, A_{i,\alpha})$  and not orthogonal to  $N_1$ . We can find  $\bar{d} \in N_1$  such that  $\text{tp}(\bar{b} \wedge \langle c \rangle, N_1)$  does not fork over  $\bar{d}$ , and  $r$  is not orthogonal to  $\bar{d}$ . As  $N_1 \subseteq_a M$  there are  $\bar{b}', c' \in N_1$  such that  $\bar{b}' \wedge \langle c' \rangle$  realizes  $\text{stp}(\bar{b} \wedge \langle c \rangle, \bar{d})$ . It is easy to see that  $r$  is not orthogonal to  $\text{stp}(c, \bar{b}')$ . So by [2] 5.11 or [5] X 7.1 there is  $c''$  realizing  $\text{stp}(c, \bar{d} \cup \bar{b} \cup \bar{b}')$  such that  $\text{tp}(c'', \bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c'\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$ . Hence  $\text{tp}(c', \bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c''\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$ , so some  $c^* \in M$ ,  $\models \varphi(c^*, \bar{b})$  and  $\text{tp}(c', \bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c^*\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$  hence  $\text{tp}(c^*, \bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c'\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$ . As  $\bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c'\} \subseteq A_{i,\alpha}$ , we get a contradiction to the choice of  $\varphi(x, \bar{b})$ .

**REMARK.** We have essentially used (and proved):

**2.2B. FACT.** Suppose  $A \subseteq B \subseteq M$ , and  $\text{tp}(\bar{a}, A)$  has an extension over  $B$  which forks over  $A$ . Then for every  $\varphi(x, \bar{b}) \in \text{tp}(\bar{a}, A)$  there is  $\bar{a}' \in |M|$  such that  $\models \varphi[\bar{a}', \bar{b}]$  and  $\text{tp}(\bar{a}', B)$  forks over  $A$ .

**2.3. CLAIM.** In 2.2, 2.2A:

(1) If  $N_{2,i} \subseteq M_i$ ,  $N_1 \cup \{a_i\} \subseteq N_{2,i}$  then  $N_{2,i} <_{N_1} M_i$ .

(2) If  $p \in S^m(M_i)$  is orthogonal to  $N_i$  and  $\text{Dp}(\text{tp}(a_i, N_i)) < \infty$ , then  $\text{Dp}(p) < \text{Dp}(\text{tp}(a_i, N_i))$ .

(3) If  $\bar{a} \in M_i$ ,  $\text{Dp}(\text{tp}(a_i, N_i)) < \infty$ ,  $N_{2,i}$  as in (1) then

$$\text{Dp}(\text{tp}(\bar{a}, N_{2,i})) < \text{Dp}(\text{tp}(a_i, N_i)).$$

PROOF. (1) By 1.5, 1.6 and  $I$ 's maximality.

(2) Let  $N_1^*$  be  $\mathbf{F}_{\kappa_0}^a$ -saturated,  $N_1 \subseteq N_1^*$ ,  $\{N_1^*, M_i\}$  independent over  $N_1$ . Now for every  $\alpha < \alpha_i$  (see 2.2)  $\text{tp}(b_{i,\alpha}, A_{i,\alpha})$  is orthogonal to every regular  $p \in S^m(N_1^*)$  with depth  $\geq 1$ . [If  $p$  is orthogonal to  $N_1$  then it is orthogonal to any type over  $N_1^*$ . Suppose  $p$  is not orthogonal to  $N_1$ , then by 2.2A, if  $\text{tp}(b_{i,\alpha}, A_{i,\alpha})$  is not orthogonal to  $p$  then  $p$  is not trivial. We finish remembering that by [2] 5.10, as  $T$  does not have the dop, any regular type of depth  $> 0$  is trivial.] So clearly  $\text{tp}_*(M_i, N_1 \cup \{a_i\})$  is orthogonal to every regular complete type of  $N_1^*$  of depth  $> 0$ . We can find  $\mathbf{F}_{\kappa_\alpha}^a$ -saturated  $M_i^*$ ,  $N_1^* \cup M_i \subseteq M_i^*$  with this property. We can apply [2] 3.2; so the conclusion of 2.3(2), (3) follows.

(3) See the proof of (2).

2.4. THE DECOMPOSITION LEMMA. Suppose  $T$  is superstable without the dop. Then for any model  $M$  there is a tree  $I(\subseteq^{w>} \|M\|)$  and  $N_\eta$  ( $\eta \in I$ ),  $a_\eta$  ( $\eta \in I$ ) such that:

- (1)  $N_\eta \subseteq_a M$  (hence  $N_\eta < M$ );
- (2)  $N_\eta \subseteq_a \cup \{N_\nu : \eta \leq \nu\}$ ;
- (3)  $p_{\eta \wedge (i)} = \text{tp}(a_{\eta \wedge (i)}, N_\eta)$  is regular;
- (4)  $\text{tp}_*(\cup \{N_\nu : \eta \wedge (i) \leq \nu, \nu \in I\} \cup \{N_\nu : \text{not } \eta \wedge (i) \leq \nu, \nu \in I\})$  does not fork over  $N_\eta$ ;
- (5)  $\{a_{\eta \wedge (i)} : \eta \wedge (i) \in I\}$  is a maximal subset of  $M$  independent of  $N_\eta$ ;
- (6)  $\text{tp}_*(\cup \{N_\nu : \eta \wedge (i) \leq \nu \in I\}, N_{\eta \wedge (i)})$  is orthogonal to  $N_\eta$ ;
- (7) if  $\text{Dp}(p_\eta) < \infty$ ,  $\eta \wedge (i) \in I$  then  $\text{Dp}(p_{\eta \wedge (i)}) < \text{Dp}(p_\eta)$ .

PROOF. Just combine the proofs of [2] 3.2 and 2.2 (and 2.3(2), (3)).

Now it is no problem to compute the number of non-isomorphic models, as in [2], [3] (using the same depth function).

### §3. Universal Theories

3.1. DEFINITION. (1) We call  $\text{cl}$  a closure operation if:

- (i) for every  $A$ ,  $A \subseteq \text{cl } A \subseteq \text{acl } A$ , and for every function symbol  $F$  (of  $\mathfrak{U}$ ) and  $\bar{a} \in A$ ,  $F^{\mathfrak{U}}(\bar{a}) \in \text{cl } A$ ;
- (ii)  $\text{cl}(\text{cl}(A)) = \text{cl } A$ , and  $A \subseteq B$  implies  $\text{cl } A \subseteq \text{cl } B$ ;

- (iii) the property " $a \in \text{cl } A$ " is preserved by an automorphism of  $\mathfrak{C}$ .
- (2) We call a closure operation  $\text{cl}$  local if in addition
- (iv) for every  $b \in \text{cl } A$ , there are a formula  $\varphi(x, \bar{y})$  and a sequence  $\bar{a} \in A$  such that  $\models \varphi[b, \bar{a}]$  and:  $\varphi(b_1, \bar{a}_1)$  implies  $b_1 \in \text{cl}(\bar{a}_1)$ .
- (3) For a set of formulas  $\Phi$  (of the form  $\varphi(\bar{x}, \bar{y})$ ) let  $\text{acl}_\Phi$  be defined by:

$$\text{acl}_\Phi^1(A) = \bigcup \{ \bar{b} : \text{for some } \bar{a} \in A, \text{ and } \varphi(\bar{x}, \bar{y}) \in \Phi, \models \varphi[\bar{b}, \bar{a}], \\ \text{and } \varphi(\bar{x}, \bar{a}) \text{ is an algebraic formula} \} \cup A,$$

$$\text{acl}_\Phi^0(A) = A,$$

$$\text{acl}_\Phi^{n+1}(A) = \text{acl}_\Phi^1(\text{acl}_\Phi^n(A)),$$

$$\text{acl}_\Phi(A) = \bigcup_n \text{acl}_\Phi^n(A).$$

- (4) We call  $\text{cl}$  a  $\Phi$ -closure operation if it is an operation and  $A \subseteq \text{cl } A \subseteq \text{acl}_\Phi A$ .

3.1A. CLAIM. (1) Every  $\text{acl}_\Phi$  is a local closure operation and is  $\text{acl}_\Psi^1$  for some  $\Psi$ .

(2) Every local closure operation is a  $\Phi$ -closure operation for some  $\Phi$ , and in Definition 3.1(2) there is a  $\varphi$  satisfying in addition: for some  $n$ ,  $\models (\forall \bar{y})(\exists^{\leq n} \bar{x})\varphi(\bar{x}, \bar{y})$  and  $\varphi(\bar{x}, \bar{a}) \vdash \text{tp}(\bar{b}, \bar{a})$ .

PROOF. Easy.

3.2. DEFINITION. We call  $\langle M_s : s \in I \rangle$  a stable system if  $M_s \subseteq \mathfrak{C}$ ,  $I$  a family of finite subsets of  $\bigcup_{s \in I} s$  closed under subsets,

$$s < t \Rightarrow M_s \subseteq M_t$$

and for every  $s \in I$ ,  $\text{tp}_*(M_s, \bigcup_{s \not\supseteq t} M_t)$  does not fork over

$$A_s \stackrel{\text{def}}{=} \bigcup \{ M_t : t \subseteq s, t \neq s \}.$$

We implicitly assume  $\text{Th}(\mathfrak{C})$  is stable.

3.3. CLAIM. (1) If  $I = \{s_\alpha : \alpha < \alpha_0\}$ ,  $[s_\alpha \subseteq s_\beta \Rightarrow \alpha \leq \beta]$ ;  $M_s < C$ , and  $\text{tp}_*(M_{s_\alpha}, \bigcup_{j < \alpha} M_{s_j})$  does not fork over  $A_{s_\alpha}$ , then  $\langle M_s : s \in I \rangle$  is a stable system of models

(2) If  $\langle M_s : s \in I \rangle$  is a stable system,  $J \subseteq I$  and  $s \in I \wedge s \subseteq \bigcup J \Rightarrow s \in J$ , then  $(\bigcup_{s \in J} M_s, \bigcup_{s \in I} M_s)$  satisfies the Tarski-Vaught condition (i.e. if  $\bar{a} \in \bigcup_{s \in J} M_s$ ,  $\bar{b} \in \bigcup_{s \in I} M_s$ ,  $\mathfrak{C} \models \varphi[\bar{a}, \bar{b}]$  then for some  $\bar{b}' \in \bigcup_{s \in J} M_s$ ,  $\mathfrak{C} \models \varphi[\bar{a}, \bar{b}']$ ).

(3) If  $\langle M_s : s \in I \rangle$  is a stable system,  $N_s < M_s$ ,  $\text{tp}(N_s, \bigcup_{t \subseteq s, t \neq s} M_t)$  does not fork

over  $\bigcup_{i \subseteq s, t \neq s} N_i$ ,  $[s < t \Rightarrow N_s \subseteq N_t]$  then  $\langle N_s : s \in I \rangle$  is a stable system and  $(\bigcup_{s \in I} N_s, \bigcup_{s \in I} M_s)$  satisfies the Tarski-Vaught condition.

PROOF. Essentially like [4] 3.5. Since we do not want to assume that the reader is familiar with [4], we prove the claim completely:

3.3A. FACT. If  $\langle M_s : s \in I \rangle$  is a stable system  $J_0 \subseteq I$ ,  $J \subseteq I$ ,  $J_0$  is closed under subsets then  $\text{tp}_*(\bigcup_{i \in J} M_i, \bigcup_{s \in J_0} M_s)$  does not fork over  $\bigcup \{M_s : s \in J_0 \text{ and } (\exists t \in J) s \subseteq t\}$ .

REMARK. If  $J$  is closed under subsets, the last set is  $\bigcup \{M_s : s \in J_0 \cap J\}$ .

PROOF. W.l.o.g.  $J$  is closed under subsets, and let  $J_1 = J \cap J_0$ . We can find a list  $\{s_\alpha : \alpha < \alpha^*\}$  of  $I$  (and  $\alpha_1 \leq \alpha_2 \leq \alpha_3 \leq \alpha$ ) such that  $[s_\alpha \subseteq s_\beta \Rightarrow \alpha \leq \beta]$ ,  $J_0 = \{s_\alpha : \alpha < \alpha_2\}$ ,  $J = \{s_\alpha : \alpha < \alpha_1, \text{ or } \alpha_2 \leq \alpha < \alpha_3\}$ . Clearly for  $\alpha < \alpha_3$ ,  $\alpha \geq \alpha_2$ ,  $\text{tp}_*(M_{s_\alpha}, \bigcup_{\beta < \alpha} M_{s_\beta})$  is included in  $\text{tp}_*(M_{s_\alpha}, \bigcup \{M_t : s_\alpha \not\subseteq t, t \in I\})$  hence does not fork over  $A_{s_\alpha}$ , but  $A_{s_\alpha} \subseteq \bigcup \{M_{s_\beta} : \beta < \alpha_1 \text{ or } \alpha_2 \leq \beta < \alpha\}$ . So  $\text{tp}_*(M_{s_\alpha}, \bigcup_{\beta < \alpha} M_{s_\beta})$  does not fork over  $\bigcup \{M_{s_\beta} : \beta < \alpha_1, \text{ or } \alpha_2 \leq \beta < \alpha\}$ . By IV 3.2(1) we can conclude that  $\text{tp}_*(\bigcup \{M_{s_\alpha} : \alpha_2 \leq \alpha < \alpha_3 \text{ or } \alpha < \alpha_1\}, \bigcup_{\beta < \alpha_2} M_{s_\beta})$  does not fork over  $\bigcup_{\beta < \alpha_1} M_{s_\beta}$ , but this is as required.

3.3B. FACT. If  $S = \langle M_s : s \in I \rangle$  is a stable system  $\bar{a}_l \in M_{s(l)}$  ( $l < n$ ),  $t \subseteq \bigcup I$  and  $\models \varphi[\bar{a}_0, \dots, \bar{a}_{n-1}]$  then we can find  $\bar{a}'_l \in M_{s(l) \cap t}$  such that  $\models \varphi[\bar{a}'_0, \dots, \bar{a}'_{n-1}]$  and  $s(l) \subseteq t \Rightarrow a'_l = a_l$ .

PROOF. W.l.o.g.  $s \subseteq s(l) \Rightarrow s \in \{s(m) : m < l\}$ . We prove it by induction on  $n$ . For  $n = 0$  there is nothing to prove, and for  $n = 1$  note  $M_{s(l) \cap t}$  is an elementary submodel  $M_{s(l)}$ . So suppose we have proved for  $n$  and we shall prove for  $n + 1$ , i.e. for given  $\bar{a}_l \in M_{s(l)}$  ( $l < n + 1$ ),  $t \subseteq \bigcup I$  and  $\varphi$ . W.l.o.g. the  $s(l)$  ( $l \leq n$ ) are distinct and  $s(n) \not\subseteq s(l)$  for  $l < n$ . We concentrate on the case  $s(n) \not\subseteq t$ . As  $\text{tp}(\bar{a}_{s(n)}, \bigcup_{l < n} M_{s(l)})$  does not fork over  $A_{s(n)}^s$  clearly  $\varphi(\bar{a}_0, \dots, \bar{a}_{n-1}, \bar{x})$  does not fork over  $A_{s(n)}^s$  hence is realized in every model which includes  $A_{s(n)}^s$ . So for some type  $p = p(\bar{x}_i)_{i < \alpha}$  over  $A_{s(n)}^s$  (infinitely many variables)  $p(\bar{x}_0, \dots, \bar{x}_i, \dots) \vdash \bigvee_{i \in \alpha} \varphi(\bar{a}_0, \dots, \bar{a}_{n-1}, \bar{x}_i)$ . So for some  $\bar{b} \subseteq A_{s(n)}^s$  and  $\psi = \psi(\bar{x}_0, \bar{x}_1, \dots, \bar{x}_k, \bar{b})$

$$\models (\exists \bar{x}_0, x_1, \dots, x_k) \psi(\bar{x}_0, \dots, \bar{x}_k, \bar{b}),$$

$$\psi(\bar{x}_0, \dots, \bar{x}_k, \bar{b}) \vdash \bigvee_{i \leq k} \varphi(\bar{a}_0, \dots, \bar{a}_{n-1}, \bar{x}_i).$$

As  $\bar{b} \subseteq A_{s(n)}^s$  ( $\forall s \subseteq s(n)$ )  $[s \in \{s(l) : l < n\}]$ , w.l.o.g.  $\bar{b} = \bar{b}_0 \wedge b_1 \wedge \dots \wedge \bar{b}_{n-1}$ ,  $\bar{b}_l \subseteq M_{s(l)}$ ,  $[s(l) \not\subseteq s(n) \Rightarrow \bar{b}_l \text{ empty}]$ . Now apply the induction hypothesis to  $\bar{a}_i \wedge \bar{b}_i \in$

$M_{s(l)}$  (for  $l < n$ ) and the formula

$$(\exists \bar{x}_0, \dots, \bar{x}_k) \psi(\bar{x}_0, \dots, \bar{x}_k, \bar{b}_0, \dots, \bar{b}_{n-1}) \wedge (\forall \bar{x}_0, \dots, \bar{x}_k) \\ \left[ \psi(\bar{x}_0, \dots, \bar{x}_k, \bar{b}_0, \dots, \bar{b}_{n-1}) \rightarrow \bigvee_{i \leq k} \varphi(\bar{a}_0, \dots, \bar{a}_{n-1}, \bar{x}_i) \right].$$

So there are  $\bar{a}'_l \wedge \bar{b}'_l \in M_{s(l) \cap t}$  ( $l < n$ ) satisfying the above formula (and as in 3.3B). Now clearly  $\bar{b}'_0 \wedge \bar{b}'_1 \wedge \dots \wedge \bar{b}'_{n-1} \subseteq A_{s(n) \cap t}^s$ , hence there are  $\bar{c}_0, \dots, \bar{c}_k \in M_{s(n) \cap t}$  such that  $\models \psi[\bar{c}_0, \dots, \bar{c}_k, \bar{b}'_0, \dots, \bar{b}'_{n-1}]$ . So for some  $i \leq k \models \varphi[\bar{a}'_0, \dots, \bar{a}'_{n-1}, \bar{c}_i]$ . So  $\bar{a}'_0, \dots, \bar{a}'_{n-1}, \bar{a}'_n \stackrel{\text{def}}{=} \bar{c}_i$  are as required.

PROOF OF 3.3. (1) An exercise in non-forking.

(2) Follows from Fact 3.3B.

(3) First we prove that  $\langle N_s : s \in I \rangle$  is a stable system. For every  $s$ ,  $\text{tp}_*(M_s, \cup \{M_t : M_t \in I, s \not\subseteq t\})$  does not fork over  $\cup \{M_t : t \subseteq s, t \neq s\}$ , hence (as  $N_s \subseteq M_s$ ) also  $\text{tp}_*(N_s, \cup \{M_t : t \in I, s \not\subseteq t\})$  does not fork over  $\cup \{M_t : t \subseteq s, t \neq s\}$ . But  $\text{tp}_*(N_s, \cup \{M_t : t \subseteq s, t \neq s\})$  does not fork over  $\cup \{N_t : t \subseteq s, t \neq s\}$ . So by III 0.1,  $\text{tp}_*(N_s, \cup \{M_t : t \in I, s \not\subseteq t\})$  does not fork over  $\cup \{N_t : t \subseteq s, t \neq s\}$ . As  $N_i \subseteq M_i$ , by monotonicity of non-forking we get the stability of the system ( $s \subseteq t \Rightarrow N_s \subseteq N_t$  was assumed, and we know  $I$  is as required).

The Tarski–Vaught condition follows by Fact 3.3B and the following fact.

Let  $j \notin \cup I$ ,  $J = I \cup \{s \cup \{j\} : s \in I\}$ , and  $N_{s \cup \{j\}} = M_s$ .

3.3C. FACT.  $\langle N_s : s \in J \rangle$  is a stable system ( $J, N_s$  as above).

PROOF. Let  $s_\alpha$  ( $\alpha < \alpha_0$ ) be as in 3.3(1), and define  $t_\alpha$  ( $\alpha < 2\alpha_0$ ) by:  $t_{2\alpha} = s_\alpha, t_{2\alpha+1} = s_\alpha \cup \{j\}$ . Clearly  $J = \{t_\alpha : \alpha < 2\alpha_0\}$  and  $t_\alpha \subseteq t_\beta \Rightarrow \alpha \leq \beta$ . Now use 3.3(1): For  $\alpha$  even ( $= 2\beta$ ) remember we have proved  $\text{tp}_*(N_{s_\beta}, \cup \{M_s : s \in I, s_\beta \not\subseteq s\})$  does not fork over  $\cup \{N_s : s \subseteq s_\beta, s \neq s_\beta\}$  and this is what we need. For  $\alpha$  odd ( $= 2\beta + 1$ ) remember  $\text{tp}_*(M_{s_\beta}, \bigcup_{\gamma < \beta} M_{s_\gamma})$  does not fork over  $\cup \{M_s : s \subseteq s_\beta, s \neq s_\beta\}$ . As  $N_{s_\beta} \subseteq M_{s_\beta}$ , by III 0.1 this gives  $\text{tp}_*(N_{t_\alpha}, \bigcup_{\gamma < \alpha} N_{t_\gamma}) = \text{tp}_*(M_{s_\beta}, \bigcup_{\gamma < \beta} M_{s_\gamma} \cup N_{s_\beta})$  does not fork over  $\cup \{M_s : s \subseteq s_\beta, s \neq s_\beta\} \cup N_{s_\beta} = \cup \{N_s : s \subseteq t_\alpha, s \neq t_\alpha\}$ , and this is what we need.

3.3D. FACT. Suppose  $\langle M_s : s \in I \rangle$  is a stable system, and each  $M_s$  is  $\kappa$ -compact,  $a_{i,j} \in M_{s(i)}$  for  $i < \alpha, j < j_\alpha, \sum_\alpha j_\alpha < \kappa$ , and  $p$  is a set of  $< \kappa$  formulas, in the variables  $x_{i,j}$  ( $i < \alpha, j < j_\alpha$ ) satisfied by the assignment  $x_{i,j} \mapsto a_{i,j}$ . If  $t \subseteq \cup I$  (not necessarily finite) then we can find  $a'_{i,j} \in M_{s(i) \cap t} [s(i) \subseteq t \Rightarrow a'_{i,j} = a_{i,j}]$  such that  $p$  is satisfied by the assignment  $x_{i,j} \mapsto a'_{i,j}$ .

PROOF. Use Fact 3.3B, and the observation: a set of formulas (not necessarily



over  $M$ ) of power  $< \kappa$  finitely satisfiable in a  $\kappa$ -compact model, is satisfiable in it.

3.3E. FACT. Suppose  $\langle M_s : s \in I \rangle$  is a stable system,  $t \subseteq \bigcup I$ , each  $M_s$  is  $\mathbf{F}_\kappa^a$ -saturated and  $a_{i,j} \in M_{s(i)}$  for  $i < \alpha$ ,  $j < j_\alpha$ ,  $\sum_{i < \alpha} j_i < \kappa$ , then we can find  $a'_{i,j} \in M_{s(i) \cap t}$  such that

$$\text{stp}_*(\langle a'_{i,j} : i < \alpha, j < j_\alpha \rangle, \{a_{i,j} : s(i) \subseteq t\}) = \text{stp}_*(\langle a_{i,j} : i < \alpha, j < j_\alpha \rangle, \{a_{i,j} : s(i) \subseteq t\}).$$

PROOF. Left to the reader.

3.3F. FACT. If  $S = \langle M_s : s \in I \rangle$  is a stable system, then for any  $t \in I$ ,  $\text{tp}_*(M_t, \bigcup \{M_s : t \not\subseteq s\})$  is definable over  $A_t^s$ .

This follows from 3.3(2).

3.4. THE MAIN THEOREM. Suppose  $T$  is stable and  $\text{cl}$  is a  $\Sigma_1$ -closure operation; then at least one of the following holds:

(A) If  $M_0 < M_1, M_2$ ,  $\{M_1, M_2\}$  independent over  $M_0$ , then  $\text{cl}(M_1 \cup M_2) < \mathfrak{U}$ .

(B) There is a set  $A = \text{cl } A$  such that the theory of  $\mathfrak{U} \upharpoonright A$  is unstable ( $\mathfrak{U} \upharpoonright A$  is the model  $\mathfrak{U}$  restricted to the set  $A$ , which by Definition 3.1(1) is closed under functions); moreover, the theory of  $\mathfrak{U} \upharpoonright A$  has the independence property (see [1] II §4). In fact, we can have  $A = \text{cl}(\bigcup_{i < i_0} M_i)$ ,  $i_0 \neq 0$ .

REMARKS. (1) If  $\text{cl}$  is as above, by 3.5(2) we can assume that in a counterexample to (A),  $M_1, M_2$  are isomorphic over  $M_0$ , hence get rid of the predicate  $P$  in the proof.

(2) Really we do not need the order  $<$  of  $L$ , but then we have to work a little more. It is also quite reasonable that we can replace stable by “without the independence property,” and then in (A) say “ $\text{tp}_*(M_2, M_1)$  is finitely satisfiable in  $M_0$ ,” but this was not checked.

(3) In conclusion 3.6(1) we shall show that when  $T$  is universal then (B) implies that some completion of  $T$  is unstable.

PROOF. Suppose  $M_0, M_1, M_2$  form a counterexample to (A), and we shall prove (B). Let  $\lambda = |T| + \|M_1\| + \|M_2\|$  and choose a model  $L = (\langle L \rangle, <, P, R)$ ,  $P$  a one place predicate,  $<$  a (linear) order,  $R$  a symmetric and reflexive two-place relation,  $L \models (\forall x, y)(P(x) \equiv P(y) \rightarrow xRy)$ , which is a  $\lambda$ -homogeneous and  $\lambda$ -universal (i.e. any isomorphism from one submodel of  $L$  onto another, both of power  $< \lambda$ , can be extended to an automorphism of  $L$ , and any model  $L'$  of power  $\leq \lambda$  satisfying the other conditions can be embedded into  $L$ ;  $L$  may have the power  $> \lambda$ ; see e.g. [1] I 1.8).

We shall now define for every  $s \in I \stackrel{\text{def}}{=} \{t : t \text{ a finite subset of } L\}$  a model  $M_s < \mathfrak{U}$  and isomorphism  $F_h$  for every  $h \in PI$  (see below),  $\text{Dom } h = s$ , such that:

(a)  $\langle M_s : s \in I \rangle$  is a stable system.

(b) Let

$$PI = \{h : \text{for some } s, t \in I, h \text{ is an isomorphism from } L \upharpoonright s \text{ on } L \upharpoonright t\}.$$

Now for any  $h \in PI$  there will be an isomorphism  $F_h$  from  $M_s$  onto  $M_t$  (where  $s = \text{Dom } h$ ,  $t = \text{Range}(h)$ ) such that:

(b,  $\alpha$ ) if  $f \subseteq h$  then  $F_f \subseteq F_h$ ,

(b,  $\beta$ ) if  $h_2 = h_0 h_1$  then  $F_{h_2} = F_{h_0} F_{h_1}$ ,

(b,  $\gamma$ ) if  $h$  is the identity on  $s = \text{Dom } h$ , then  $F_h$  is the identity mapping on  $M_s$ .

(c)  $M_\emptyset = M_0$ ,  $M_{\{u\}}$  is isomorphic to  $M_1$  over  $M_0$  if  $u \in P$  and  $M_{\{u\}}$  is isomorphic to  $M_2$  over  $M_0$  if  $u \notin P$  (note that  $\emptyset \in I$  and we deal exactly with the  $u \in L$ ).

We denote  $F_h^n = F_h \upharpoonright \cup \{M_s : s \subseteq \text{Dom } h, |s| \leq n\}$ . Now the definition is as follows: we define by induction on  $n$ ,  $M_s$  ( $s \in I, |s| = n$ ) and  $F_h^n$  ( $h \in PI$ ) such that (a), (b), (c) hold in the relevant cases (restricting to  $F_h^n$  when appropriate). In the end we shall let  $F_h = \bigcup_n F_h^n$ .

For  $n = 0$  use (c), for  $n = 1$  use (c) and the facts on forking (see III 0.1). For  $n > 1$  use 3.3: let  $\{t_i : i < i(*)\}$  be a list of all  $t \in I, |t| = n$ .

Now we define by induction on  $i < i(*)$  the model  $M_{t_i}$ . If there is no  $j < i$  such that  $L \upharpoonright t_i \cong L \upharpoonright t_j$ , choose  $M_{t_i}$  as any  $M < \mathfrak{U}$ ,

$$A_{t_i} \stackrel{\text{def}}{=} \bigcup_{\substack{s \subseteq t_i \\ s \neq t_i}} M_s \subseteq M, \quad \|M\| \leq |T| + |A_{t_i}| \leq \lambda,$$

and such that  $\text{tp}_*(M_{t_i}, \cup \{M_s : s \in I, |s| < n\} \cup \bigcup_{j < i} M_{t_j})$  does not fork over  $A_{t_i}$  (which is possible by the extension property of non-forking, see III 0.1). If there is  $j < i$  such that  $L \upharpoonright t_i \cong L \upharpoonright t_j$ , choose minimal  $j = j(i)$ , and let  $h_i^n$  be the isomorphism from  $L \upharpoonright t_i$  onto  $L \upharpoonright t_j$  (it is unique as  $<$  is a linear order of  $L$  and  $t_i$  is finite). Now there is an elementary mapping  $H_i^n$  extending  $F_{h_i}^{n-1}$  (see above) and whose domain is  $M_{t_{j(i)}}$  and  $\text{tp}_*(\text{Range}(H_i^n), \cup \{M_s : s \in I, |s| < n\} \cup \bigcup_{\alpha < i} M_{t_\alpha})$  does not fork over  $\text{Range}(F_{h_i}^{n-1})$ ; but note  $\text{Range}(F_{h_i}^{n-1}) = A_{t_j}$ . Now let  $M_{t_i} = \text{Range}(H_i^n)$ .

So we have defined all the  $M_{t_i}$  but still have to define  $F_h^n$  for  $h \in PI$ . Let  $\alpha(0) < \dots < \alpha(k-1)$  be a list of  $\{i : t_i \subseteq \text{Dom } h\}$ , and let  $t_{\beta(i)}$  be the range of  $h \upharpoonright t_{\alpha(i)}$ . Now we can define  $F_h^n$  as  $F_h^{n-1} \cup \bigcup_{i < k} H_{\beta(i)}^n (H_{\alpha(i)}^{n-1})^{-1}$ . It is easy to check that this is a well-defined one-to-one function with the suitable range and

domain.  $F_h^n$  is an elementary mapping by 3.3F and if  $g$  extends  $h$ , both in  $PI$  then  $F_g^n$  extend  $F_h^n$ . So we have defined  $M_s, F_h$  ( $s \in I, h \in PI$ ) as required.

Now let

$$J = \{s \in I : \text{for every } u \neq v \in s, L \models uRv\}$$

and for any  $S \subseteq L$ , let  $B_S = \text{cl} \cup \{M_s : s \subseteq S, s \in J\}$ . By the hypothesis (that  $M_0, M_1, M_2$  exemplify the failure of 3.4(A)), if  $u \in P, v \notin P, u, v \in L$  but not  $uRv$  then  $\mathcal{U} \upharpoonright B_{\{u,v\}}$  is not an elementary submodel of  $\mathcal{U}$ . By 3.3(2), the pair  $(B_{\{u,v\}}, B_L)$  satisfies the Tarski-Vaught condition (inside  $\mathcal{U}$ ). Hence  $\mathcal{U} \upharpoonright B_L$  is not an elementary submodel of  $\mathcal{U}$ . So there are  $\bar{a} \in B_L$ , and a first order formula  $\varphi$ , such that  $\mathcal{U} \models \varphi[\bar{a}]$ ,  $B_L \upharpoonright \mathcal{U} \models \neg \varphi[\bar{a}]$ , so for some  $n < \omega$  and finite  $t \subseteq L$ ,  $\varphi$  is a  $\Sigma_n$ -formula or  $\Pi_n$ -formula and  $\bar{a} \subseteq B_t$ . Among all possible  $\bar{a}, \varphi, n, t$  choose an example with minimal  $n$ , and for the fixed  $n$ , a minimal  $|t|$ , and for the minimal  $n$  and  $|t|$ , maximal  $|\{s : s \subseteq t, s \in J\}|$ . It is easy to prove that  $\varphi$  cannot be quantifier-free, nor a  $\Pi_n$ -formula, so  $n \geq 1$ ,  $\varphi(\bar{x}) = (\exists \bar{y})\psi(\bar{y}, \bar{x})$ ,  $\psi$  a  $\Pi_{n-1}$ -formula. It is also very easy to see that necessarily for some  $u, v \in t$ ,  $\neg uRv$ . Hence  $B_t = \text{cl}(B_{t-\{u\}} \cup B_{t-\{v\}})$ . So w.l.o.g. there are  $\bar{b} \in B_{t-\{v\}}, \bar{c} \in B_{t-\{u\}}$  so that  $\bar{a}$  is algebraic over  $\bar{b} \wedge \bar{c}$ , in fact for some  $\Sigma_1$ -formula  $\Theta, \models \Theta[\bar{a}, \bar{b}, \bar{c}]$  and  $\Theta(\bar{x}, \bar{b}, \bar{c})$  is algebraic. Let  $k$  be the number of  $\bar{a}'$  satisfying  $\Theta(\bar{x}, \bar{b}, \bar{c}) \wedge \varphi(\bar{x})$  (in  $\mathcal{U}$ ) and let  $\varphi^*(\bar{y}, \bar{z})$  be defined as

$$(\exists^{\approx k} \bar{x})(\Theta(\bar{x}, \bar{y}, \bar{z}) \wedge \varphi(\bar{x})).$$

Now for every set  $Y \subseteq \lambda \times \lambda$  we can define in  $L$  elements  $u_i, v_i$  ( $i < \lambda$ ) such that:

( $\alpha$ ) In  $L$ ,  $u, u_i$  realize the same quantifier-free type over  $t - \{u, v\}$ ; and similarly  $v, v_i$  realize the same quantifier-free type over  $t - \{u, v\}$ .

( $\beta$ )  $u_i R v_j$  holds iff  $\langle i, j \rangle \in Y$ .

( $\gamma$ ) If  $u < v$  then for every  $i, j < \lambda$ ,  $u_i < v_j$  (in  $L$ ) and if  $v < u$  then for every  $i, j < \lambda$ ,  $v_j < u_i$  (in  $L$ ).

Let  $g_i$  be the function with domain  $t - \{v\}$ ,  $g_i(u) = u_i$ ,  $g_i(u') = u'$  for  $u' \in t - \{u, v\}$ . Let  $h_i$  be the function with domain  $t - \{u\}$ ,  $h_i(v) = v_i$ ,  $h_i(v') = v'$ , for  $v' \in t - \{u, v\}$ . It is easy to check that  $g_i, h_i \in PI$  and  $g_i \cup h_j \in PI$  iff  $\neg u_i R v_j$  (iff  $\langle i, j \rangle \notin Y$ ). Let  $\bar{b}_i = F_{g_i}(\bar{b})$ ,  $\bar{c}_j = F_{h_j}(\bar{c})$ .

FACT A.  $\text{tp}(\bar{b} \wedge \bar{c}, \emptyset) = \text{tp}(\bar{b}_i \wedge \bar{c}_j, \emptyset)$  (in  $\mathcal{U}$ ).

This is because  $F_{g_i} \cup F_{h_j}$  is an elementary mapping (by 3.3F), and  $F_{g_i} \cup F_{h_j}(\bar{b} \wedge \bar{c}) = \bar{b}_i \wedge \bar{c}_j$ .

FACT B.  $(\mathcal{U} \upharpoonright B_L) \models \varphi^*[\bar{b}_i, \bar{c}_j]$  if  $u_i R v_j$ .

We have chosen  $k$  such that  $\mathcal{U} \models (\exists^{\leq k} \bar{x})[\Theta(\bar{x}, \bar{b}, \bar{c}) \times \varphi(\bar{x})]$ , i.e.  $\mathcal{U} \models \varphi^*[\bar{b}, \bar{c}]$ , hence by Fact A,  $\mathcal{U} \models \varphi^*[\bar{b}_i, \bar{c}_i]$ .

Now if  $\mathcal{U} \restriction B_L \models \neg \varphi^*[\bar{b}_i, \bar{c}_i]$ , then  $\bar{b}_i \wedge \bar{c}_i, \varphi^*, n, t^* = t \cup \{u_i, v_i\} - \{u, v\}$  contradict the choice of  $\bar{a}, \varphi, n, t$ , i.e.  $\varphi^*$  is a  $\Sigma_n$ -formula as  $\Theta$  is a  $\Sigma_1$ -formula,  $\varphi$  a  $\Sigma_n$ -formula and  $n > 0$ , also  $|t^*| = |t|$ , however the maximality of  $|\{s : s \subseteq t, s \in J\}|$  is contradicted.

FACT C.  $(\mathcal{U} \restriction B_L) \models \neg \varphi^*[\bar{b}_i, \bar{c}_i]$  if  $\neg u_i R v_i$ .

If  $\bar{a}' \in B_L$ ,  $\mathcal{U} \restriction B_L \models \Theta[\bar{a}', \bar{b}, \bar{c}] \wedge \varphi[\bar{a}']$  then  $\mathcal{U} \models \Theta[\bar{a}', \bar{b}, \bar{c}]$ ,  $\mathcal{U} \models \varphi[\bar{a}']$  (by the minimality of  $n$ , as  $\varphi$  is a  $\Sigma_n$ -formula; for  $\Theta$  — trivially). Hence the set of  $\bar{a}' \in B_L$  for which  $(\mathcal{U} \restriction B_L) \models \Theta[\bar{a}', \bar{b}, \bar{c}] \wedge \varphi[\bar{a}']$  is a subset of the set of  $\bar{a}' \in B_L$  for which  $\mathcal{U} \models \Theta[\bar{a}', \bar{b}, \bar{c}] \wedge \varphi[\bar{a}']$  which is a proper subset of the set of  $\bar{a}' \in \mathcal{U}$  for which  $\mathcal{U} \models \Theta[\bar{a}', \bar{b}, \bar{c}] \wedge \varphi[\bar{a}']$  (as witnessed by  $\bar{a}$ ). So we have just proved  $(\mathcal{U} \restriction B_L) \models \neg \varphi^*[\bar{b}, \bar{c}]$ . So it is enough to find an automorphism of  $\mathcal{U} \restriction B_L$  taking  $\bar{b} \wedge \bar{c}$  to  $\bar{b}_i \wedge \bar{c}_i$ . Now, as we have noted before,  $g_i \cup h_j \in PI$ , and  $F_{g_i \cup h_j}(\bar{b} \wedge \bar{c}) = \bar{b}_i \wedge \bar{c}_i$ . By choice of  $L$ ,  $g_i \cup h_j$  can be extended to an automorphism  $f$  of  $L$ . Now  $\{F_{f|t} : t \in PI\}$  is a directed family of elementary mapping, hence its union,  $F$ , is an elementary mapping. We shall prove that  $F^* = F \restriction B_L$  is as required.

(i)  $F^*(\bar{b} \wedge \bar{c}) = \bar{b}_i \wedge \bar{c}_i$  : as  $F_{g_i \cup h_j}$  belong to the family.

(ii)  $F$  maps  $\bigcup_{s \in J} M_s$  onto itself by the properties of the  $F_f$ 's, and as  $\text{cl}(\bigcup_{s \in J} M_s) \subseteq \text{Dom } F$ , clearly by the properties of  $\text{cl}$ ,  $F^*$  has to map  $\text{cl}(\bigcup_{s \in J} M_s)$  onto itself, so it is an automorphism of  $\mathcal{U} \restriction B_L$ .

So we have proved that

$$\mathcal{U} \restriction B_L \models \varphi^*[\bar{b}_i, \bar{c}_i] \quad \text{iff } \langle i, j \rangle \in Y$$

where  $Y \subseteq \lambda \times \lambda$  was arbitrary. So for some  $Y$ , we get that  $\mathcal{U} \restriction B_L$  has a theory with the independence properties.

We can get some more relevant facts.

3.5. CLAIM. (1) In 3.4 we can replace “ $\text{cl}$  is a  $\Sigma_1$ -closure operation” by “ $\text{cl}$  is a  $\Phi$ -closure operation,” provided that:

(\*) if  $A = \text{cl}(\bigcup_{i < \alpha} M_i)$ ,  $\bar{a} \in A$ ,  $\varphi(\bar{x}) \in \Phi$  and  $\mathcal{U} \restriction A \models \varphi[\bar{a}]$  then  $\mathcal{U} \models \varphi[\bar{a}]$ .

(2) If  $\text{cl}$  is local, then “ $(A, B)$  satisfies the Tarski–Vaught condition” implies “ $(\text{cl } A, \text{cl } B)$  satisfies the Tarski–Vaught condition.”

(3) If 3.4(A) holds then for every such  $M_0 < M_1, M_2$ ,  $\text{cl}(M_1 \cup M_2) = \text{acl}(M_1 \cup M_2)$ .

PROOF. (1) We replace in the proof  $\Sigma_n, \Pi_n$  by  $\Sigma'_n, \Pi'_n$  defined below by induction on  $n$ , and then the proof of 3.4 proves the assertion.

$\Sigma'_n$ : if  $n = 0$ ,  $\Sigma'_n$  is the set of quantifier-free formulas;  
 if  $n = 1$ ,  $\Sigma'_n$  is the set  $\Phi$  of formulas satisfying  $(*)$  of 3.5(1);  
 if  $n > 1$ ,  $\Sigma'_n$  is the closure of  $\Pi_{n-1}$  by conjunctions, disjunctions and existential quantifiers.

$\Pi'_n$  is the set of negations of  $\Sigma'_n$  formulas.

(2) Trivial.

(3) Trivial.

3.6. CONCLUSION. (1) If  $T$  is a stable universal theory (i.e., every completion of  $T$  is stable), then for every model  $\mathcal{U}$  of  $T$ , 3.4(A) holds.

(2) We can replace "universal" by  $\Sigma_2$  (i.e. every sentence of  $T$  is a  $\Sigma_2$ , or even  $\Sigma_2^{\text{al}}$  (defined below)).

3.7. DEFINITION. (1) Let  $\Pi_1^{\text{al}}$  be the closure of the family of quantifier-free formulas by: universal quantification, conjunctions, disjunctions, and  $(\exists^{1,k}x) \dots$  which mean: there are at least one but no more than  $k$   $\bar{x}$ 's satisfying  $\dots$ .

(2) Let  $\Sigma_2^{\text{al}} = \{(\exists \bar{y})\varphi(\bar{y}, \bar{x}) : \varphi(\bar{y}, \bar{x}) \in \Pi_1^{\text{al}}\}$ ,  $\Sigma_1^{\text{al}} = \{\neg \varphi : \varphi \in \Pi_1^{\text{al}}\}$ .

PROOF. So  $\text{Th}(\mathcal{U} \upharpoonright B_L)$  is unstable, so what? Here comes the main use of the assumption that the theory  $T$  is universal (this is the only use of that fact). Remember the definition of  $B_L$ : it was obtained by applying  $\text{cl}$  on a subset of  $\mathcal{U}$ . By part (1) (i) of Definition 3.1 and part (1) (ii) the set  $B_L$  is closed under the functions of the model  $\mathcal{U}$  hence it is a submodel of  $\mathcal{U}$ . But since  $T$  has  $\Pi_1$  axiomatization, really  $\mathcal{U} \upharpoonright B_L \models T$ . So we found a model of  $T$  which is unstable (even has the independence property), namely a completion of  $T (= \text{Th}(\mathcal{U} \upharpoonright B_L))$  which is not stable.

(2) By adding constants, translating the  $\Sigma_2$  to  $\Pi_1$  axioms, i.e., when applying 3.5, the set  $B_L$  such that  $\text{Th}(\mathcal{U} \upharpoonright B_L)$  has the independence property, satisfies:  $B_L = \text{cl } B_L$  and  $B_L$  extend some elementary submodel of  $\mathcal{U}$ . We can conclude that  $\mathcal{U} \upharpoonright B_L$  is a model of  $T$ .

The proof for  $\Sigma_2^{\text{al}}$  is similar.

## §4. Examples

### 4.1. Unstable

Let the language have one two-place function, the theory: empty. This  $T$  is a variety, not stable and  $I(\aleph_\alpha, T) = 2^{\aleph_\alpha}$ .

#### 4.2. *Stable not Superstable*

Let the language have  $\omega$  one-place function  $F_n$ , and  $T$  be

$$\{(\forall x)[F_n(F_m(x)) = F_m(F_n(x)) = F_n(x)] : n \leq m < \omega\}.$$

Clearly  $T$  is a variety stable not superstable,  $I(\aleph_\alpha, T) = 2^{\aleph_\alpha}$ . A “natural” model is  $(\omega^\omega, F_0, \dots, F_n, \dots)$ ,  $F_n(\eta) = \eta \upharpoonright n$ .

#### 4.3. *Superstable with the dop*

Let us define a model  $M$ : let  $G$  be an abelian group of order 2,  $|M| = G \times \omega$ , with the functions:

$$\begin{array}{ll} c & \langle 0, 0 \rangle \quad (c \text{ an individual constant}), \\ D & D(\langle a, n \rangle, \langle b, m \rangle, \langle c, k \rangle) = \langle a - b + c, 0 \rangle \quad (D \text{ a three-place function}), \\ F_k & F_k(\langle a, n \rangle) = \langle a, k \rangle \quad (F_k \text{ a one-place function for each } k < \omega). \end{array}$$

Let  $T$  be the set of sentences  $(\forall x_1 \dots)(\tau = \sigma)$ ,  $(\forall x_1 \dots)(\tau \neq \sigma)$  for terms  $\tau, \sigma$ , which  $M$  satisfies.

$T$  is universal, superstable with the dop.

#### 4.4. *Superstable without the dop, but Deep*

If  $L$  contains just one-place function  $F$ ,  $T$  is empty, then  $T$  is a variety and as required.

4.5. DEFINITION. Call  $T$  suitable if  $I(\aleph_\alpha, T) \geq |\alpha| + \aleph_0$ ,  $I(\aleph_\alpha, T) \leq I(\aleph_\beta, T)$  for  $\alpha \leq \beta$ .

4.6. CLAIM. (1) For universal  $T_i$  ( $i \leq \alpha$ ) there is a universal  $T$ ,  $|T| = \sum_{i < \alpha} |T_i|$  (the power of  $T_i$  is  $|L(T_i)| + \aleph_0$ ) such that: if each  $T_i$  is suitable then  $T$  is too and

$$I(\lambda, T) = \sum_{i < \alpha} I(\lambda, T_i).$$

(2) Similarly with  $I(\lambda, T) = \prod_{i < \alpha} I(\lambda, T_i)$ .

REMARK. (1) The suitability hypothesis is just to simplify the computation, and anyhow we here encounter only such  $T$ 's.

(2) E.g. in (2) we should of course write  $\text{Min}\{2^\lambda, \prod_{i < \alpha} I(\lambda, T_i)\}$ , but we shall ignore this.

PROOF. (1) W.l.o.g.  $L(T_i)$  ( $i < \alpha$ ) are pairwise disjoint. Let  $c_i$  ( $i < \alpha$ ) be new individual constants,  $P$  a monadic predicate and let  $T$  consist of the following

sentences:

$$\begin{array}{ll}
 c_i \neq c_\alpha \rightarrow c_j = c_\alpha & \text{for } i \neq j < \alpha. \\
 \neg P(c_i) & \\
 c_i = c_\alpha \rightarrow (\forall \bar{x}) R(\bar{x}) & R \in L(T_i) \text{ a predicate.} \\
 c_i = c_\alpha \rightarrow (\forall x_1, \dots) [F(x_1, \dots) = x_1] & F \in L(T_i) \text{ a function symbol.} \\
 c_i \neq c_\alpha \rightarrow \psi^P & \psi \in T_i.
 \end{array}$$

$$R(x_1, \dots, x_n) \rightarrow \bigwedge_{i=1}^n P(x_i) \quad R \in L(T_i) \text{ a predicate.}$$

$$F(x_1, x_2, \dots) \neq x_1 \rightarrow \bigwedge_{i=1}^n P(x_i) \quad F \in L(T_i) \text{ a function symbol.}$$

(2) Let  $\{P_i : i < \alpha\}$  be new monadic predicates.  $T$  will say: the  $P_i$ 's are pairwise disjoint,  $P_i$  is a model of  $T_i$ , and the predicates and function symbols of  $T_i$  are trivial when applied to elements not all of which are in  $P_i$ .

4.7. CLAIM. *For a universal  $T_0$  there is a universal  $T_1$ ,  $|T_1| = |T_0|$ , such that if  $T_0$  is suitable,  $T_1$  is suitable too; and*

$$I(\lambda, T_1) = 2^{I(\lambda T_0)}.$$

PROOF. W.l.o.g.  $T_0$  has no individual constants (replace  $c$  by  $F_c(x)$ , adding  $(\forall x, y) [F_c(x) = F_c(y)]$  to  $T_0$ ).  $T_1$  "says":  $E$  is an equivalence relation, each equivalence class is a model of  $T_0$ , and

$$\begin{aligned}
 P(x_1, \dots, x_n) &\rightarrow \bigwedge_{i=1}^n x_i E x_1, \\
 F(x_1, \dots, x_n) \neq x_1 &\rightarrow \bigwedge_{i=1}^n x_i E x_1
 \end{aligned}$$

for every predicate  $P$  and function symbol  $F$  of  $T_0$ .

4.8. CLAIM. *Claims 4.6 and 4.7 hold for "quasi-varieties" and for "totally transcendental universal" and for "total transcendental quasi-varieties" instead of "universal" provided that for the quasi-variety cases the language has no function symbols.*

PROOF. The problem is taking care of the "quasi-variety".

Case I. 4.6(1) We assume w.l.o.g.  $L(T_i)$  are pairwise disjoint and with no individual constant. Let  $P_i$  ( $i < \alpha$ ) be new monadic predicates.  $T$  will say:

$$(a) \quad P_i(x) \wedge P_j(y) \rightarrow \varphi \quad \text{for } i \neq j, \varphi \text{ an atomic formula,}$$

- (b)  $P_i(x) \rightarrow P_i(y)$  for  $i < \alpha$ ,  
 (c)  $R(x_1, \dots, x_n) \rightarrow P_i(y)$  for  $R$  a predicate of  $L(T_i)$ ,  
 (d)  $\psi$   $\psi \in T_i$  for some  $i < \alpha$ .

Of course  $L(T) = \bigcup_{i < \alpha} L(T_i) \cup \{P_i : i < \alpha\}$ . Let  $M$  be a model of  $T$ . If  $P_i^M = \emptyset$  for  $i < \alpha$  by (c), (d) all the relations of  $M$  are empty so  $M$  is determined up to isomorphism by its cardinality.

If  $P_{i(0)}^M \neq \emptyset$ ,  $P_{i(1)}^M \neq \emptyset$  and  $i(0) \neq i(1)$  the model  $M$  is again determined up to isomorphism by its cardinality: by (a).

Lastly suppose  $P_{i(0)}^M \neq \emptyset$ ,  $P_j^M = \emptyset$  for  $j \neq i(0)$ . By (b)  $P_{i(0)}^M \models |M|$ , by (c)  $R^M = \emptyset$  for  $R \in L(T_j)$ ,  $j \neq i(0)$ , and by (d)  $M \upharpoonright L(T_{i(0)})$  is a model of  $T_{i(0)}$ . As in the cases with inclusion we can conclude:

$$I(\lambda, T) = 2 + \sum_{i < \alpha} I(\lambda T_i).$$

*Case II.* 4.6(2) Let  $T$  say:

- (a)  $P_i(x) \times P_j(x) \rightarrow \varphi$ ,  $i < j < \alpha$ ,  $\varphi$  any atomic formula,  
 (b)  $R(x_1, \dots, x_n) \rightarrow P_i(x_l)$ ,  $R$  an  $n$ -place predicate is  $L(T_i)$ ,  $l \in \{1, \dots, n\}$ ,  
 (c)  $n \wedge_m P_i(x_m) \wedge \bigwedge_{i=1}^m \psi_i \rightarrow \psi$ ,  $\bigwedge_{i=1}^n \psi_i \rightarrow \psi$  an axiom of  $T_i$ .

*Case III.* 4.7 Similar.

REMARK. We can add to the definition of *suitable*

- (iii)  $I(\lambda, T \cup \{c = c\}) = I(\lambda, T)$  for  $c$  a new individual constant.

Then also for quasi-varieties we can allow the languages to have individual constants and get similar results.

4.9. CLAIM. (1) *There is a variety  $T_\kappa$  ( $\kappa \geq \aleph_0$  a cardinal)  $|T_\kappa| = \kappa$ ,  $I(\aleph_\alpha, T) = (|\alpha| + \aleph_0)^{2^\kappa}$ .*

(2) *There is a variety  $T$ ,  $I(\aleph_\alpha, T)$  is 1 for  $\alpha > 0$  and  $\aleph_0$  for  $\alpha = 0$ ,  $|T| = \aleph_0$ .*

(3) *There is a variety  $T$ ,  $I(\aleph_\alpha, T) = 1$ ,  $|T| = \aleph_0$ .*

(4) *For any  $n < \omega$ , there is a variety  $T$ ,  $I(\aleph_\alpha, T) = \beth_n(|\alpha| + \aleph_0)$ .*

PROOF. (1) Let the language have  $\kappa$  one-place functions  $F_i$  ( $i < \kappa$ ), and

$$T = \{(\forall x) F_i(F_i(x)) = x : i < \kappa\}.$$

(2) The axioms of  $T$  are those of a vector space over the rationals (i.e., for each rational number there is a one-place function symbol for multiplication by it and of course we have the addition function).

(3)  $T$  is empty.



(4) When  $n > 0$  let  $T$  be  $(\forall x)F^{n+1}(x) = c$ . For  $n = 0$  let  $T = \{(\forall x)F(F(x)) = x\}$ . A model of  $T$ ,  $M = \langle |M|, F^M \rangle$  is determined up to isomorphism by the following two cardinals:  $|\{a \in M : F(a) = a\}|$  and  $|\{a \in M : F(a) \neq a\}|$ .

4.10. CLAIM. (1) *There is a countable variety  $T$ ,  $I(\aleph_\alpha, T) = \beth_m((|\alpha| + \aleph_0)^{\aleph_0})$ .*

(2) *There is a countable variety  $T$ ,  $I(\aleph_\alpha, T) = \beth_m((|\alpha| + \aleph_0)^{\aleph_0})$ .*

PROOF. (1) For  $m = 0$ , let  $T$  consist of

$$(\forall x)F_n(F_n(x)) = x \quad (\text{for } n < \omega).$$

For  $m > 0$  let  $T$  consist of  $(G, F_n$  are unary functions):

$$G^{m+2}(x) = G^{m+1}(x),$$

$$F_n(F_n(x)) = x,$$

$$F_n(G(x)) = G(x),$$

$$G(F_n(x)) = G(x) \text{ for } n < \omega.$$

(2) First let  $m = 0$ , and  $T_0$  will consist of  $(\forall x)[F(F^{-1}(x)) = x = F^{-1}(F(x))]$  (i.e., essentially one unary one-to-one function). So a model  $M$  is characterized by the cardinals

$$\lambda_n(M) = |\{x \in M : n \geq 1 \text{ is minimal such that } F^n(x) = x\} / E|$$

where

$$xEy \stackrel{\text{def}}{=} (\exists n) \left( x = F^n(y) \vee y = F^n(x) \right).$$

( $E$  is an equivalence relation by our axioms,  $F^n$  is the  $n$ th power of  $F$ .)

For  $m > 0$ ,  $T = T_m$  will consist of  $(F, G, G^{-1}$  are unary functions)

$$F^{m+1}(x) = F^{m+2}(x),$$

$$G(F(x)) = F(G(x)),$$

$$F(G(x)) = F(x),$$

$$G(G^{-1}(x)) = x = G^{-1}(G(x)).$$

Note that if  $M \models T$ ,  $A_i = \text{Range}(F^i)$ , then  $G$  is the identity on  $A_1$ : for  $x \in A_1$ , there is  $y$ ,  $x = F(y)$  so  $x = F(y) = F(G(y)) = G(F(y)) = G(x)$ .

4.11. CONJECTURE. For every variety  $T$ , either  $I(\lambda, T) = 2^\lambda$  for  $\lambda > |T|$  or  $I(\aleph_\alpha, T) < \beth_\omega(|\alpha| + |T|)$  for every  $\alpha$ .

## §5. Counterexample

5.1. EXAMPLE. For each  $n < \omega$  there is a universal theory  $T_n$ , which has an unstable completion, but if  $\varphi(\bar{x}, \bar{y})$  is an unstable formula in such a completion, then  $\varphi$  is not  $\Sigma_n$ .

We first prove:

5.2. CLAIM. We can define by induction on  $n$  a theory  $T_n$  such that:

(A)  $T_n$  is universal, countable  $\aleph_0$ -stable, and in fact every completion of it has a finite depth (see [2, §4]), in fact depth  $n$ , and  $I(\aleph_n, T_n) = \beth_n(|\alpha| + \aleph_0)$ .

(B)  $T_n$  is complete for  $\bigcup_{m < n} (\Sigma_m \cup \Pi_m)$  sentences.

(C)  $T_n$  is not complete for  $\Sigma_n$ -sentences, but only sentences from  $\Sigma_n \cup \Pi_m$  are needed for completion, i.e. if  $T$  is a completion of  $T_n$  then

$$T_n \cup \{\varphi : \varphi \in T, \varphi \text{ a } \Sigma_n \text{ or } \Pi_n \text{ sentence}\}$$

is complete.

(D) For every completion  $T$  of  $T_n$ , there is a complete universal countable  $\aleph_0$ -stable theory  $S(T)$  of depth  $n$  (maybe in a larger language),  $T \subseteq S(T)$ ,  $S(T)$  as in (A), such that  $I(\aleph_n, S(T)) \leq \beth_n(|\alpha| + \aleph_0)$ .

(E)  $T_n^* = T_n \cup \{\Theta : \Theta \text{ a } \Sigma_n\text{-sentence consistent with } T_n\}$  is a consistent theory.

(F)  $L_n$ , the language of  $T_n$ , as well as the languages of  $S(T)$  from (D), consists of predicates and one-place functions only.

(G) The language of  $S(T)$  (for  $T$  a completion of  $T_n$ ) is  $L_n^+$  (the language does not depend on the theory  $T$ , only on  $n$ ), and for every sentence  $\psi \in L(T_n)$  there is a sentence  $\Theta_\psi^* \in L_n$  such that for every completion  $T$  of  $T_n$   $[\psi \in T \leftrightarrow \Theta_\psi^* \in S(T)]$ .

PROOF. For  $n = 1$ . We let  $L_0$  contain the equality sign only, and  $T_0 = \{(\forall x)x = x\}$ .

Clearly  $T_n^*$  is (equivalent to)  $\{(\exists x_1 \cdots x_k) \wedge_{i < j} x_i \neq x_j : k < \omega\}$  (hence  $T_n^*$  is consistent, i.e. (E) holds). Also any other completion of  $T_0$  is the theory of a finite model, i.e. it is

$$T_0^l = \left\{ \left( (\exists x_1 \cdots x_l) \wedge_{i < j} x_i \neq x_j \right) \cup \left\{ (\forall x_1 \cdots x_{l+1}) \vee_{i < j} x_i = x_j \right\} \right\} \quad \text{for some } l.$$

It is now easy to prove (D): e.g., for  $T = T_0^k$

$$S(T) = \{(\forall x)(\forall y)F_n(x) = F_n(y) : n < k\}$$

$$\cup \left\{ (\forall x_1, \dots, x_{l+1}), \vee_{i < j} x_i = x_j \right\}$$

$$\cup \{(\forall x)F_n(x) \neq F_m(x) : n < m < k\}.$$

So  $L_n^+ = L_n \cup \{F_n : n < \omega\}$ . (We use one place functions instead of individual constants just for the convenience of the induction, formally — so that (F) holds.) The other parts are very easy, too.

For  $n+1$ . So  $T_n, L_n, T_n^*, S(T), L_n^+$  are defined and satisfy (A)–(F).

Let  $S^l$  ( $l < \omega$ ) be each a copy of some theories  $S(T)$  such that every sentence in  $L_n$  consistent with  $T_n$  appears in infinitely many  $S^l$ , their languages are disjoint outside  $L_n$ , and  $F_n^m, E_n \in S^l$  ( $l, m < \omega$ ). Now  $L_{n+1}$  consists of  $L_n \cup \{E_n\} \cup \{F_l^i : l < \omega\} \cup \bigcup_{l < \omega} L(S^l)$ ,  $E_n$  a two place relation,  $F_n^m$  a one place function symbol (so it is still countable).  $T_{n+1}$  will consist of sentences saying the following:

(i)  $E$  is an equivalence relation, and for any function symbol  $F$  of  $L_n$  (necessarily one place by (F)), and the choice of  $F_l^i$

$$(\forall x)[xE_n F(x)].$$

(ii) Each  $E_n$ -equivalence class is a model of  $T_n$ , i.e. if  $\forall x_1 \cdots x_{k-1} \psi \in T_n$  ( $\psi$  quantifier-free) then

$$(\forall x_0 \cdots x_{k-1}) \left( \bigwedge_{l < k} x_0 E_n x_l \rightarrow \psi \right).$$

(iii)  $F_n^l$  is really a hidden individual constant, i.e.

$$\forall x \forall y (F_n^l(x) = F_n^l(y))$$

and those “constants” are not  $E_n$ -equivalent, i.e.

$$(\forall x)[\neg F_n^l(x) E_n F_n^k(x)] \quad \text{for } l < k < \omega.$$

(iv) The  $E_n$ -equivalence class of  $F_n^l(x)$  is a model of  $S^l$ . It is easy to check that  $T_{n+1}$  is as required.

(v) Every function symbol of  $S^l$  not in  $L_n$  is the identity outside the equivalence class which is a model of  $S^l$  (and similarly for other non-logical symbols).

REMARK. In 5.2(A) we could add that  $T_n$  has only countably many completions.

PROOF OF 5.1. Let  $P$  be a place predicate,  $F$  a two-place function,  $E$  a two-place predicate, all not in  $T_n$  from 5.2. Now  $T$  will say:

(a)  $F$  is a pairing function from  $P$  into  $\neg P$ :

$$(\forall xy) \left[ \neg P(x) \vee \neg P(y) \rightarrow F(x, y) = x \right],$$

$$\begin{aligned}
 & (\forall xy) \left[ P(x) \wedge P(y) \rightarrow \neg P(F(x, y)) \right], \\
 & (\forall x_1 x_2 y_1 y_2) \left[ P(x_1) \wedge P(y_1) \wedge P(x_2) \wedge P(y_2) \rightarrow \right. \\
 & \quad \left. \left[ F(x_1, y_1) = F(x_2, y_2) \equiv (x_1 = x_2 \wedge y_1 = y_2) \right] \right].
 \end{aligned}$$

(b)  $E$  is an equivalence relation on  $\neg P$ , and each equivalence class contains at most one  $F(x, y)$ .

(c) Each equivalence class is a model of  $T_n$  and the unary functions of  $L(T_n)$  are the identity on  $\neg P$ .

(d) If  $\neg P(F(x, y))$ , then  $F(x, y)$  is  $F_{n-1}^0(F(x, y))$ , i.e. it is one of individual constants of  $T_n$  which we have hidden for technical reasons as one place functions.

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